

CP/M-86[®]
Operating
System
System Guide



Texas Instruments Professional Computer

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CP/M-86® Operating System System Guide
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Foreword

The *CP/M-86 Operating System System Guide* presents the system programming aspects of CP/M-86®, a single-user operating system for the Intel® 8086 and 8088 16-bit microprocessors. The discussion assumes that the reader is familiar with CP/M®, the Digital Research eight-bit operating system. To clarify specific differences with CP/M-86, this document refers to the eight-bit version of CP/M as CP/M-80™. Elements common to both systems are simply called CP/M features.

This system guide presents an overview of the conventions of the CP/M-86 programming interface. It also describes procedures for adapting CP/M-86 to a custom hardware environment.

Section 1 gives an overview of CP/M-86 and summarizes the differences between it and CP/M-80. Section 2 describes the general execution environment, while Section 3 tells how to generate command files. Sections 4 and 5 respectively define the programming interfaces to the Basic Disk Operating System (BDOS) and the Basic Input/Output System (BIOS). Section 6 discusses alteration of the BIOS to support custom disk configurations, and Section 7 describes the loading operation and the organization of the CP/M-86 system file.

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CP/M-86 System Overview

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GENERAL CHARACTERISTICS OF CP/M-86

CP/M-86 contains all of the facilities of CP/M-80 with additional features to account for increased processor address space of up to a megabyte (1,048,576 bytes) of main memory. CP/M-86 also maintains file compatibility with all previous versions of CP/M. The file structure of Version 2 of CP/M is used, allowing as many as sixteen drives with up to eight megabytes on each drive. Because of this, CP/M-80 and CP/M-86 systems may exchange files without modifications of the file format.

CP/M-86 resides in the file CPM.SYS, which is loaded into memory by a cold-start loader during system initialization. The cold-start loader resides on the first track of the system disk. CPM.SYS contains three program modules: the Console Command Processor (CCP), the Basic Disk Operating System (BDOS), and the user-configurable Basic I/O System (BIOS). The CCP and BDOS portions occupy approximately 10K bytes, while the size of the BIOS varies with the implementation. The operating system executes in any portion of memory above the reserved interrupt locations, while the remainder of the address space is partitioned into as many as eight non-contiguous regions (as defined in a BIOS table). Unlike CP/M-80, the CCP area cannot be used as a data area subsequent to transient program load; all CP/M-86 modules remain in memory at all times, and are not reloaded during a warm start.

NOTE

One K equals 1,024 bytes.

In a way similar to CP/M-80, CP/M-86 loads and executes memory image files from disk. Memory image files are preceded by a header record, defined in this document, which provides information required for proper program loading and execution. Memory image files under CP/M-86 are identified by a CMD file type.

Unlike CP/M-80, CP/M-86 does not use absolute locations for system entry or default variables. The BDOS entry takes place through a reserved software interrupt, while entry to the BIOS is provided by a new BDOS call. Two variables maintained in low memory under CP/M-80, the default disk number and I/O byte, are placed in the CCP and BIOS, respectively. Dependence upon absolute addresses is minimized in CP/M-86 by maintaining initial base page values, such as the default FCB and default command buffer, in the transient program data area.

Utility programs such as ED, PIP, STAT and SUBMIT operate in the same manner under both CP/M-86 and CP/M-80. In its operation, DDT-86TM resembles DDTTM supplied with CP/M-80. It allows interactive debugging of 8086 and 8088 machine code. Similarly, ASM-86TM allows assembly language programming and development for the 8086 and 8088 using Intel-like mnemonics.

The GENCMD (Generate CMD) utility corresponds to the LOAD program of CP/M-80, and converts the hexadecimal (hex) files produced by ASM-86 or Intel utilities into memory image format suitable for execution under CP/M-86. The LDCOPY (Loader Copy) program replaces SYSGEN, and is used to copy the cold start loader from a system disk for replication. A variation of GENCMD, called LMCMD, converts output from the Intel LOC86 utility into CMD format. Finally, GENDEF (Generate DISKDEF) is provided as an aid in producing custom disk parameter tables. ASM-86, GENCMD, LMCMD, and GENDEF are also supplied in COM file format for cross-development under CP/M-80.

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Several terms used throughout this manual are defined in the following table:

CP/M-86 Terms

Term	Meaning
Nibble	4-bit half-byte
Byte	8-bit value
Word	16-bit value
Double Word	32-bit value
Paragraph	16 contiguous bytes
Paragraph Boundary	An address divisible evenly by 16 (low order nibble 0)
Segment	Up to 64K contiguous bytes
Segment Register	One of CS, DS, ES, or SS
Offset	16-bit displacement from a segment register
Group	A segment-register-relative relocatable program unit
Address	The effective memory address derived from the composition of a segment register value with an offset value

A group consists of segments that are loaded into memory as a single unit. Since a group may consist of more than 64K bytes, it is the responsibility of the application program to manage segment registers when code or data beyond the first 64K segment is accessed.

CP/M-86 supports eight program groups: the code, data, stack and extra groups as well as four auxiliary groups. When a code, data, stack or extra group is loaded, CP/M-86 sets the respective segment register (CS, DS, SS or ES) to the base of the group. CP/M-86 can also load four auxiliary groups. A transient program manages the location of the auxiliary groups using values stored by CP/M-86 in the user's base page.

DIFFERENCES BETWEEN CP/M-80 AND CP/M-86

The structure of CP/M-86 is as close to CP/M-80 as possible in order to provide a familiar programming environment which allows application programs to be transported to the 8086 and 8088 processors with minimum effort. This section points out the specific differences between CP/M-80 and CP/M-86 in order to reduce your time in scanning this manual if you are already familiar with CP/M-80. The terms and concepts presented in this section are explained in detail throughout this manual, so you will need to refer to the table of contents to find relevant sections which provide specific definitions and information.

Due to the nature of the 8086 processor, the fundamental difference between CP/M-80 and CP/M-86 is found in the management of the various relocatable groups. Although CP/M-80 references absolute memory locations by necessity, CP/M-86 takes advantage of the static relocation capabilities inherent in the 8086 processor. The operating system itself is usually loaded directly above the interrupt locations, at location 0400H, and relocatable transient programs load in the best-fit memory region. However, you can load CP/M-86 into any portion of memory without changing the operating system (for this reason, there is no MOVCPM utility with CP/M-86), and transient programs will load and run in any non-reserved region.

Three general memory models are presented below: the 8080 Model, the Small Model, and the Compact Model. If you are converting 8080 programs to CP/M-86, you can use either the 8080 Model or the Small Model, and leave the Compact Model for a time when your addressing needs increase. You will use GENCMD, described in Operation of GENCMD, Chapter 3, to produce an executable program file from a hex file. GENCMD parameters allow you to specify which memory model your program requires.

CP/M-86 itself is constructed as an 8080 Model. This means that all of the segment registers are placed at the base of CP/M-86, and your customized BIOS is identical, in most respects, to that of CP/M-80 (with changes in instruction mnemonics, of course). In fact, the only additions are found in the SETDMAB, GETSEGB, SETIOB, and GETIOB entry points in the BIOS. Your warm-start subroutine is simpler, since you are not required to reload the CCP and BDOS under CP/M-86. One other point: if you implement the IOBYTE facility, you will have to define the variable in your BIOS. Taking these changes into account, you need only perform a simple translation of your CP/M-80 BIOS into 8086 code in order to implement your 8086 BIOS.

If you have implemented CP/M-80 Version 2, you already have disk definition tables which will operate properly with CP/M-86. You may wish to attach different disk drives, or experiment with sector skew factors to increase performance. If so, you can use the new GENDEF utility which performs the same function as the DISKDEF macro used by MAC under CP/M-80. You will find, however, that GENDEF provides you with more information and checks error conditions better than the DISKDEF macro.

Although generating a CP/M-86 system is generally easier than generating a CP/M-80 system, complications arise if you are using single-density floppy disks. CP/M-86 is too large to fit in the two-track system area of a single-density disk, so the bootstrap operation must perform two steps to load CP/M-86: first the bootstrap must load the cold start loader, then the cold start loader loads CP/M-86 from a system file. The cold start loader includes a LDBIOS which is identical to your CP/M-86 BIOS with the exception of the INIT entry point. You can simplify the LDBIOS if you wish, because the loader need not write to the disk. If you have a double-density disk, or reserve enough tracks on a single-density disk, you can load CP/M-86 without a two-step boot.

To make a BDOS system call, use the reserved software interrupt #244. The jump to the BDOS at location 0005 found in CP/M-80 is not present in CP/M-86. However, the address field at offset 0006 is present so that programs which size available memory using this word value will operate without change. CP/M-80 BDOS functions use certain 8080 registers for entry parameters and returned values. CP/M-86 BDOS functions use a table of corresponding 8086 registers. For example, the 8086 registers CH and CL correspond to the 8080 registers B and C. Look through the list of BDOS function numbers in Chapter 4, and you will find that functions 0, 27, and 31 have changed slightly. Several new functions have been added, but they do not affect existing programs.

One major philosophical difference is that in CP/M-80, all addresses sent to the BDOS are simply 16-bit values in the range 0000H to 0FFFFH. In CP/M-86, however, the addresses are really just 16-bit offsets from the DS (Data Segment) register which is set to the base of your data area. If you translate an existing CP/M-80 program to the CP/M-86 environment, your data segment will be less than 64K bytes. In this case, the DS register need not be changed following initial load, and thus all CP/M-80 addresses become simple DS-relative offsets in CP/M-86.

Under CP/M-80, programs terminate in one of three ways: by returning directly to the CCP, by calling BDOS function 0, or by transferring control to absolute location 0000H. CP/M-86, however, supports only the first two methods of program termination. This has the side effect of not providing the automatic disk system reset following the jump to 0000H which, instead, is accomplished by entering a CTRL at the CCP level.

You will find many new facilities in CP/M-86 that will simplify your programming and expand your application programming capability. However, we have designed CP/M-86 to make it easy to get started: in short, if you are converting from CP/M-80 to CP/M-86, there will be no major changes beyond the translation to 8086 machine code. Further programs you design for CP/M-86 are upward-compatible with MP/M-86™, the Digital Research multitasking operating system, as well as CP/NET-86™ which provides a distributed operating system in a network environment.

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Command Setup and Execution Under CP/M-86

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INTRODUCTION

This section discusses the operation of the Console Command Processor (CCP), the format of transient programs, CP/M-86 memory models, and memory image formats.

CCP BUILT-IN AND TRANSIENT COMMANDS

The operation of the CP/M-86 CCP is similar to that of CP/M-80. Upon initial cold start, the CP/M sign-on message is printed, drive A is automatically logged in, and the standard prompt is issued at the console. CP/M-86 then waits for input command lines from the console, which may include one of the following built-in commands:

DIR ERA REN TYPE USER

Note that SAVE is not supported under CP/M-86, since the equivalent function is performed by DDT-86.

Alternatively, the command line may begin with the name of a transient program with the assumed file type CMD denoting a command file. The CMD file type differentiates transient command files used under CP/M-86 from COM files which operate under CP/M-80.

The CCP allows multiple programs to reside in memory, providing facilities for background tasks. A transient program such as a debugger may load additional programs for execution under its own control. Thus, for example, a background printer spooler could first be loaded, followed by an execution of DDT-86. DDT-86 may, in turn, load a test program for a debugging session and transfer control to the test program between breakpoints. CP/M-86 keeps account of the order in which programs are loaded and, upon encountering **CTRL-C**, discontinues execution of the most recent program activated at the CCP level. **CTRL-C** at the DDT-86 command level aborts DDT-86 and its test program. A second **CTRL-C** at the CCP level aborts the background printer spooler. A third **CTRL-C** resets the disk system. Note that program abort due to

CTRL-C does not reset the disk system, as is the case in CP/M-80. A disk reset does not occur unless the **CTRL-C** occurs at the CCP command input level with no programs residing in memory.

When CP/M-86 receives a request to load a transient program from the CCP or another transient program, it checks the program's memory requirements. If sufficient memory is available, CP/M-86 assigns the required amount of memory to the program and loads the program. Once loaded, the program can request additional memory from the BDOS for buffer space. When the program is terminated, CP/M-86 frees both the program memory area and any additional buffer space.

TRANSIENT PROGRAM EXECUTION MODELS

The initial values of the segment registers are determined by one of three memory models used by the transient program and described in the CMD file header. The three memory models are summarized in the following table:

CP/M-86 Memory Models

Model	Group Relationships
8080 Model	Code and Data Groups Overlap
Small Model	Independent Code and Data Groups
Compact Model	Three or More Independent Groups

The 8080 Model supports programs which are directly translated from CP/M-80 when code and data areas are intermixed. The 8080 model consists of one group which contains all the code, data, and stack areas. Segment registers are initialized to the starting address of the region containing this group. The segment registers can, however, be managed by the application program during execution so that multiple segments within the code group can be addressed.

The Small Model is similar to that defined by Intel, where the program consists of an independent code group and a data group. The Small Model is suitable for use by programs taken from CP/M-80 where code and data is easily separated. Note again that the code and data groups often consist of, but are not restricted to, single 64K byte segments.

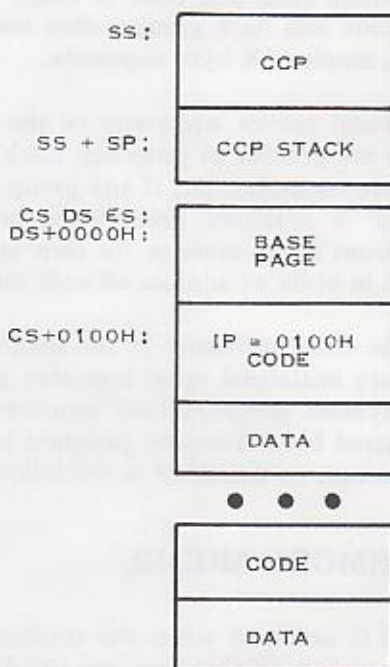
The Compact Model occurs when any of the extra, stack, or auxiliary groups are present in program. Each group may consist of one or more segments, but if any group exceeds one segment in size, or if auxiliary groups are present, then the application program must manage its own segment registers during execution in order to address all code and data areas.

The three models differ primarily in the manner in which segment registers are initialized upon transient program loading. The operating system program load function determines the memory model used by a transient program by examining the program group usage, as described in the following sections.

THE 8080 MEMORY MODEL

The 8080 Model is assumed when the transient program contains only a code group. In this case, the CS, DS, and ES registers are initialized to the beginning of the code group, while the SS and SP registers remain set to a 96-byte stack area in the CCP. The Instruction Pointer Register (IP) is set to 100H, similar to CP/M-80 thus allowing base page values at the beginning of the code group. Following program load, the 8080

Model appears as shown in the following figure, where low addresses are shown at the top of the diagram:



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The intermixed code and data regions are indistinguishable. The following base page values, are identical to CP/M-80, allowing simple translation from 8080, 8085, or Z80® code into the 8086 and 8088 environment. The following ASM-86 example shows how to code an 8080 model transient program.

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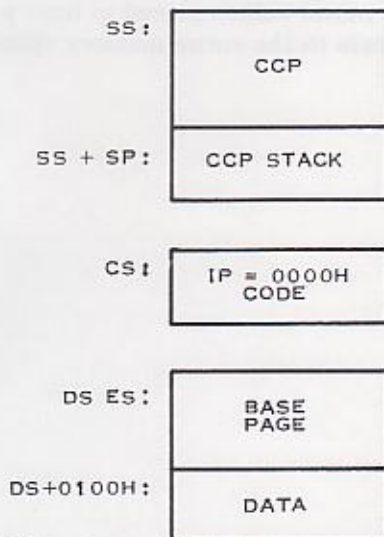
```

    eseg
    org     100h
    .
    .
    .       (code)
endcs     equ     $
    dseg
    org     offset endcs
    .
    .
    .       (data)
    end

```

THE SMALL MEMORY MODEL

The Small Model is assumed when the transient program contains both a code and data group. (In ASM-86, all code is generated following a CSEG directive, while data is defined following a DSEG directive with the origin of the data segment independent of the code segment.) In this model, CS is set to the beginning of the code group, the DS and ES are set to the start of the data group, and the SS and SP registers remain in the CCP's stack area as shown in the following figure:



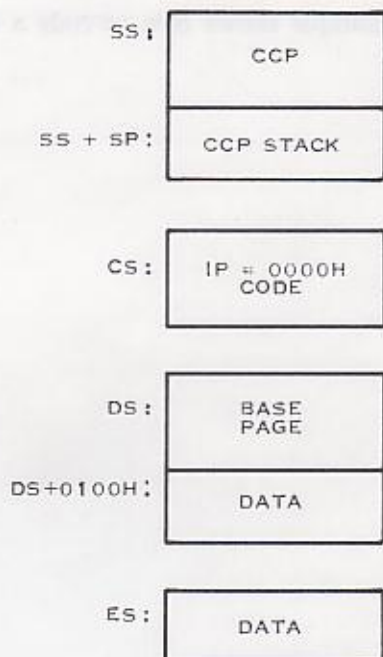
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The machine code begins at CS + 0000H, the base page values begin at DS + 0000H, and the data area starts at DS + 0100H. The following ASM-86 example shows how to code a small model transient program.

```
cseg
.
.      (code)
dseg
org    100h
.
.
.      (data)
end
```

THE COMPACT MEMORY MODEL

The Compact Model is assumed when code and data groups are present, along with one or more of the remaining stack, extra, or auxiliary groups. In this case, the CS, DS, and ES registers are set to the base addresses of their respective areas. The following figure shows the initial configuration of segment registers in the Compact Model. The values of the various segment registers can be programmatically changed during execution by loading from the initial values placed in base page by the CCP, thus allowing access to the entire memory space.



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If the transient program intends to use the stack group as a stack area, the SS and SP registers must be set upon entry. The SS and SP registers remain in the CCP area, even if a stack group is defined. Although it may appear that the SS and SP registers should be set to address the stack group, there are two contradictions. First, the transient program may be using the stack group as a data area. In that case, the Far Call instruction used by the CCP to transfer control to the transient program could overwrite data in the stack area. Second, the SS register would logically be set to the base of the group, while the SP would be set to the offset of the end of the group. However, if the stack group exceeds 64K, the address range from the base to the end of the group exceeds a 16-bit offset value.

The following ASM-86 example shows how to code a compact model transient program.

```
cseg
.
.
.      (code)
dseg
org   100h
.
.
.      (data)
eseg
.
.
.      (more data)
sseg
.
.
.      (stack area)
end
```

BASE PAGE INITIALIZATION

In a manner similar to CP/M-80, the CP/M-86 base page contains default values and locations initialized by the CCP and used by the transient program. The base page occupies the regions from offset 0000H through 00FFH relative to the DS register. The values in the base page for CP/M-86 include those of CP/M-80, and appear in the same relative positions, as shown in the following figure:

DS + 0000 :	LC0	LC1	LC2
DS + 0003 :	BC0	BC1	M80
DS + 0006 :	LD0	LD1	LD2
DS + 0009 :	BD0	BD1	XXX
DS + 000C :	LE0	LE1	LE2
DS + 000F :	BE0	BE1	XXX
DS + 0012 :	LS0	LS1	LS2
DS + 0015 :	BS0	BS1	XXX
DS + 0018 :	LX0	LX1	LX2
DS + 001B :	BX0	BX1	XXX
DS + 001E :	LX0	LX1	LX2
DS + 0021 :	BX0	BX1	XXX
DS + 0024 :	LX0	LX1	LX2
DS + 0027 :	BX0	BX1	XXX
DS + 002A :	LX0	LX1	LX2
DS + 002D :	BX0	BX1	XXX

DS + 0030 :
DS + 005B :

NOT
CURRENTLY
USED

DS + 005C :
DS + 0080 :
DS + 0100 :

DEFAULT FCB
DEFAULT BUFFER
BEGIN USER DATA

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Each byte is indexed by 0, 1, and 2, corresponding to the standard Intel storage convention of low, middle, and high-order (most significant) byte. The value xxx in the preceding figure marks unused bytes. LC is the last code group location (24-bits, where the 4 high-order bits equal zero).

In the 8080 Model, the low order bytes of LC (LC0 and LC1) never exceed 0FFFFH and the high order byte (LC2) is always zero. BC is base paragraph address of the code group (16-bits). LD and BD provide the last position and paragraph base of the data group. The last position is one byte less than the group length. It should be noted that bytes LD0 and LD1 appear in the same relative positions of the base page in both CP/M-80 and CP/M-86, thus easing the program translation task. The M80 byte is equal to 1 when the 8080 Memory Model is in use. LE and BE provide the length and paragraph base of the optional extra group, while LS and BS give the optional stack group length and base. The bytes marked LX and BX correspond to a set of four optional independent groups which may be required for programs which execute using the Compact Memory Model. The initial values for these descriptors are derived from the header record in the memory image file, described in the following section.

TRANSIENT PROGRAM LOAD AND EXIT

In a manner similar to CP/M-80, the CCP parses up to two filenames following the command and places the properly formatted FCBs at locations 005CH and 006CH in the base page relative to the DS register. Under CP/M-80, the default DMA address is initialized to 0080H in the base page. Due to the segmented memory of the 8086 and 8088 processors, the DMA address is divided into two parts: the DMA segment address and the DMA offset. Therefore, under CP/M-86, the default DMA base is initialized to the value of DS, and the default DMA offset is initialized to 0080H. Thus, CP/M-80 and CP/M-86 operate in the same way: both assume the default DMA buffer occupies the second half of the base page.

The CCP transfers control to the transient program through an 8086 Far Call. The transient program may choose to use the 96-byte CCP stack and optionally return directly to the CCP upon program termination by executing a Far Return. Program termination also occurs when BDOS function zero is executed. Note that function zero can terminate a program without removing the program from memory or changing the memory allocation state (see Simple BDOS Calls in Chapter 4). The operator may terminate program execution by typing a single **CTRL-C** during line edited input, which has the same effect as the program executing BDOS function zero. Unlike the operation of CP/M-80, no disk reset occurs, and the CCP and BDOS modules are not reloaded from disk upon program termination.

Command (CMD) File Generation

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INTRODUCTION

As mentioned previously, two utility programs are provided with CP/M-86, called GENCMD and LMCMD, which are used to produce CMD memory image files suitable for execution under CP/M-86. GENCMD accepts Intel 8086 hex format files as input, while LMCMD reads Intel L-module files output from the standard Intel LOC86 Object Code Locator utility. GENCMD is used to process output from the Digital Research ASM-86 assembler and Intel's OH86 utility, while LMCMD is used when Intel-compatible developmental software is available for generation of programs targeted for CP/M-86 operation.

INTEL 8086 HEX FILE FORMAT

GENCMD input is in Intel hex format, which is produced by both the Digital Research ASM-86 assembler and the standard Intel OH86 utility program (see Intel document #9800639-03 entitled *MCS[®]-86 Software Development Utilities Operating Instructions for ISIS-II Users*). The CMD file produced by GENCMD contains a header record which defines the memory model and memory size requirements for loading and executing the CMD file.

An Intel hex file consists of the traditional sequence of ASCII records in the following format:

:	L	L	A	A	A	A	T	T	D	D	D	.	.	.	D	C	C
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

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where the beginning of the record is marked by an ASCII colon, and each subsequent digit position contains an ASCII hexadecimal digit in the range 0 - 9 or A - F. The fields are defined in the following table:

Intel Hex Field Definitions

Field	Contents
11	Record Length 00 - FF (0 - 255 in decimal)
aaaa	Load Address
tt	Record Type:
00	data record, loaded starting at offset aaaa from current base paragraph
01	end of file, cc = FF
02	extended address, aaaa is paragraph base for subsequent data records
03	start address is aaaa (ignored, IP set according to memory model in use)

The following are output from ASM-86 only:

81	same as 00, data belongs to code segment
82	same as 00, data belongs to data segment
83	same as 00, data belongs to stack segment
84	same as 00, data belongs to extra segment
85	paragraph address for absolute code segment
86	paragraph address for absolute data segment
87	paragraph address for absolute stack segment
88	paragraph address for absolute extra segment
d	Data Byte
cc	Check Sum (00 - Sum of Previous Digits)

OPERATION OF GENCMD

The GENCMD utility is invoked at the CCP level by using the following syntax:

```
GENCMD filename parameter-list
```

where the filename corresponds to the hex input file with an assumed (and unspecified) file type of H86. GENCMD accepts optional parameters to specifically identify the 8080 Memory Model and to describe memory requirements of each segment group. The GENCMD parameters are listed following the filename, as shown in the command line above where the parameter-list consists of a sequence of keywords and values separated by commas or blanks. The keywords are:

```
8080 CODE DATA EXTRA STACK X1 X2 X3 X4
```

The 8080 keyword forces a single code group so that the BDOS load function sets up the 8080 Memory Model for execution, thus allowing intermixed code and data within a single segment. The form of this command is:

```
GENCMD filename 8080
```

The remaining keywords follow the filename or the 8080 option and define specific memory requirements for each segment group, corresponding one-to-one with the segment groups defined in the previous section. In each case, the values corresponding to each group are enclosed in square brackets and separated by commas. Each value is a hexadecimal number

representing a paragraph address or segment length in paragraph units denoted by hhhh, prefixed by a single letter which defines the meaning of each value, as follows:

Value	Meaning
Ahhhh	Load the group at absolute location hhhh
Bhhhh	The group starts at hhhh in the hex file
Mhhhh	The group requires a minimum of $hhhh * 16$ bytes
Xhhhh	The group can address a maximum of $hhhh * 16$ bytes

Generally, the CMD file header values are derived directly from the hex file, and the parameters shown above need not be included. The following situations, however, require the use of GENCMD parameters.

- The 8080 keyword is included whenever ASM-86 is used in the conversion of 8080 programs to the 8086/8088 environment when code and data are intermixed within a single 64K segment, regardless of the use of CSEG and DSEG directives in the source program.
- An absolute address (A value) must be given for any group which must be located at an absolute location. Normally, this value is not specified, since CP/M-86 cannot generally ensure that the required memory region is available, in which case the CMD file cannot be loaded.
- The B value is used when GENCMD processes a hex file produced by Intel's OH86, or similar utility program that contains more than one group. The output from OH86 consists of a sequence of data records with no information to identify code, data, extra, stack, or auxiliary groups. In this case, the B value marks the beginning address of the group named by the keyword, causing GENCMD to load

data following this address to the named group (see the examples below). Thus, the B value is normally used to mark the boundary between code and data segments when no segment information is included in the hex file. Files produced by ASM-86 do not require the use of the B value since segment information is included in the hex file.

- The minimum memory value (M value) is included only when the hex records do not define the minimum memory requirements for the named group. Generally, the code group size is determined precisely by the data records loaded into the area. That is, the total space required for the group is defined by the range between the lowest and highest data byte addresses. The data group, however, may contain uninitialized storage at the end of the group and thus no data records are present in the hex file which define the highest referenced data item. The highest address in the data group can be defined within the source program by including a DB 0 as the last data item. Alternatively, the M value can be included to allocate the additional space at the end of the group. Similarly, the stack, extra, and auxiliary group sizes must be defined using the M value, unless the highest addresses within the groups are implicitly defined by data records in the hex file.
- The maximum memory size, given by the X value, is generally used when additional free memory may be needed for such purposes as I/O buffers or symbol tables. If the data area size is fixed, then the X parameter need not be included. In this case, the X value is assumed to be the same as the M value. The value XFFFF allocates the largest memory region available but, if used, the transient program must be aware that a three-byte length field is produced in the base page for this group where the high order byte may be non-zero. Programs converted directly from CP/M-80 or programs that use a 2-byte pointer to address buffers should restrict this value to XFFF or less, producing a maximum allocation length of 0FFF0H bytes.

The following GENCMD command line transforms the file X.H86 into the file X.CMD with the proper header record:

```
gencmd x code [a40] data[m30,xfff]
```

In this case, the code group is forced to paragraph address 40H, or equivalently, byte address 400H. The data group requires a minimum of 300H bytes, but can use up to 0FFF0H bytes, if available.

Assuming that a file Y.H86 exists which contains Intel hex records with no interspersed segment information, and also assuming that this file is on drive B, the command:

```
gencmd b:y data[b30,m20] extra [b50] stack [m40] xl[m40]
```

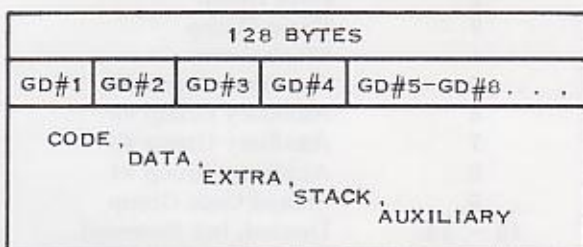
produces the file Y.CMD on drive B by selecting records beginning at address 0000H for the code segment, with records starting at 300H allocated to the data segment. The extra segment is filled from records beginning at 500H, while the stack and auxiliary segment #1 are uninitialized areas requiring a minimum of 400H bytes each. In this example, the data area requires a minimum of 200H bytes. Note again, that the B value need not be included if the Digital Research ASM-86 assembler is used.

OPERATION OF LMCMD

The LMCMD utility operates in exactly the same manner as GENCMD, with the exception that LMCMD accepts an Intel L-module file as input. The primary advantage of the L-module format is that the file contains internally coded information which defines values which would otherwise be required as parameters to GENCMD, such the beginning address of the group's data segment. Currently, however, the only language processors which use this format are the standard Intel development packages, although various independent vendors will be likely to take advantage of this format in the future.

COMMAND (CMD) FILE FORMAT

The CMD file produced by GENCMD and LMCMD consists of the 128-byte header record followed immediately by the memory image. Under normal circumstances, the format of the header record is of no consequence to a programmer. For completeness, however, the various fields of this record are shown in the following figure:



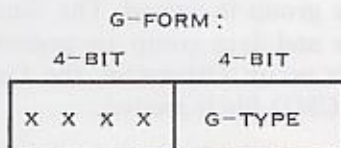
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In the previous figure, GD#2 through GD#8 represent Group Descriptors. Each Group Descriptor corresponds to an independently-loaded program unit and has the following fields:



2284777

where G-Form describes the group format, or has the value zero if no more descriptors follow. If G-Form is non-zero, then the 8-bit value is parsed as two fields:



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The G-Type field determines the Group Descriptor type. The valid Group Descriptors have a G-Type in the range of 1 through 9, as shown in the following table:

Group Descriptors	
G-Type	Group Type
1	Code Group
2	Data Group
3	Extra Group
4	Stack Group
5	Auxiliary Group #1
6	Auxiliary Group #2
7	Auxiliary Group #3
8	Auxiliary Group #4
9	Shared Code Group
10 - 14	Unused, but Reserved
15	Escape Code for Additional Types

All remaining values in the group descriptor are given in increments of 16-byte paragraph units with an assumed low-order 0 nibble to complete the 20-bit address. G-Length gives the number of paragraphs in the group. Given a G-length of 0080H, for example, the size of the group is 00800H = 2048D bytes. A-Base defines the base paragraph address for a non-relocatable group while G-Min and G-Max define the minimum and maximum size of the memory area to allocate to the group. G-Type 9 marks a pure code group for use under MP/M-86 and future versions of CP/M-86. Presently a Shared Code Group is treated as a non-shared Program Code Group under CP/M-86.

The memory model described by a header record is implicitly determined by the Group Descriptors. The 8080 Memory Model is assumed when only a code group is present, since no independent data group is named. The Small Model is implied when both a code and data group are present, but no additional group descriptors occur. Otherwise, the Compact Model is assumed when the CMD file is loaded.

Basic Disk Operating System (BDOS) Functions

Introduction	4-3
BDOS Parameters and Function Codes	4-3
Simple BDOS Calls	4-6
BDOS File Operations	4-11
BDOS Memory Management and Load	4-31

INTRODUCTION

This section presents the interface conventions which allow transient program access to CP/M-86 BDOS and BIOS functions. The BDOS calls correspond closely to CP/M-80 Version 2 in order to simplify translation of existing CP/M-80 programs for operation under CP/M-86. BDOS entry and exit conditions are described first, followed by a presentation of the individual BDOS function calls.

BDOS PARAMETERS AND FUNCTION CODES

Entry to the BDOS is accomplished through the 8086 software interrupt #224, which is reserved by Intel Corporation for use by CP/M-86 and MP/M-86. The function code is passed in register CL with byte parameters in DL and word parameters in DX. Single byte values are returned in AL, word values in both AX and BX, and double word values in ES and BX. All segment registers, except ES, are saved upon entry and restored upon exit from the BDOS (corresponding to PL/M-86 conventions). The following summarizes input and output parameter passing:

BDOS Parameter Summary

BDOS Entry Registers

CL Function Code
DL Byte Parameter
DX Word Parameter
DS Data Segment

BDOS Return Registers

Byte value returned in AL
Word value returned in both
AX and BX
Double-word value returned with offset
in BX and segment in ES

Note that the CP/M-80 BDOS requires an information address as input to various functions. This address usually provides buffer or File Control Block information used in the system call. In CP/M-86, however, the information address is derived from the current DS register combined with the offset given in the DX register. That is, the DX register in CP/M-86 performs the same function as the DE pair in CP/M-80, with the assumption that DS is properly set. This poses no particular problem for programs which use only a single data segment (as is the case for programs converted from CP/M-80), but when the data group exceeds a single segment, you must ensure that the DS register is set to the segment containing the data area related to the call. It should also be noted that zero values are returned for function calls which are out-of-range.

A list of CP/M-86 calls is given in the following table, with an asterisk following functions which differ from or are added to the set of CP/M-80 Version 2 functions.

CP/M-86 BDOS Functions

F#	Result
0*	System Reset
1	Console Output
2	Console Output
3	Reader Input
4	Punch Output
5	List Output
6*	Direct Console I/O
7	Get I/O Byte
8	Set I/O Byte
9	Print String
10	Read Console Buffer
11	Get Console Status
12	Return Version Number
13	Reset Disk System
14	Select Disk
15	Open File

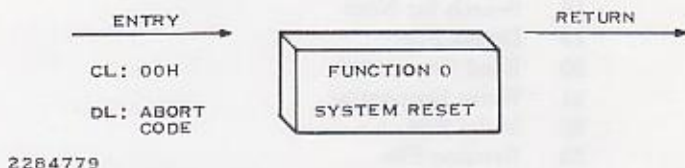
CP/M-86 BDOS Functions (Continued)

F#	Result
16	Close File
17	Search for First
18	Search for Next
19	Delete File
20	Read Sequential
21	Write Sequential
22	Make File
23	Rename File
24	Return Login Vector
26	Set DMA Address
26	Set DMA Address
27*	Get Addr (Alloc)
28	Write Protect Disk
29	Get Addr (R/O Vector)
30	Set File Attributes
31*	Get Addr (Disk Parms)
32	Set/Get User Code
33	Read Random
34	Write Random
35	Compute File Size
36	Set Random Record
37*	Reset Drive
40	Write Random with Zero Fill
50*	Direct BIOS Call
51*	Set DMA Segment Base
52*	Get DMA Segment Base
53*	Get Max Memory Available
54*	Get Max Memory at Abs Location
55*	Get Memory Region
56*	Get Absolute Memory Region
57*	Free memory region
58*	Free all memory
59*	Program load

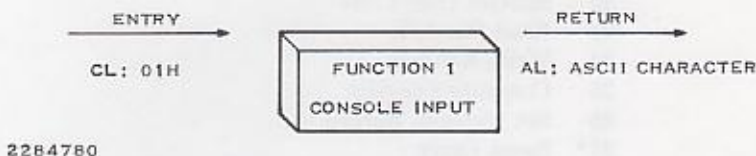
The individual BDOS functions are described in the following three sections, which cover the simple functions, file operations, and extended operations for memory management and program loading.

SIMPLE BDOS CALLS

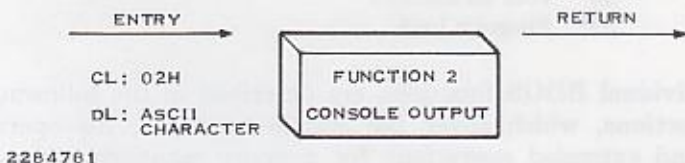
The first set of BDOS functions covers the range 0 through 12, and perform simple functions such as system reset and single character I/O.



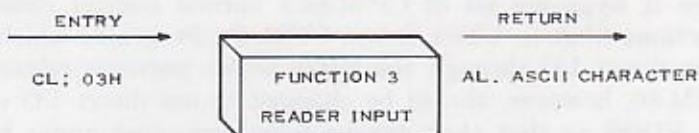
The system reset function returns control to the CP/M operating system at the CCP command level. The abort code in DL has two possible values: if DL = 00H, then the currently active program is terminated, and control is returned to the CCP. If DL is a 01H, the program remains in memory, and the memory allocation state remains unchanged.



The console input function reads the next character from the logical console device (CONSOLE) to register AL. Graphic characters, along with carriage return, line feed, and backspace (CTRL-H) are echoed to the console. Tab characters (CTRL-I) are expanded in columns of eight characters. The BDOS does not return to the calling program until a character has been typed, thus suspending execution if a character is not ready.

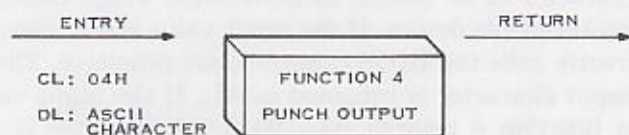


The ASCII character from DL is sent to the logical console. Tab characters expand in columns of eight characters. In addition, a check is made for start/stop scroll (**CTRL-S**).



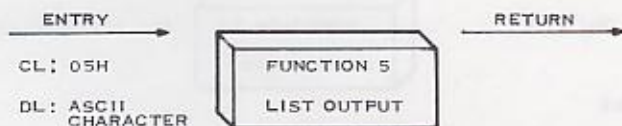
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The Reader Input function reads the next character from the logical reader (READER) into register AL. Control does not return until the character has been read.



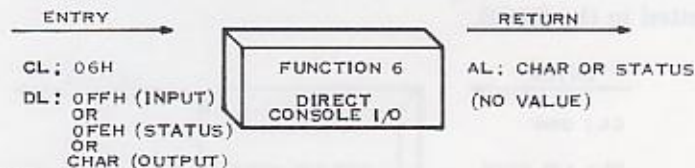
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The Punch Output function sends the character from register DL to the logical punch device (PUNCH).



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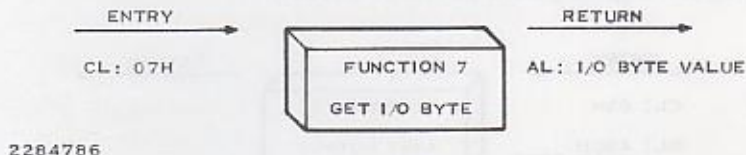
The List Output function sends the ASCII character in register DL to the logical list device (LIST).



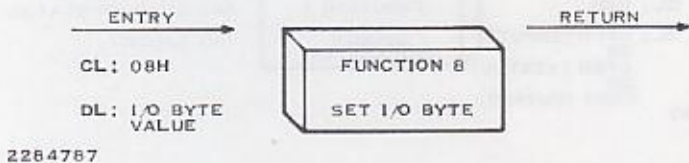
2284785

Direct console I/O is supported under CP/M-86 for those specialized applications where unadorned console input and output is required. Use of this function should, in general, be avoided since it bypasses all of CP/M-86's normal control character functions (that is, **CTRL-S** and **CTRL-P**). Programs which perform direct I/O through the BIOS under previous releases of CP/M-80, however, should be changed to use direct I/O under the BDOS so that they can be fully supported under future releases of MP/M™ and CP/M.

Upon entry to function 6, register DL either contains (1) a hexadecimal FF, denoting a CONSOLE input request, or (2) a hexadecimal FE, denoting a CONSOLE status request, or (3) an ASCII character to be output to CONSOLE where CONSOLE is the logical console device. If the input value is FF, then function 6 directly calls the BIOS console input primitive. The next console input character is returned in AL. If the input value is FE, then function 6 returns AL=00 if no character is ready and AL=FF otherwise. If the input value in DL is not FE or FF, then function 6 assumes that DL contains a valid ASCII character which is sent to the console.

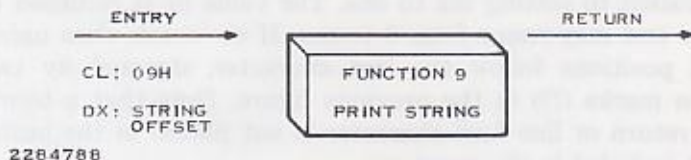


The Get I/O Byte function returns the current value of IOBYTE in register AL. The IOBYTE contains the current assignments for the logical devices CONSOLE, READER, PUNCH, and LIST provided the IOBYTE facility is implemented in the BIOS.

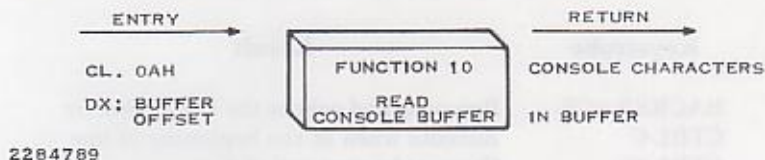


MP/M is a trademark of Digital Research.

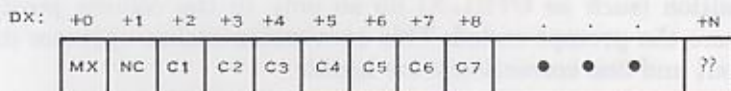
The Set I/O Byte function changes the system IOBYTE value to that given in register DL. This function allows transient program access to the IOBYTE in order to modify the current assignments for the logical devices CONSOLE, READER, PUNCH, and LIST.



The Print String function sends the character string stored in memory at the location given by DX to the logical console device (CONSOLE), until a dollar sign (\$) is encountered in the string. Tabs are expanded as in function 2, and checks are made for start/stop scroll and printer echo.



The Read Buffer function reads a line of edited console input into a buffer addressed by register DX from the logical console device (CONSOLE). Console input is terminated when either the input buffer is filled or when a return (CTRL-M) or a line feed (CTRL-J) character is entered. The input buffer addressed by DX takes the form:



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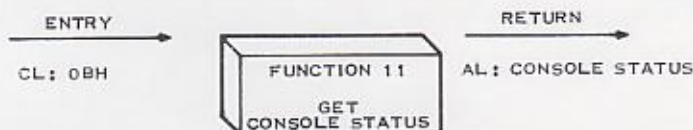
where mx is the maximum number of characters which the buffer will hold, and nc is the number of characters placed in the buffer. The characters entered by the operator follow the nc value. The value mx must be set prior to making a function 10 call and may range in value from 1 to 255. Setting mx to zero is equivalent to setting mx to one. The value nc is returned to the user and may range from 0 to mx . If $nc < mx$, then uninitialized positions follow the last character, denoted by two question marks (??) in the previous figure. Note that a terminating return or line feed character is not placed in the buffer and not included in the count nc .

A number of editing control functions are supported during console input under function 10. These are summarized in the following table:

Line Editing Controls

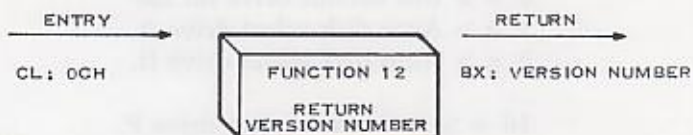
Keystroke	Result
BACKSPACE	Removes and echoes the last character
CTRL-C	Reboots when at the beginning of line
CTRL-E	Causes physical end of line
CTRL-H	Backspaces one character position
CTRL-J	Terminates input line (line feed)
CTRL-M	Terminates input line (return)
CTRL-R	Retypes the current line after new line
CTRL-U	Removes current line after new line
CTRL-X	Backspaces to beginning of current line

Certain functions which return the carriage to the leftmost position (such as **CTRL-X**) do so only to the column position where the prompt ended. This convention makes operator data input and line correction more legible.



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The Console Status function checks to see if a character has been typed at the logical console device (CONSOLE). If a character is ready, the value 01H is returned in register AL. Otherwise a 00H value is returned.



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Function 12 provides information which allows version independent programming. A two-byte value is returned, with BH = 00 designating the CP/M release (BH = 01 for MP/M), and BL = 00 for all releases previous to 2.0. CP/M 2.0 returns a hexadecimal 20 in register BL, with subsequent version 2 releases in the hexadecimal range of 21 through 2F. To provide version number compatibility, the initial release of CP/M-86 returns a 2.2.

BDOS File Operations

Functions 12 through 52 are related to disk file operations under CP/M-86. In many of these operations, DX provides the DS- relative offset to a file control block (FCB). The File Control Block (FCB) data area consists of a sequence of 33 bytes for sequential access, or a sequence of 36 bytes in the case that the file is accessed randomly. The default file control block normally located at offset 005CH from the DS register can be used for random access files, since bytes 007DH, 007EH, and 007FH are available for this purpose. Here is the FCB format, followed by definitions of each of its fields:

DR	F1	F2	//	F8	T1	T2	T3	EX	S1	S2	RC	D0	//	DN	CR	RD	R1	R2
00	01	02	...	08	09	10	11	12	13	14	15	16	...	31	32	33	34	35

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where:

dr	Drive code (0-16) 0 => Use default drive for file 1 => Auto disk select drive A, 2 => Auto disk select drive B, ... 16 => Auto disk select drive P.
f1...f8	Contain the file name in ASCII uppercase, with high bit = 0
t1,t2,t3	Contain the file type in ASCII uppercase, with high bit = 0 t1', t2', and t3' denote the high bit of these positions, t1' = 1 => Read/Only file, t2' = 1 => SYS file, no DIR list
ex	Contains the current extent number, nor- mally set to 00 by the user, but in range 0 - 31 during file I/O
s1	Reserved for internal system use
s2	Reserved for internal system use, set to zero on call to OPEN, MAKE, SEARCH
rc	Record count for extent x, takes on values from 0 - 128
d0...dn	Filled in by CP/M, reserved for system use
cr	Current record to read or write in a sequen- tial file operation, normally set to zero by user

r0,r1,r2 Optional random record number in the range 0-65535, with overflow to r2, r0, r1 constituting a 16-bit value with low byte r0, and high byte r1

For users of earlier versions of CP/M, it should be noted in passing that both CP/M Version 2 and CP/M-86 perform directory operations in a reserved area of memory that does not affect write buffer content, except in the case of Search and Search Next where the directory record is copied to the current DMA address.

There are three error situations that the BDOS may encounter during file processing, initiated as a result of a BDOS File I/O function call. When one of these conditions is detected, the BDOS issues the following message to the console:

BDOS ERR ON x: error

where x is the drive name of the drive selected when the error condition is detected, and the word error is one of the following three messages:

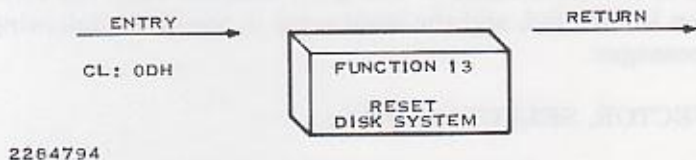
BAD SECTOR, SELECT, or R/O

These error situations are trapped by the BDOS, and thus the executing transient program is temporarily halted when the error is detected. No indication of the error situation is returned to the transient program.

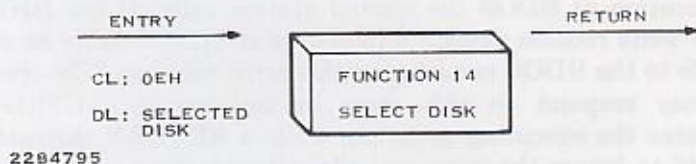
The BAD SECTOR error is issued as the result of an error condition returned to the BDOS from the BIOS module. The BDOS makes BIOS sector read and write commands as part of the execution of BDOS file related system calls. If the BIOS read or write routine detects a hardware error, it returns an error code to the BDOS resulting in this error message. The operator may respond to this error in two ways: a **CTRL-C** terminates the executing program, while a **RETURN** instructs CP/M-86 to ignore the error and allow the program to continue execution.

The SELECT error is also issued as the result of an error condition returned to the BDOS from the BIOS module. The BDOS makes a BIOS disk select call prior to issuing any BIOS read or write to a particular drive. If the selected drive is not supported in the BIOS module, it returns an error code to the BDOS resulting in this error message. CP/M-86 terminates the currently running program and returns to the command level of the CCP following any input from the console.

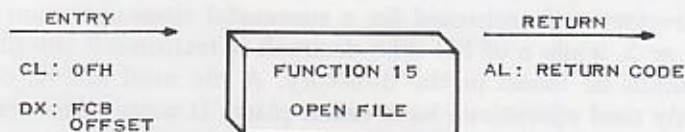
The R/O message occurs when the BDOS receives a command to write to a drive that is in read-only status. Drives may be placed in read-only status explicitly as the result of a STAT command or BDOS function call, or implicitly if the BDOS detects that disk media has been changed without performing a warm start. The ability to detect changed media is optionally included in the BIOS, and exists only if a checksum vector is included for the selected drive. Upon entry of any character at the keyboard, the transient program is aborted, and control returns to the CCP.



The Reset Disk Function is used to programmatically restore the file system to a reset state where all disks are set to read/write (see functions 28 and 29), only disk drive A is selected. This function can be used, for example, by an application program which requires disk changes during operation. Function 37 (Reset Drive) can also be used for this purpose.



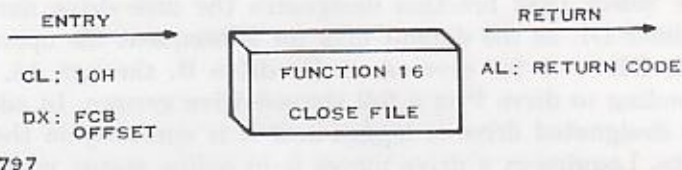
The Select Disk function designates the disk drive named in register DL as the default disk for subsequent file operations, with DL = 0 for drive A, 1 for drive B, through 15, corresponding to drive P in a full sixteen-drive system. In addition, the designated drive is logged-in if it is currently in the reset state. Logging-in a drive places it in online status which activates the drive's directory until the next cold start, warm start, disk system reset, or drive reset operation. FCBs which specify drive code zero (dr = 00H) automatically reference the currently selected default drive. Drive code values between 1 and 16, however, ignore the selected default drive and directly reference drives A through P.



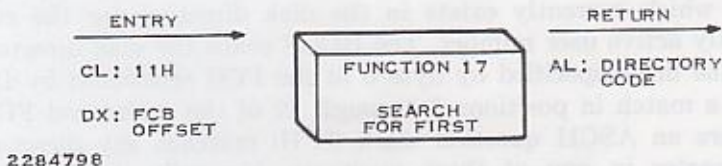
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The Open File operation is used to activate a FCB specifying a file which currently exists in the disk directory for the currently active user number. The BDOS scans the disk directory of the drive specified by byte 0 of the FCB referenced by DX for a match in positions 1 through 12 of the referenced FCB, where an ASCII question mark (3FH) matches any directory character in any of these positions. Normally, no question marks are included and, further, byte x of the FCB is set to zero before making the open call.

If a directory element is matched, the relevant directory information is copied into bytes d0 through dn of the FCB, thus allowing access to the files through subsequent read and write operations. Note that an existing file must not be accessed until a successful open operation is completed. Further, an FCB not activated by either an open or make function must not be used in BDOS read or write commands. Upon return, the open function returns a directory code with the value 0 through 3 if the open was successful, or 0FFH (255 decimal) if the file cannot be found. If question marks occur in the FCB then the first matching FCB is activated. Note that the current record (cr) must be zeroed by the program if the file is to be accessed sequentially from the first record.



The Close File function performs the inverse of the open file function. Given that the FCB addressed by DX has been previously activated through an open or make function (see functions 15 and 22), the close function permanently records the new FCB in the referenced disk directory. The FCB matching process for the close is identical to the open function. The directory code returned for a successful close operation is 0, 1, 2, or 3, while a 0FFH (255 decimal) is returned if the file name cannot be found in the directory. A file need not be closed if only read operations have taken place. If write operations have occurred, however, the close operation is necessary to permanently record the new directory information.



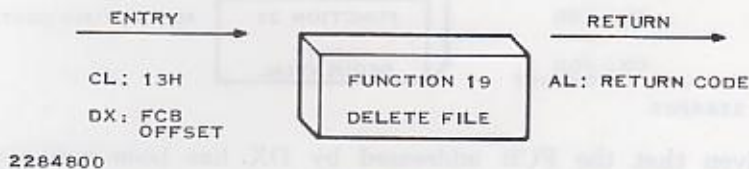
Search First scans the directory for a match with the file given by the FCB addressed by DX. The value 255 (hexadecimal FF) is returned if the file is not found, otherwise 0, 1, 2, or 3 is returned indicating the file is present. In the case that the file is found, the buffer at the current DMA address is filled with the record containing the directory entry, and its relative starting position is $AL * 32$ (that is, rotate the AL register left 5 bits). Although not normally required for application programs, the directory information can be extracted from the buffer at this position.

An ASCII question mark (63 decimal, 3F hexadecimal) in any position from f1 through x matches the corresponding field of any directory entry on the default or auto-selected disk drive. If the dr field contains an ASCII question mark, then the auto

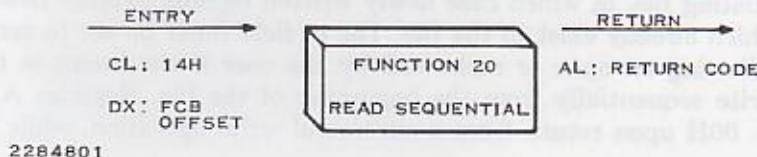
disk select function is disabled, the default disk is searched, with the search function returning any matched entry, allocated or free, belonging to any user number. This latter function is not normally used by application programs, but does allow complete flexibility to scan all current directory values. If the dr field is not a question mark, the s2 byte is automatically zeroed.



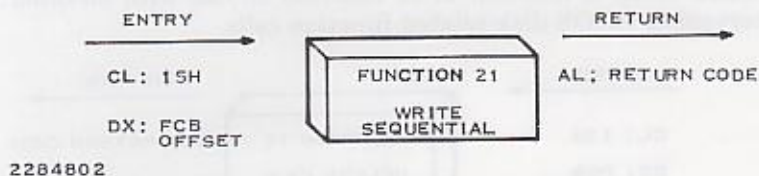
The Search Next function is similar to the Search First function, except that the directory scan continues from the last matched entry. In a way similar to function 17, function 18 returns the decimal value 255 in A when no more directory items match. In terms of execution sequence, a function 18 call must follow either a function 17 or function 18 call with no other intervening BDOS disk related function calls.



The Delete File function removes files which match the FCB addressed by DX. The filename and type may contain ambiguous references (that is, question marks in various positions), but the drive select code cannot be ambiguous, as in the Search and Search Next functions. Function 19 returns a 0FFH (decimal 255) if the referenced file or files cannot be found, otherwise a value of zero is returned.



Given that the FCB addressed by DX has been activated through an open or make function (numbers 15 and 22), the Read Sequential function reads the next 128 byte record from the file into memory at the current DMA address. The record is read from position cr of the extent, and the cr field is automatically incremented to the next record position. If the cr field overflows then the next logical extent is automatically opened and the cr field is reset to zero in preparation for the next read operation. The cr field must be set to zero following the open call by the user if the intent is to read sequentially from the beginning of the file. The value 00H is returned in the AL register if the read operation was successful, while a value of 01H is returned if no data exists at the next record position of the file. Normally, the no data situation is encountered at the end of a file. However, it can also occur if an attempt is made to read a data block which has not been previously written, or an extent which has not been created. These situations are usually restricted to files created or appended by use of the BDOS Write Random command (function 34).

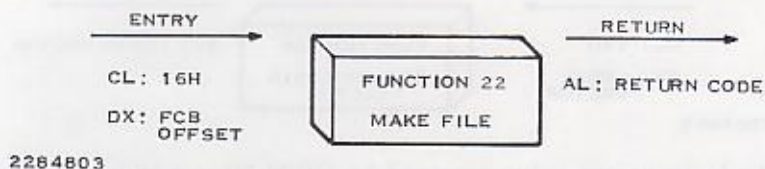


Given that the FCB addressed by DX has been activated through an open or make function (numbers 15 and 22), the Write Sequential function writes the 128 byte data record at the current DMA address to the file named by the FCB. The record is placed at position cr of the file, and the cr field is automatically incremented to the next record position. If the cr field overflows then the next logical extent is automatically opened and the cr field is reset to zero in preparation for the next write operation. Write operations can take place into an existing file, in which case newly written records overlay those which already exist in the file. The cr field must be set to zero following an open or make call by the user if the intent is to write sequentially from the beginning of the file. Register AL = 00H upon return from a successful write operation, while a

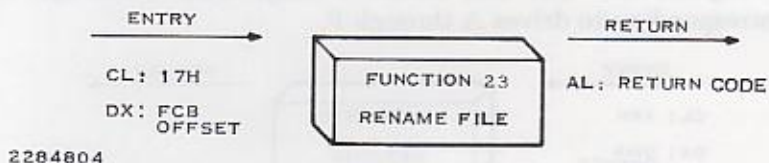
non-zero value indicates an unsuccessful write due to one of the following conditions:

01 No available directory space — This condition occurs when the write command attempts to create a new extent that requires a new directory entry and no available directory entries exist on the selected disk drive.

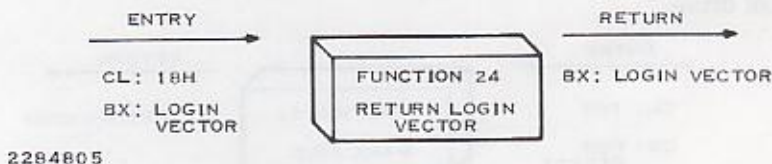
02 No available data block — This condition is encountered when the write command attempts to allocate a new data block to the file and no unallocated data blocks exist on the selected disk drive.



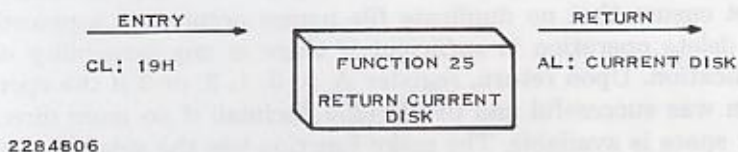
The Make File operation is similar to the open file operation except that the FCB must name a file which does not exist in the currently referenced disk directory (that is, the one named explicitly by a non-zero dr code, or the default disk if dr is zero). The BDOS creates the file and initializes both the directory and main memory value to an empty file. The programmer must ensure that no duplicate file names occur, and a preceding delete operation is sufficient if there is any possibility of duplication. Upon return, register A = 0, 1, 2, or 3 if the operation was successful and 0FFH (255 decimal) if no more directory space is available. The make function has the side-effect of activating the FCB and thus a subsequent open is not necessary.



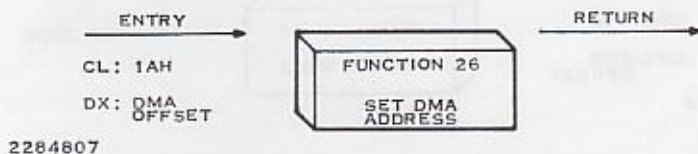
The Rename function uses the FCB addressed by DX to change all directory entries of the file specified by the file name in the first 16 bytes of the FCB to the file name in the second 16 bytes. It is the user's responsibility to insure that the file names specified are valid CP/M unambiguous file names. The drive code dr at position 0 is used to select the drive, while the drive code for the new file name at position 16 of the FCB is ignored. Upon return, register AL is set to a value of zero if the rename was successful, and 0FFH (255 decimal) if the first file name could not be found in the directory scan.



The login vector value returned by CP/M-86 is a 16-bit value in BX, where the least significant bit corresponds to the first drive A, and the high order bit corresponds to the sixteenth drive, labelled P. A 0 bit indicates that the drive is not online, while a 1 bit marks a drive that is actively online due to an explicit disk drive selection, or an implicit drive select caused by a file operation which specified a non-zero dr field.



Function 25 returns the currently selected default disk number in register AL. The disk numbers range from 0 through 15 corresponding to drives A through P.



DMA is an acronym for Direct Memory Address, which is often used in connection with disk controllers which directly access the memory of the mainframe computer to transfer data to and from the disk subsystem. Although many computer systems use non-DMA access (that is, the data is transferred through programmed I/O operations), the DMA address has, in CP/M, come to mean the address at which the 128 byte data record resides before a disk write and after a disk read. In the CP/M-86 environment, the Set DMA function is used to specify the offset of the read or write buffer from the current DMA base. Therefore, to specify the DMA address, both a function 26 call and a function 51 call are required. Thus, the DMA address becomes the value specified by DX plus the DMA base value until it is changed by a subsequent Set DMA or set DMA base function.

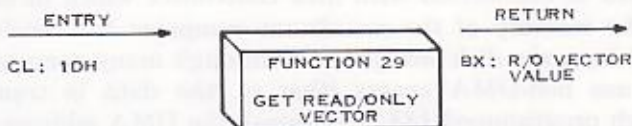


An allocation vector is maintained in main memory for each on-line disk drive. Various system programs use the information provided by the allocation vector to determine the amount of remaining storage (see the STAT program). Function 27 returns the segment base and the offset address of the allocation vector for the currently selected disk drive. The allocation information may, however, be invalid if the selected disk has been marked read/only.



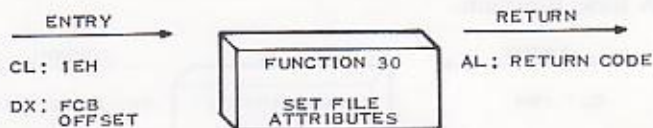
The disk write protect function provides temporary write protection for the currently selected disk. Any attempt to write to the disk, before the next cold start, warm start, disk system reset, or drive reset operation produces the following message:

Bdos Err on d: R/O



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Function 29 returns a bit vector in register BX which indicates drives which have the temporary read/only bit set. In a manner similar to function 24, the least significant bit corresponds to drive A, while the most significant bit corresponds to drive P. The R/O bit is set either by an explicit call to function 28, or by the automatic software mechanisms within CP/M-86 which detect changed disks.



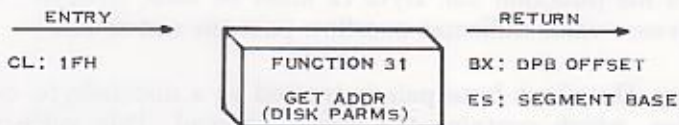
2284811

The Set File Attributes function allows programmatic manipulation of permanent indicators attached to files. In particular, the R/O, System and Archive attributes (t1', t2', and t3') can be set or reset. The DX pair addresses a FCB containing a file name with the appropriate attributes set or reset. It is the user's responsibility to insure that an ambiguous file name is not specified. Function 30 searches the default disk drive directory area for directory entries that belong to the current user number and that match the FCB specified name and type fields. All matching directory entries are updated to contain the selected indicators. Indicators f1' through f4' are not presently used, but may be useful for applications programs, since they are not involved in the matching process during file open and close operations. Indicators f5' through f8' are reserved for future system expansion. The currently assigned attributes are defined as follows:

- t1s': The R/O attribute indicates, if set, that the file is in read/only status. BDOS will not allow write commands to be issued to files in R/O status.

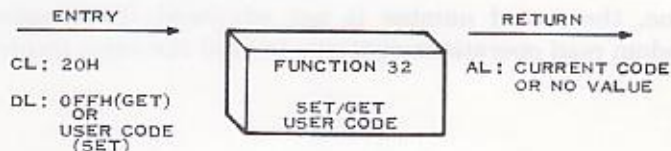
- t2\$: The System attribute is referenced by the CP/M DIR utility. If set, DIR will not display the file in a directory display.
- t3\$: The Archive attribute is reserved but not actually used by CP/M-86. If set it indicates that the file has been written to back up storage by a user written archive program. To implement this facility, the archive program sets this attribute when it copies a file to back up storage; any programs updating or creating files reset this attribute. Further, the archive program backs up only those files that have the Archive attribute reset. Thus, an automatic back up facility restricted to modified files can be easily implemented.

Function 30 returns with register AL set to 0fFH (255 decimal) if the referenced file cannot be found, otherwise a value of zero is returned.



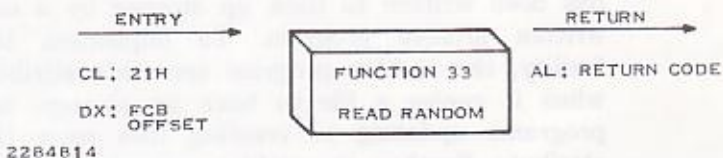
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The offset and the segment base of the BIOS resident disk parameter block of the currently selected drive are returned in BX and ES as a result of this function call. This control block can be used for either of two purposes. First, the disk parameter values can be extracted for display and space computation purposes, or transient programs can dynamically change the values of current disk parameters when the disk environment changes, if required. Normally, application programs will not require this facility. The paragraph entitled GENDEF Output, in Chapter 6, defines the BIOS disk parameter block.



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An application program can change or interrogate the currently active user number by calling function 32. If register DL = 0FFH, then the value of the current user number is returned in register AL, where the value is in the range 0 to 15. If register DL is not 0FFH, then the current user number is changed to the value of DL (modulo 16).



The Read Random function is similar to the sequential file read operation of previous releases, except that the read operation takes place at a particular record number, selected by the 24-bit value constructed from the three byte field following the FCB (byte positions r0 at 33, r1 at 34, and r2 at 35). Note that the sequence of 24 bits is stored with the least significant byte first (r0), the middle byte next (r1), and the high byte last (r2). CP/M does not reference byte r2, except in computing the size of a file (function 35). Byte r2 must be zero, however, since a non-zero value indicates overflow past the end of file.

Thus, the r0, r1 byte pair is treated as a double-byte, or word value, which contains the record to read. This value ranges from 0 to 65535, providing access to any particular record of any size file. In order to access a file using the Read Random function, the base extent (extent 0) must first be opened. Although the base extent may or may not contain any allocated data, this ensures that the FCB is properly initialized for subsequent random access operations. The selected record number is then stored in the random record field (r0,r1), and the BDOS is called to read the record. Upon return from the call, register AL either contains an error code, as listed below, or the value 00 indicating the operation was successful. In the latter case, the buffer at the current DMA address contains the randomly accessed record. Note that contrary to the sequential read operation, the record number is not advanced. Thus, subsequent random read operations continue to read the same record.

Upon each random read operation, the logical extent and current record values are automatically set. Thus, the file can be sequentially read or written, starting from the current randomly accessed position. Note, however, that in this case, the last randomly read record will be re-read as you switch from random mode to sequential read, and the last record will be re-written as you switch to a sequential write operation. You can, of course, simply advance the random record position following each random read or write to obtain the effect of a sequential I/O operation.

Error codes returned in register AL following a random read are listed in the following table:

Function 33 (Read Random) Error Codes

Code	Meaning
01	Reading unwritten data — This error code is returned when a random read operation accesses a data block which has not been previously written.
02	(Not returned by the Random Read command)
03	Cannot close current extent — This error code is returned when BDOS cannot close the current extent prior to moving to the new extent containing the record specified by bytes r0, r1 of the FCB. This error can be caused by an overwritten FCB or a read random operation on an FCB that has not been opened.
04	Seek to unwritten extent — This error code is returned when a random read operation accesses an extent that has not been created. This error situation is equivalent to error 01.
05	(Not returned by the Random Read command)
06	Random record number out of range — This error code is returned whenever byte r2 of the FCB is non-zero.

Normally, non-zero return codes can be treated as missing data, with zero return codes indicating operation complete.



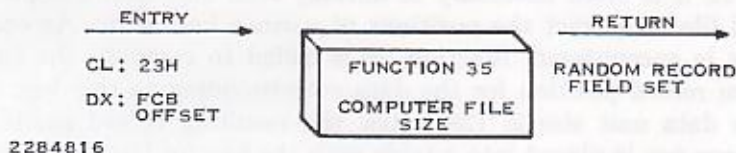
The Write Random operation is initiated in a manner similar to the Read Random call, except that data is written to the disk from the current DMA address. Further, if the disk extent or data block which is the target of the write has not yet been allocated, the allocation is performed before the write operation continues. As in the Read Random operation, the random record number is not changed as a result of the write. The logical extent number and current record positions of the file control block are set to correspond to the random record which is being written. Sequential read or write operations can commence following a random write, with the note that the currently addressed record is either read or rewritten again as the sequential operation begins. You can also simply advance the random record position following each write to get the effect of a sequential write operation. In particular, reading or writing the last record of an extent in random mode does not cause an automatic extent switch as it does in sequential mode.

In order to access a file using the Write Random function, the base extent (extent 0) must first be opened. As in the Read Random function, this ensures that the FCB is properly initialized for subsequent random access operations. If the file is empty, a Make File function must be issued for the base extent. Although the base extent may or may not contain any allocated data, this ensures that the file is properly recorded in the directory, and is visible in DIR requests.

Upon return from a Write Random call, register AL either contains an error code, as listed in the following table, or the value 00, indicating that the operation was successful.

Function 34 (WRITE RANDOM) Error Codes

Code	Meaning
01	(Not returned by the Random Write command)
02	No available data block — This condition is encountered when the Write Random command attempts to allocate a new data block to the file and no unallocated data blocks exist on the selected disk drive.
03	Cannot close current extent — This error code is returned when BDOS cannot close the current extent prior to moving to the new extent containing the record specified by bytes r0, r1 of the FCB. This error can be caused by an overwritten FCB or a write random operation on an FCB that has not been opened.
04	(Not returned by the Random Write command)
05	No available directory space — This condition occurs when the write command attempts to create a new extent that requires a new directory entry and no available directory entries exist on the selected disk drive.
06	Random record number out of range — This error code is returned whenever byte r2 of the FCB is non-zero.



When computing the size of a file, the DX register addresses an FCB in random mode format (bytes r0, r1, and r2 are present). The FCB contains an unambiguous file name which is used in the directory scan. Upon return, the random record bytes contain the virtual file size which is, in effect, the record address of the record following the end of the file. If, following a call to function 35, the high record byte r2 is 01, then the file contains the maximum record count 65,536. Otherwise, bytes r0 and r1 constitute a 16-bit value (r0 is the least significant byte, as before) which is the file size.

Data can be appended to the end of an existing file by simply calling function 35 to set the random record position to the end of file, then performing a sequence of random writes starting at the preset record address.

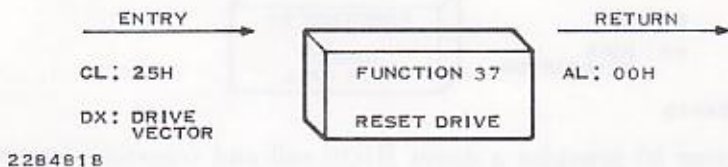
The virtual size of a file corresponds to the physical size when the file is written sequentially. If, instead, the file was created in random mode and holes exist in the allocation, then the file may in fact contain fewer records than the size indicates. If, for example, a single record with record number 65535 (CP/M's maximum record number) is written to a file using the Write Random function, then the virtual size of the file is 65536 records, although only one block of data is actually allocated.



The Set Random Record function causes the BDOS to automatically produce the random record position of the next record to be accessed from a file which has been read or written sequentially to a particular point. The function can be useful in two ways.

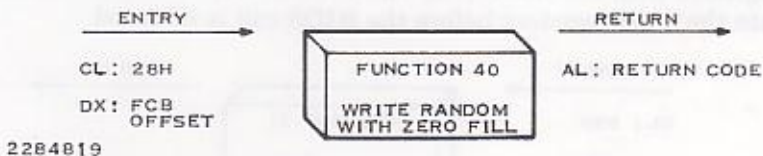
First, it is often necessary to initially read and scan a sequential file to extract the positions of various key fields. As each key is encountered, function 36 is called to compute the random record position for the data corresponding to this key. If the data unit size is 128 bytes, the resulting record position minus one is placed into a table with the key for later retrieval. After scanning the entire file and tabulating the keys and their record numbers, you can move instantly to a particular keyed record by performing a random read using the corresponding random record number which was saved earlier. The scheme is easily generalized when variable record lengths are involved, since the program need only store the buffer-relative byte position along with the key and record number in order to find the exact starting position of the keyed data at a later time.

A second use of function 36 occurs when switching from a sequential read or write to random read or write. A file is sequentially accessed to a particular point in the file, function 36 is called which sets the record number, and subsequent random read and write operations continue from the next record in the file.



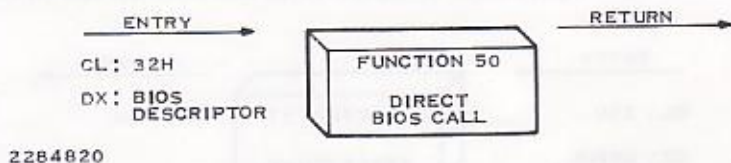
The Reset Drive function is used to programmatically restore specified drives to the reset state (a reset drive is not logged-in and is in read/write status). The passed parameter in register DX is a 16 bit vector of drives to be reset, where the least significant bit corresponds to the first drive, A, and the high order bit corresponds to the sixteenth drive, labelled P. Bit values of 1 indicate that the specified drive is to be reset.

In order to maintain compatibility with MP/M™, CP/M returns a zero value for this function.

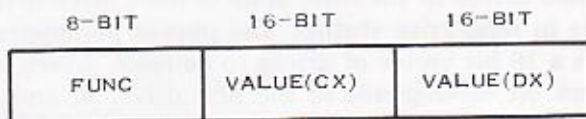


The Write Random With Zero Fill function is similar to the Write Random function (function 34) with the exception that a previously unallocated data block is initialized to records filled with zeros before the record is written. If this function has been used to create a file, records accessed by a read random

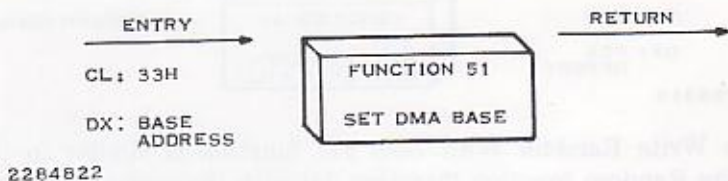
operation that contain all zeros identify unwritten random record numbers. Unwritten random records in allocated data blocks of files created using the Write Random function contain uninitialized data.



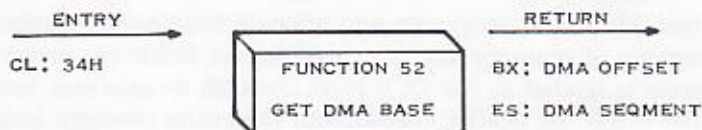
Function 50 provides a direct BIOS call and transfers control through the BDOS to the BIOS. The DX register addresses a five-byte memory area containing the BIOS call parameters:



where Func is a BIOS function number, (see the first table in Chapter 5) and value (CX) and value (DX) are the 16-bit values which would normally be passed directly in the CX and DX registers with the BIOS call. The CX and DX values are loaded into the 8086 registers before the BIOS call is initiated.



Function 51 sets the base register for subsequent DMA transfers. The word parameter in DX is a paragraph address and is used with the DMA offset to specify the address of a 128 byte buffer area to be used in the disk read and write functions. Note that upon initial program load, the default DMA base is set to the address of the user's data segment (the initial value of DS) and the DMA offset is set to 0080H, which provides access to the default buffer in the base page.



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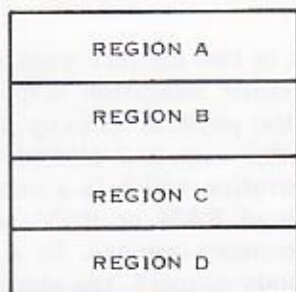
Function 52 returns the current DMA Base Segment address in ES, with the current DMA Offset in DX.

BDOS MEMORY MANAGEMENT AND LOAD

Memory is allocated in two distinct ways under CP/M-86. The first is through a static allocation map, located within the BIOS, that defines the physical memory which is available on the host system. In this way, it is possible to operate CP/M-86 in a memory configuration which is a mixture of up to eight non-contiguous areas of RAM or ROM, along with reserved, missing, or faulty memory regions. In a simple RAM-based system with contiguous memory, the static map defines a single region, usually starting at the end of the BIOS and extending to the end of available memory.

Once memory is physically mapped in this manner, CP/M-86 performs the second level of dynamic allocation to support transient program loading and execution. CP/M-86 again allows dynamic allocation of memory into eight regions. A request for allocation takes place either implicitly, through a program load operation, or explicitly through the BDOS calls given in this section. Programs themselves are loaded in two ways: through a command entered at the CCP level, or through the BDOS Program Load operation (function 59). Multiple programs can be loaded at the CCP level, as long as each program executes a System Reset (function 0) and remains in memory (DL = 01H). Multiple programs of this type only receive control by intercepting interrupts, and thus under normal circumstances there is only one transient program in memory at any given time. If, however, multiple programs are present in memory, then CTRL-C characters entered by the operator delete these programs in the opposite order in which they were loaded no matter which program is actively reading the console.

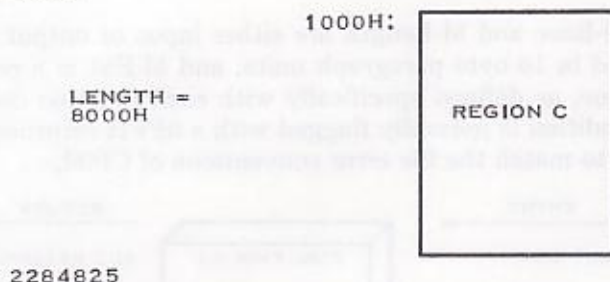
Any given program loaded through a CCP command can, itself, load additional programs and allocate data areas. Suppose four regions of memory are allocated in the following order: a program is loaded at the CCP level through an operator command. The CMD file header is read, and the entire memory image consisting of the program and its data is loaded into region A, and execution begins. This program, in turn, calls the BDOS Program Load function (59) to load another program into region B, and transfers control to the loaded program. The region B program then allocates an additional region C, followed by a region D. The order of allocation is shown in the following figure:



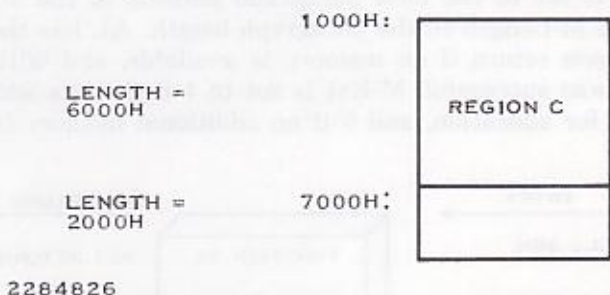
2284824

There is a hierarchical ownership of these regions: the program in A controls all memory from A through D. The program in B also controls regions B through D. The program in A can release regions B through D, if desired, and reload yet another program. DDT-86, for example, operates in this manner by executing the Free Memory call (function 57) to release the memory used by the current program before loading another test program. Further, the program in B can release regions C and D if required by the application. It must be noted, however, that if either A or B terminates by a System Reset (BDOS function 0 with DL = 00H) then all four regions A through D are released.

A transient program may release a portion of a region, allowing the released portion to be assigned on the next allocation request. The released portion must, however, be at the beginning or end of the region. Suppose, for example, the program in region B above receives 800H paragraphs at paragraph location 100H following its first allocation request as shown in the following figure:

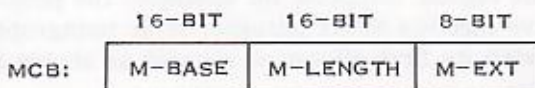


Suppose further that region D is then allocated. The last 200H paragraphs in region C can be returned without affecting region D by releasing the 200H paragraphs beginning at paragraph base 700H, resulting in the memory arrangement shown in the following figure:

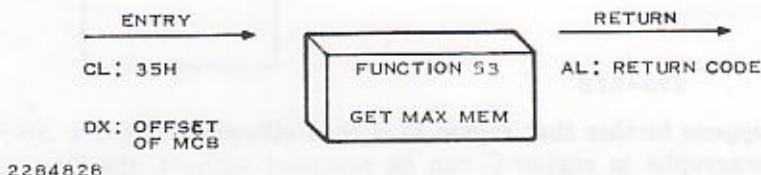


The region beginning at paragraph address 700H is now available for allocation in the next request. Note that a memory request will fail if eight memory regions have already been allocated. Normally, if all program units can reside in a contiguous region, the system allocates only one region.

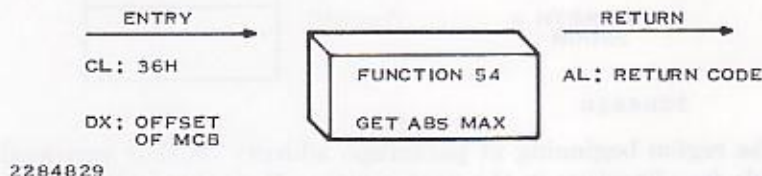
Memory management functions beginning at 53 reference a Memory Control Block (MCB), defined in the calling program, which takes the form:



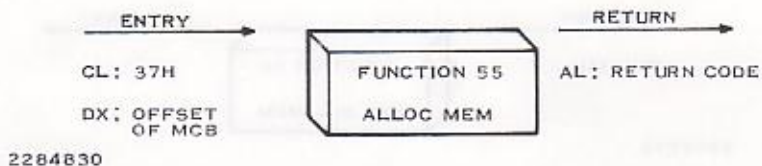
where M-Base and M-Length are either input or output values expressed in 16-byte paragraph units, and M-Ext is a returned byte value, as defined specifically with each function code. An error condition is normally flagged with a 0FFH returned value in order to match the file error conventions of CP/M.



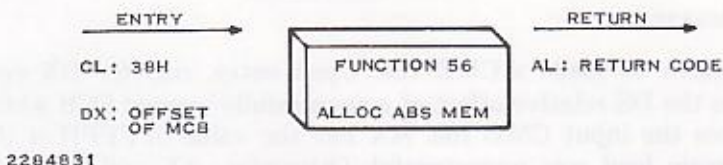
Function 53 finds the largest available memory region which is less than or equal to M-Length paragraphs. If successful, M-Base is set to the base paragraph address of the available area, and M-Length to the paragraph length. AL has the value 0FFH upon return if no memory is available, and 00H if the request was successful. M-Ext is set to 1 if there is additional memory for allocation, and 0 if no additional memory is available.



Function 54 is used to find the largest possible region at the absolute paragraph boundary given by M-Base, for a maximum of M- Length paragraphs. M-Length is set to the actual length if successful. AL has the value 0FFH upon return if no memory is available at the absolute address, and 00H if the request was successful.



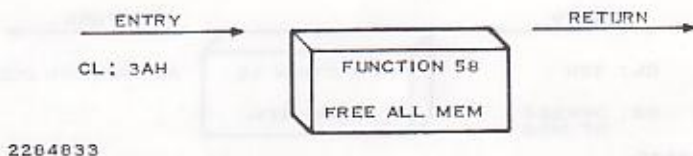
The allocate memory function allocates a memory area according to the MCB addressed by DX. The allocation request size is obtained from M-Length. Function 55 returns, in the user's MCB, the base paragraph address of the allocated region. Register AL contains a 00H if the request was successful, and a 0FFH if the memory could not be allocated.



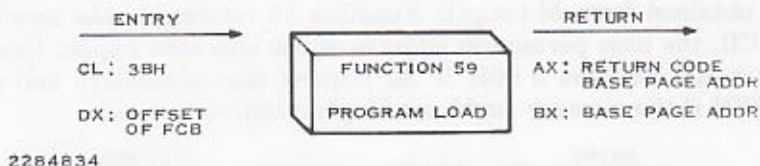
The allocate absolute memory function allocates a memory area according to the MCB addressed by DX. The allocation request size is obtained from M-Length and the absolute base address from M-Base. Register AL contains a 00H if the request was successful and a 0FFH if the memory could not be allocated.



Function 57 is used to release memory areas allocated to the program. The value of the M-Ext field controls the operation of this function: if M-Ext = 0FFH then all memory areas allocated by the calling program are released. Otherwise, the memory area of length M-Length at location M-Base given in the MCB addressed by DX is released (the M-Ext field should be set to 00H in this case). As described above, either an entire allocated region must be released, or the end of a region must be released: the middle section cannot be returned under CP/M-86.



Function 58 is used to release all memory in the CP/M-86 environment (normally used only by the CCP upon initialization).



Function 59 loads a CMD file. Upon entry, register DX contains the DS relative offset of a successfully opened FCB which names the input CMD file. AX has the value 0FFFFH if the program load was unsuccessful. Otherwise, AX and BX both contain the paragraph address of the base page belonging to the loaded program. The base address and segment length of each segment is stored in the base page. Note that upon program load at the CCP level, the DMA base address is initialized to the base page of the loaded program, and the DMA offset address is initialized to 0080H. However, this is a function of the CCP, and a function 59 does not establish a default DMA address. It is the responsibility of the program which executes function 59 to execute function 51 to set the DMA base and to function 26 to set the DMA offset before passing control to the loaded program.

Basic I/O System (BIOS) Organization

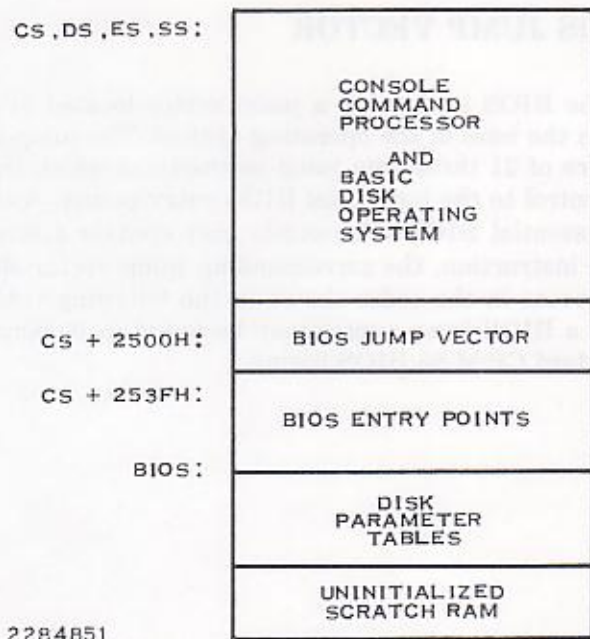
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INTRODUCTION

The distribution version of CP/M-86 is set up for operation with the Intel SBC 86/12 microcomputer and an Intel 204 diskette controller. All hardware dependencies are, however, concentrated in subroutines which are collectively referred to as the Basic I/O System, or BIOS. A CP/M-86 system implementor can modify these subroutines, as described below, to tailor CP/M-86 to fit nearly any 8086 or 8088 operating environment. This section describes the actions of each BIOS entry point, and defines variables and tables referenced within the BIOS. The discussion of Disk Definition Tables is, however, treated separately in the next section of this manual.

ORGANIZATION OF THE BIOS

The BIOS portion of CP/M-86 resides in the topmost portion of the operating system (highest addresses), and takes the general form shown in the following figure:



As described in the following sections, the CCP and BDOS are supplied with CP/M-86 in hex file form as CPM.H86. In order to implement CP/M-86 on non-standard hardware, you must create a BIOS which performs the functions listed below and concatenate the resulting hex file to the end of the CPM.H86 file. The GENCMD utility is then used to produce the CPM.SYS file for subsequent load by the cold start loader. The cold start loader that loads the CPM.SYS file into memory contains a simplified form of the BIOS, called the LDBIOS (Loader BIOS). It loads CPM.SYS into memory at the location defined in the CPM.SYS header (usually 0400H). The procedure to follow in construction and execution of the cold start loader and the CP/M-86 Loader is given in a later section.

Appendix D contains a listing of the standard CP/M-86 BIOS for the Intel SBC 86/12 system using the Intel 204 Controller Board. Appendix E shows a sample skeletal BIOS, called CBIOS, that contains the essential elements with the device drivers removed. You may wish to review these listings in order to determine the overall structure of the BIOS.

THE BIOS JUMP VECTOR

Entry to the BIOS is through a jump vector located at offset 2500H from the base of the operating system. The jump vector is a sequence of 21 three-byte jump instructions which transfer program control to the individual BIOS entry points. Although some non-essential BIOS subroutines may contain a single return (RET) instruction, the corresponding jump vector element must be present in the order shown in the following table. An example of a BIOS jump vector may be found in Appendix D, in the standard CP/M-86 BIOS listing.

Parameters for the individual subroutines in the BIOS are passed in the CX and DX registers, when required. CX receives the first parameter; DX is used for a second argument. Return values are passed in the registers according to type: byte values are returned in AL. Word values (16 bits) are returned in BX. Specific parameters and returned values are described with each subroutine.

BIOS Jump Vector

Offset from Beginning of Bios	Suggested BIOS Instruction	BIOS F#	Description
2500H	JMP INIT	0	Arrive Here from Cold Boot
2503H	JMP WBOOT	1	Arrive Here for Warm Start
2506H	JMP CONST	2	Check for Console Char Ready
2509H	JMP CONIN	3	Read Console Character In
250CH	JMP CONOUT	4	Write Console Character Out
250FH	JMP LIST	5	Write Listing Character Out
2512H	JMP PUNCH	6	Write Char to Punch Device
2515H	JMP READER	7	Read Reader Device
2518H	JMP HOME	8	Move to Track 00
251BH	JMP SELDSK	9	Select Disk Drive
251EH	JMP SETTRK	10	Set Track Number
2521H	JMP SETSEC	11	Set Sector Number
2524H	JMP SETDMA	12	Set DMA Offset Address
2527H	JMP READ	13	Read Selected Sector
252AH	JMP WRITE	14	Write Selected Sector
252DH	JMP LISTST	15	Return List Status
2530H	JMP SECTTRAN	16	Sector Translate
2533H	JMP SETDMAB	17	Set DMA Segment Address
2536H	JMP GETSEGB	18	Get MEM DESC Table Offset
2539H	JMP GETIOB	19	Get I/O Mapping Byte
253CH	JMP SETIOB	20	Set I/O Mapping Byte

There are three major divisions in the BIOS jump table: system (re)initialization subroutines, simple character I/O subroutines, and disk I/O subroutines.

SIMPLE PERIPHERAL DEVICES

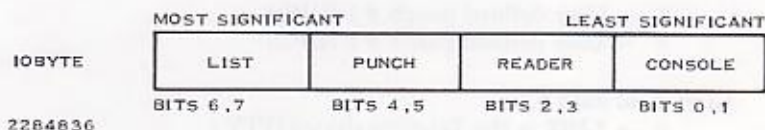
All simple character I/O operations are assumed to be performed in ASCII, uppercase and lowercase, with high order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII **CTRL-Z** (1AH). Peripheral devices are seen by CP/M-86 as logical devices, and are assigned to physical devices within the BIOS. Device characteristics are defined in the following table:

CP/M-86 Logical Device Characteristics

Device Name	Characteristics
CONSOLE	The principal interactive console which communicates with the operator, accessed through CONST, CONIN, and CONOUT. Typically, the CONSOLE is a device such as a CRT or Teletype.
LIST	The principal listing device, if it exists on your system, which is usually a hard-copy device, such as a printer or Teletype.
PUNCH	The principal tape punching device, if it exists, which is normally a high-speed paper tape punch or Teletype.
READER	The principal tape reading device, such as a simple optical reader or teletype.

Note that a single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, your CBIOS should give an appropriate error message so that the system does not hang if the device is accessed by PIP or some other transient program. Alternately, the PUNCH and LIST subroutines can just simply return, and the READER subroutine can return with a 1AH (CTRL-Z) in register A to indicate immediate end-of-file.

For added flexibility, you can optionally implement the IOBYTE function which allows reassignment of physical and logical devices. The IOBYTE function creates a mapping of logical to physical devices which can be altered during CP/M-86 processing (see the STAT command). The definition of the IOBYTE function corresponds to the Intel standard as follows: a single location in the BIOS is maintained, called IOBYTE, which defines the logical to physical device mapping which is in effect at a particular time. The mapping is performed by splitting the IOBYTE into four distinct fields of two bits each, called the CONSOLE, READER, PUNCH, and LIST fields, as shown below:



The value in each field can be in the range 0 - 3, defining the assigned source or destination of each logical device. The values which can be assigned to each field are given in the following table:

IOBYTE Field Definitions

CONSOLE field (bits 0,1)

- 0 - Console is assigned to the console printer (TTY:)
- 1 - Console is assigned to the CRT device (CRT:)
- 2 - Batch mode: use the READER as the CONSOLE input, and the LIST device as the CONSOLE output (BAT:)
- 3 - User defined console device (UC1:)

READER field (bits 2,3)

- 0 - READER is the Teletype device (TTY:)
- 1 - READER is the high-speed reader device (RDR:)
- 2 - User defined reader # 1 (UR1:)
- 3 - User defined reader # 2 (UR2:)

PUNCH field (bits 4,5)

- 0 - PUNCH is the Teletype device (TTY:)
- 1 - PUNCH is the high speed punch device (PUN:)
- 2 - User defined punch # 1 (UP1:)
- 3 - User defined punch # 2 (UP2:)

LIST field (bits 6,7)

- 0 - LIST is the Teletype device (TTY:)
- 1 - LIST is the CRT device (CRT:)
- 2 - LIST is the line printer device (LPT:)
- 3 - User defined list device (UL1:)

Note again that the implementation of the IOBYTE is optional, and affects only the organization of your CBIOS. No CP/M-86 utilities use the IOBYTE except for PIP which allows access to the physical devices, and STAT which allows logical-physical assignments to be made and displayed. In any case, you should omit the IOBYTE implementation until your basic CBIOS is fully implemented and tested, then add the IOBYTE to increase your facilities.

BIOS SUBROUTINE ENTRY POINTS

The actions which must take place upon entry to each BIOS subroutine are given below. It should be noted that disk I/O is always performed through a sequence of calls on the various disk access subroutines. These set up the disk number to access, the track and sector on a particular disk, and the direct memory access (DMA) offset and segment addresses involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation. Note that there is often a single call to SELDSK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there may be a call to set the DMA segment base and a call to set the DMA offset followed by several calls which read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

The READ and WRITE subroutines should try several times (10 is standard) before reporting the error condition to the BDOS. The HOME subroutine may or may not actually perform the track 00 seek, depending upon your controller characteristics; the important point is that track 00 has been selected for the next operation, and is often treated in exactly the same manner as SETTRK with a parameter of 00.

BIOS Subroutine Summary

Subroutine	Description
INIT	This subroutine is called directly by the CP/M-86 loader after the CPM.SYS file has been read into memory. The procedure is responsible for any hardware initialization not performed by the bootstrap loader, setting initial values for BIOS variables (including IOBYTE), printing a sign-on message, and initializing the interrupt vector to point to the BDOS offset (0B11H) and base. When this routine completes, it jumps to the CCP offset (0H). All segment registers should be initialized at this time to contain the base of the operating system.
WBOOT	This subroutine is called whenever a program terminates by performing a BDOS function #0 call. Some re-initialization of the hardware or software may occur here. When this routine completes, it jumps directly to the warm start entry point of the CCP (06H).
CONST	Sample the status of the currently assigned console device and return 0FFH in register AL if a character is ready to read, and 00H in register AL if no console characters are ready.
CONIN	Read the next console character into register AL, and set the parity bit (high order bit) to zero. If no console character is ready, wait until a character is typed before returning.

BIOS Subroutine Summary (Continued)

Subroutine	Description
CONOUT	Send the character from register CL to the console output device. The character is in ASCII, with high order parity bit set to zero. You may want to include a time-out on a line feed or carriage return, if your console device requires some time interval at the end of the line (such as a TI Silent 700™ terminal). You can, if you wish, filter out control characters which have undesirable effects on the console device.
LIST	Send the character from register CL to the currently assigned listing device. The character is in ASCII with zero parity.
PUNCH	Send the character from register CL to the currently assigned punch device. The character is in ASCII with zero parity.
READER	Read the next character from the currently assigned reader device into register AL with zero parity (high order bit must be zero). An end of file condition is reported by returning an ASCII CTRL-Z (1AH).
HOME	Return the disk head of the currently selected disk to the track 00 position. If your controller does not have a special feature for finding track 00, you can translate the call into a call to SETTRK with a parameter of 0.

Silent 700 is a trademark of Texas Instruments Incorporated.

BIOS Subroutine Summary (Continued)

Subroutine	Description
SELDSK	Select the disk drive given by register CL for further operations, where register CL contains 0 for drive A, 1 for drive B, and so on up to 15 for drive P (the standard CP/M-86 distribution version supports two drives). On each disk select, SELDSK must return in BX the base address of the selected drive's Disk Parameter Header. For standard floppy disk drives, the content of the header and associated tables does not change. The sample BIOS included with CP/M-86 called CBIOS contains an example program segment that performs the SELDSK function. If there is an attempt to select a non-existent drive, SELDSK returns BX = 0000H as an error indicator. Although SELDSK must return the header address on each call, it is advisable to postpone the actual physical disk select operation until an I/O function (seek, read or write) is performed. This is due to the fact that disk select operations may take place without a subsequent disk operation and thus disk access may be substantially slower using some disk controllers. On entry to SELDSK it is possible to determine whether it is the first time the specified disk has been selected. Register DL, bit 0 (least significant bit) is a zero if the drive has not been previously selected. This information is of interest in systems which read configuration information from the disk in order to set up a dynamic disk definition table.

BIOS Subroutine Summary (Continued)

Subroutine	Description
SETTRK	Register CX contains the track number for subsequent disk accesses on the currently selected drive. You can choose to seek the selected track at this time, or delay the seek until the next read or write actually occurs. Register CX can take on values in the range 0-76 corresponding to valid track numbers for standard floppy disk drives, and 0-65535 for non-standard disk subsystems.
SETSEC	Register CX contains the translated sector number for subsequent disk accesses on the currently selected drive (see SECTTRAN, below). You can choose to send this information to the controller at this point, or instead delay sector selection until a read or write operation occurs.
SETDMA	Register CX contains the DMA (disk memory access) offset for subsequent read or write operations. For example, if CX = 80H when SETDMA is called, then all subsequent read operations read their data into 80H through 0FFH offset from the current DMA segment base, and all subsequent write operations get their data from that address, until the next calls to SETDMA and SETDMAB occur. Note that the controller need not actually support direct memory access. If, for example, all data is received and sent through I/O ports, the CBIOS which you construct will use the 128 byte area starting at the selected DMA offset and base for the memory buffer during the following read or write operations.

BIOS Subroutine Summary (Continued)

Subroutine	Description
READ	<p>Assuming the drive has been selected, the track has been set, the sector has been set, and the DMA offset and segment base have been specified, the READ subroutine attempts to read one sector based upon these parameters, and returns the following error codes in register AL:</p> <ul style="list-style-type: none">0 no errors occurred1 non-recoverable error condition occurred <p>Currently, CP/M-86 responds only to a zero or non-zero value as the return code. That is, if the value in register AL is 0 then CP/M-86 assumes that the disk operation completed properly. If an error occurs, however, the CBIOS should attempt at least 10 retries to see if the error is recoverable. When an error is reported the BDOS will print the message <code>BDOS ERR ON x: BAD SECTOR</code>. The operator then has the option of typing <code>RETURN</code> to ignore the error, or <code>CTRL-C</code> to abort.</p>
WRITE	<p>Write the data from the currently selected DMA buffer to the currently selected drive, track, and sector. The data should be marked as non-deleted data to maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register AL, with error recovery attempts as described above.</p>
LISTST	<p>Return the ready status of the list device. The value 00 is returned in AL if the list device is not ready to accept a character, and 0FFH if a character can be sent to the printer.</p>

BIOS Subroutine Summary (Continued)

Subroutine	Description
SECTRAN	Performs logical to physical sector translation to improve the overall response of CP/M-86. Standard CP/M-86 systems are shipped with a skew factor of 6, where five physical sectors are skipped between sequential read or write operations. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In computer systems that use fast processors, memory and disk subsystems, the skew factor may be changed to improve overall response. Note, however, that you should maintain a single density IBM® compatible version of CP/M-86 for information transfer into and out of your computer system, using a skew factor of 6. In general, SECTRAN receives a logical sector number in CX. This logical sector number may range from 0 to the number of sectors -1. SECTRAN also receives a translate table offset in DX. The sector number is used as an index into the translate table, with the resulting physical sector number in BX. For standard systems, the tables and indexing code is provided in the CBIOS and need not be changed. If DX = 0000H no translation takes place, and CX is simply copied to BX before returning. Otherwise, SECTRAN computes and returns the translated sector number in BX. Note that SECTRAN is called when no translation is specified in the Disk Parameter Header.
SETDMAB	Register CX contains the segment base for subsequent DMA read or write operations. The BIOS will use the 128 byte buffer at the memory address determined by the DMA base and the DMA offset during read and write operations.

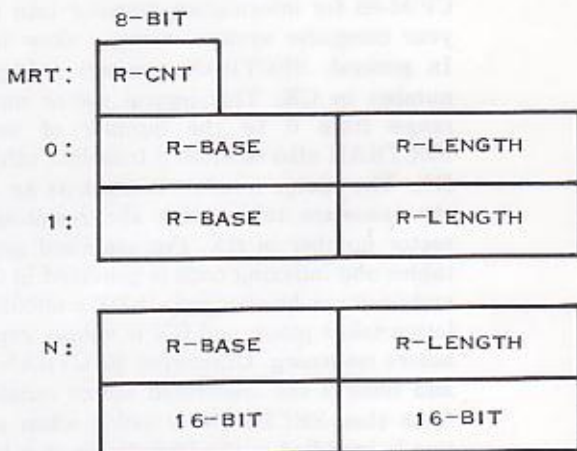
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BIOS Subroutine Summary (Continued)

Subroutine	Description
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GETSEGB	Returns the address of the Memory Region Table (MRT) in BX. The returned value is the offset of the table relative to the start of the operating system. The table defines the location and extent of physical memory which is available for transient programs.
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Memory areas reserved for interrupt vectors and the CP/M-86 operating system are not included in the MRT. The Memory Region Table takes the form:



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where R-Cnt is the number of Memory Region Descriptors (equal to $n+1$ in the diagram above), while R-Base and R-Length give the paragraph base and length of each physically contiguous area of memory. Again, the reserved interrupt locations, normally 0-3FFH, and the CP/M-86 operating system are not included in this map, because the map contains regions available to transient programs. If all memory is contiguous, the R-Cnt field is 1 and $n = 0$, with only a single Memory Region Descriptor which defines the region.

BIOS Subroutine Summary (Continued)

Subroutine	Description
GETIOB	Returns the current value of the logical to physical input/output device byte (IOBYTE) in AL. This eight-bit value is used to associate physical devices with CP/M-86's four logical devices.
SETIOB	Use the value in CL to set the value of the IOBYTE stored in the BIOS.

The following section describes the exact layout and construction of the disk parameter tables referenced by various subroutines in the BIOS.

BIOS Disk Definition Tables

Introduction	6-3
Disk Parameter Table Format	6-3
Table Generation Using GENDEF	6-10
GENDEF Output	6-16

INTRODUCTION

In a manner similar to CP/M-80, CP/M-86 is a table-driven operating system with a separate field-configurable Basic I/O System (BIOS). By altering specific subroutines in the BIOS (presented in the previous chapter), CP/M-86 can be customized for operation on any RAM-based 8086 or 8088 microprocessor system.

The purpose of this section is to present the organization and construction of tables within the BIOS that define the characteristics of a particular disk system used with CP/M-86. These tables can be either hand-coded or automatically generated using the GENDEF utility provided with CP/M-86. The elements of these tables are presented below.

DISK PARAMETER TABLE FORMAT

In general, each disk drive has an associated (16-byte) disk parameter header which both contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The format of the disk parameter header for each drive is as follows:

DISK PARAMETER HEADER							
XLT	0000	0000	0000	DIRBUF	DPB	CSV	ALV
16B	16B	16B	16B	16B	16B	16B	16B

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where each element is a word (16-bit) value. The meaning of each Disk Parameter Header (DPH) element is given in the following table:

Disk Parameter Header Elements

Element	Description
XLT	Offset of the logical to physical translation vector, if used for this particular drive, or the value 0000H if no sector translation takes place (i.e. the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.
0000	Scratchpad values for use within the BDOS (initial value is unimportant).
DIRBUF	Offset of a 128 byte scratchpad, area for directory operations within BDOS. All DPH's address the same scratchpad area.
DPB	Offset of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.
CSV	Offset of a scratchpad area used for software check for changed disks. This offset is different for each DPH.
ALV	Offset of a scratchpad area used by the BDOS to keep disk storage allocation information. This offset is different for each DPH.

Given n disk drives, the DPHs are arranged in a table whose first row of 16 bytes corresponds to drive 0, with the last row corresponding to drive $n-1$. The table thus appears as the following:

00	XLT 00	0000	0000	0000	DIRBUF	DBP 00	CSV 00	ALV 00
01	XLT 01	0000	0000	0000	DIRBUF	DBP 01	CSV 01	ALV 01
$n-1$	XLT $n-1$	0000	0000	0000	DIRBUF	DBP $n-1$	CSV $n-1$	ALV $n-1$

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where the label DPBASE defines the offset of the DPH table relative to the beginning of the operating system.

A responsibility of the SELDSK subroutine, defined in the previous section, is to return the offset of the DPH from the beginning of the operating system for the selected drive. The following sequence of operations returns the table offset, with a 0000H returned if the selected drive does not exist.

NDISKS EQU 4 ;NUMBER OF DISK DRIVES

.....

SELDISK:

```

;SELECT DISK N GIVEN BY CL
MOV BX,0000H ;READY FOR ERR
CPM CL,NDISKS ;N BEYOND MAX DISKS?
JNB RETURN ;RETURN IF SO
;0 <= N < NDISKS
MOV CH,0 ;DOUBLE (N)
MOV BX,CX ;BX = N
MOV CL,4 ;READY FOR * 16
SHL BX,CL ;N = N * 16
MOV CX,OFFSET DPBASE
ADD BX,CX ;DPBASE + N * 16
RETURN: RET ;BX = .DPH (N)

```

The translation vectors (XLT 00 through XLT n-1) are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count-1. The Disk Parameter Block (DPB) for each drive is more complex. A particular DPB, which is addressed by one or more DPHs, takes the general form:

SPT	BSH	BLM	EXM	DSM	DRM	AL0	AL1	CKS	OFF
16B	8B	8B	8B	16B	16B	8B	8B	16B	16B
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where each is a byte or word value, as shown by the 8b or 16b indicator below the field. The fields are defined in the following table:

Disk Parameter Block Fields

Field	Definition
SPT	The total number of sectors per track
BSH	The data allocation block shift factor, determined by the data block allocation size
BLM	The block mask which is also determined by the data block allocation size
EXM	The extent mask, determined by the data block allocation size and the number of disk blocks
DSM	Determines the total storage capacity of the disk drive
DRM	Determines the total number of directory entries which can be stored on this drive
AL0.AL1	Determine reserved directory blocks
CKS	The size of the directory check vector
OFF	The number of reserved tracks at the beginning of the (logical) disk

Although these table values are produced automatically by GENDEF, it is worthwhile reviewing the derivation of each field so that the values may be cross-checked when necessary. The values of BSH and BLM determine (implicitly) the data allocation size BLS, which is not an entry in the disk parameter block. Given that you have selected a value for BLS, the values of BSH and BLM are shown in the following table, with all values in decimal:

BSH and BLM Values for Selected BLS

BLS	BSH	BLM
1,024	3	7
2,048	4	15
4,096	5	31
8,192	6	63
16,384	7	127

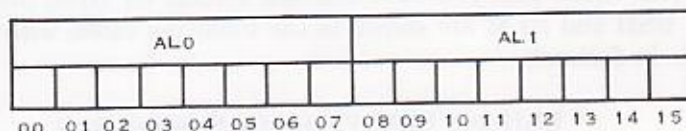
The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255, as shown in the following table.

Maximum EXM Values

BLS	DSM < 256	DSM > 255
1,024	0	N/A
2,048	1	0
4,096	3	1
8,192	7	3
16,384	15	7

The value of DSM is the maximum data block number supported by this particular drive, measured in BLS units. The product BLS times (DSM + 1) is the total number of bytes held by the drive and, of course, must be within the capacity of the physical disk, not counting the reserved operating system tracks.

The DRM entry is one less than the total number of directory entries, which can take on a 16-bit value. The values of AL0 and AL1, however, are determined by DRM. The two values AL0 and AL1 can together be considered a string of 16-bits, as follows:



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where position 00 corresponds to the high order bit of the byte labeled AL0, and 15 corresponds to the low order bit of the byte labeled AL1. Each bit position reserves a data block for a number of directory entries, thus allowing a total of 16 data blocks to be assigned for directory entries (bits are assigned starting at 00 and filled to the right until position 15). Each directory entry occupies 32 bytes, as shown in the following table:

BLS and Number of Directory Entries

BLS	Directory Entries
1,024	32 times # bits
2,048	64 times # bits
4,096	128 times # bits
8,192	256 times # bits
16,384	512 times # bits

Thus, if DRM = 127 (128 directory entries), and BLS = 1024, then there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high order bits of AL0 are set, resulting in the values AL0 = 0F0H and AL1 = 00H.

The CKS value is determined as follows: if the disk drive media is removable, then $CKS = (DRM + 1) / 4$, where DRM is the last directory entry number. If the media is fixed, then set CKS = 0 (no directory records are checked in this case).

Finally, the OFF field determines the number of tracks which are skipped at the beginning of the physical disk. This value is automatically added whenever SETTRK is called, and can be used as a mechanism for skipping reserved operating system tracks, or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the DPB, recall that several DPHs can address the same DPB if their drive characteristics are identical. Further, the DPB can be dynamically changed when a new drive is addressed by simply changing the pointer in the DPH since the BDOS copies the DPB values to a local area whenever the SELDSK function is invoked.

Returning back to the DPH for a particular drive, note that the two address values CSV and ALV remain. Both addresses reference an area of uninitialized memory following the BIOS. The areas must be unique for each drive, and the size of each area is determined by the values in the DPB.

The size of the area addressed by CSV is CKS bytes, which is sufficient to hold the directory check information for this particular drive. If $CKS = (DRM + 1) / 4$, then you must reserve $(DRM + 1) / 4$ bytes for directory check use. If $CKS = 0$, then no storage is reserved.

The size of the area addressed by ALV is determined by the maximum number of data blocks allowed for this particular disk, and is computed as $(DSM / 8) + 1$.

The BIOS shown in Appendix D demonstrates an instance of these tables for standard 8-in single-density drives. It may be useful to examine this program, and compare the tabular values with the definitions given in the preceding paragraphs.

TABLE GENERATION USING GENDEF

The GENDEF utility supplied with CP/M-86 greatly simplifies the table construction process. GENDEF reads a file:

x.DEF

containing the disk definition statements, and produces an output file:

x.LIB

containing assembly language statements which define the tables necessary to support a particular drive configuration. The form of the GENDEF command is:

GENDEF x parameter list

where x has an assumed (and unspecified) filetype of DEF. The parameter list may contain zero or more of the symbols defined in the following table:

GENDEF Optional Parameters

Parameter	Effect
\$C	Generate Disk Parameter Comments
\$O	Generate DPBASE OFFSET \$
\$Z	Z80, 8080, 8085 Override
\$COZ	(Any of the Above)

The C parameter causes GENDEF to produce an accompanying comment line, similar to the output from the STAT DSK: utility which describes the characteristics of each defined disk. Normally, the DPBASE is defined as:

DPBASE EQU \$

which requires a `MOV CX,OFFSET DPBASE` in the `SELDSK` subroutine shown above. For convenience, the `$O` parameter produces the definition

```
DPBASE EQU OFFSET $
```

allowing a `MOV CX,DPBASE` in `SELDSK`, in order to match your particular programming practices. The `$Z` parameter is included to override the standard 8086/8088 mode in order to generate tables acceptable for operation with Z80, 8080, and 8085 assemblers.

The disk definition contained within `x.DEF` is composed with the CP/M text editor, and consists of disk definition statements identical to those accepted by the `DISKDEF` macro supplied with CP/M-80 Version 2. A BIOS disk definition consists of the following sequence of statements:

```
DISKS      n
DISKDEF    0, ...
DISKDEF    1, ...
.....
DISKDEF    n-1
.....
ENDEF
```

Each statement is placed on a single line, with optional embedded comments between the keywords, numbers, and delimiters.

The `DISKS` statement defines the number of drives to be configured with your system, where `n` is an integer in the range of 1 through 16. A series of `DISKDEF` statements then follow which define the characteristics of each logical disk, 0 through `n-1`, corresponding to logical drives A through P. Note that the `DISKS` and `DISKDEF` statements generate the in-line fixed data tables described in the previous section, and thus must be placed in a non-executable portion of your BIOS, typically at the end of your BIOS, before the start of uninitialized RAM.

The `ENDEF` (End of Diskdef) statement generates the necessary uninitialized RAM areas which are located beyond initialized RAM in your BIOS.

The form of the DISKDEF statement is

DISKDEF dn,fsc,lsc, [skf],bls,dks,dir,cks,ofs, [0]

where:

dn is the logical disk number, 0 to n-1
fsc is the first physical sector number (0 or 1)
lsc is the last sector number
skf is the optional sector skew factor
bls is the data allocation block size
dks is the disk size in bls units
dir is the number of directory entries
cks is the number of checked directory entries
ofs is the track offset to logical track 00
[0] is an optional 1.4 compatibility flag

The value dn is the drive number being defined with this DISKDEF statement. The fsc parameter accounts for differing sector numbering systems, and is usually 0 or 1. The lsc parameter is the last numbered sector on a track. When present, the skf parameter defines the sector skew factor which is used to create a sector translation table according to the skew. If the number of sectors is less than 256, a single-byte table is created, otherwise each translation table element occupies two bytes. No translation table is created if the skf parameter is omitted or equal to 0.

The bls parameter specifies the number of bytes allocated to each data block, and takes on the values 1024, 2048, 4096, 8192, or 16384. Generally, performance increases with larger data block sizes because there are fewer directory references. Also, logically connected data records are physically close on the disk. Further, each directory entry addresses more data and the amount of BIOS work space is reduced. The dks parameter specifies the total disk size in bls units. That is, if bls = 2048 and dks = 1000, then the total disk capacity is 2,048,000 bytes. If dks is greater than 255, then the block size parameter bls must be greater than 1024. The value of dir is the total number of directory entries which may exceed 255, if desired.

The cks parameter determines the number of directory items to check on each directory scan, and is used internally to detect changed disks during system operation, where an intervening cold start or system reset has not occurred (when this situation is detected, CP/M-86 automatically marks the disk read/only so that data is not subsequently destroyed). As stated in the previous section, the value of cks = dir when the media is easily changed, as is the case with a floppy disk subsystem. If the disk is permanently mounted, then the value of cks is typically 0, since the probability of changing disks without a restart is quite low.

The ofs value determines the number of tracks to skip when this particular drive is addressed, which can be used to reserve additional operating system space or to simulate several logical drives on a single large capacity physical drive. Finally, the [0] parameter is included when file compatibility is required with versions of CP/M-80, version 1.4, which have been modified for higher density disks (typically double density). This parameter ensures that no directory compression takes place, which would cause incompatibilities with these non-standard CP/M 1.4 versions. Normally, this parameter is not included.

For convenience and economy of table space, the special form

DISKDEF i, i

gives disk i the same characteristics as a previously defined drive i. A standard four-drive single density system, which is compatible with CP/M-80 Version 1.4, and upwardly compatible with CP/M-80 Version 2 implementations, is defined using the following statements:

```
DISKS      4
DISKDEF    0,1,26,6,1024,243,64,64,2
DISKDEF    1,0
DISKDEF    2,0
DISKDEF    3,0
ENDEF
```

with all disks having the same parameter values of 26 sectors per track (numbered 1 through 26), with a skew of 6 between sequential accesses, 1024 bytes per data block, 243 data blocks for a total of 243K byte disk capacity, 64 checked directory entries, and two operating system tracks.

The DISKS statement generates n Disk Parameter Headers (DPHs), starting at the DPH table address DPBASE generated by the statement. Each disk header block contains sixteen bytes, as described above, and corresponds one-for-one to each of the defined drives. In the four drive standard system, for example, the DISKS statement generates a table of the form:

```
DPBASE EQU $
DPE0    DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,
           CSV0,ALV0
DPE1    DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,
           CSV1,ALV1
DPE2    DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,
           CSV2,ALV2
DPE3    DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,
           CSV3,ALV3
```

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive 0 through 3. The values contained within the disk parameter header are described in detail earlier in this section. The check and allocation vector addresses are generated by the ENDEF statement for inclusion in the RAM area following the BIOS code and tables.

Note that if the *skf* (skew factor) parameter is omitted (or equal to 0), the translation table is omitted, and a 0000H value is inserted in the XLT position of the disk parameter header for the disk. In a subsequent call to perform the logical to physical translation, SECTTRAN receives a translation table address of DX = 0000H, and simply returns the original logical sector from CX in the BX register. A translate table is constructed when the *skf* parameter is present, and the (non-zero) table address is placed into the corresponding DPHs. The table shown below, for example, is constructed when the standard skew factor *skf* = 6 is specified in the DISKDEF statement call:

```
XLT0 EOU OFFSET  $
DB 1,7,13,19,25,5,11,17,23,3,9,15,21
DB 2,8,14,20,26,6,12,18,24,4,10,16,22
```

Following the ENDEF statement, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS which is loaded upon cold start, but must be available between the BIOS and the end of operating system memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF statement. For a standard four-drive system, the ENDEF statement might produce:

```
1C72 = BEGDAT EQU OFFSET  $
(data areas)
1DB0 = ENDDAT EQU OFFSET  $
013C = DATSIZ EQU OFFSET  $-BEGDAT
```

which indicates that uninitialized RAM begins at offset 1C72H, ends at 1DB0H-1, and occupies 013CH bytes. You must ensure that these addresses are free for use after the system is loaded.

After modification, you can use the STAT program to check your drive characteristics, since STAT uses the disk parameter block to decode the drive information. The comment included in the LIB file by the \$C parameter to GENCMD will match the output from STAT. The STAT command form

STAT d: DSK:

decodes the disk parameter block for drive d (d = A, . . . ,P) and displays the values shown below:

r:	128 Byte Record Capacity
k:	Kilobyte Drive Capacity
d:	32 Byte Directory Entries
c:	Checked Directory Entries
e:	Records/Extent
b:	Records/Block
s:	Sectors/Track
t:	Reserved Tracks

GENDEF OUTPUT

GENDEF produces a listing of the statements included in the DEF file at the user console (**CTRL-P** can be used to obtain a printed listing, if desired). Each source line is numbered, and any errors are shown below the line in error, with a question mark (?) beneath the item which caused the condition. The source errors produced by GENCMD are listed in the following table, and by errors that can occur when producing input and output files are listed in the table following that.

GENDEF Source Error Messages

Message	Meaning
Bad Val	More than 16 disks defined in DISKS statement
Convert	Number cannot be converted, must be constant in binary, octal, decimal, or hexadecimal as in ASM-86
Delimit	Missing delimiter between parameters
Duplic	Duplicate definition for a disk drive
Extra	Extra parameters occur at the end of line
Length	Keyword or data item is too long
Missing	Parameter required in this position
No Disk	Referenced disk not previously defined
No Stmt	Statement keyword not recognized
Numeric	Number required in this position
Range	Number in this position is out of range
Too Few	Not enough parameters provided
Quote	Missing end quote on current line

GENDEF Input and Output Error Messages

Message	Meaning
Cannot Close .LIB File	LIB™ file close operation unsuccessful, usually due to hardware write protect
LIB Disk Full	No space for LIB file
No Input File present	Specified DEF file not found
No .LIB Directory Space	Cannot create LIB file due to too many files on LIB disk
Premature End-of-File	End of DEF file encountered unexpectedly

Given the file TWO.DEF containing the following statements:

```
disks      2
diskdef    0,1,26,6,2048,256,128,128,2
diskdef    1,1,58,2048,1024,300,0,2
endef
```

the command:

```
gencmd two $c
```

produces the console output:

```
DISKDEF Table Generator, Vers 1.0
1 DISKS 2
2 DISKDEF 0,1,58,2048,256,128,128,2
3 DISKDEF 1,1,58,2048,1024,300,0,2
4 ENDEF
No Error(s)
```

The resulting TWO.LIB file is brought into the following skeletal assembly language program, using the ASM-86 INCLUDE directive. The ASM-86 output listing is truncated on the right, but can be easily reproduced using GENDEF and ASM-86.

LIB is a trademark of Digital Research.

```

=           ;           2048: Kilobyte Drive Capacity
=           ;           300: 32 Byte Directory Entries
=           ;           0: Checked Directory Entries
=           ;           128: Records / Extent
=           ;           16: Records / Block
=           ;           58: Sectors / Track
=           ;           2: Reserved Tracks
=           ;
= 004C      dpbl      equ      offset $           ;Disk P
= 004C 3A 00      dw       58                   ;Sector
= 004E 04      db       4                       ;Block
= 004F 0F      db       15                      ;Block
= 0050 00      db       0                       ;Extnt
= 0051 FF 03      dw       1023                  ;Disk S
= 0053 2B 01      dw       299                   ;Direct
= 0055 F8      db       248                      ;Alloc0
= 0056 00      db       0                       ;Alloc1
= 0057 00 00     dw       0                       ;Check
= 0059 02 00     dw       2                       ;Offset
= 0000      xltl      equ      0                   ;No Tra
= 0080      als1      equ      128                  ;Alloca
= 0000      css1      equ      0                   ;Check
=           ;           ENDEF
=           ;
=           ;           Uninitialized Scratch Memory Fo
=           ;
= 005B      beqdat    equ      offset $           ;Start
= 005B      dirbuf    rs       128                 ;Direct
= 00DB      alv0      rs       als0                ;Alloc
= 00FB      csv0      rs       css0                ;Check
= 011B      alv1      rs       als1                ;Alloc
= 019B      csv1      rs       css1                ;Check
= 019B      enddat    equ      offset $           ;End of
= 0140      datsiz    equ      offset $-beqdat     ;Size o
= 019B 00      db       0                       ;Marks
=           ;           END

```


CP/M-86 Bootstrap and Adaptation Procedures

Introduction	7-3
The Cold Start Load Operation	7-4
Organization of CPM. SYS	7-8

INTRODUCTION

This section describes the components of the standard CP/M-86 distribution disk, the operation of each component, and the procedures to follow in adapting CP/M-86 to non-standard hardware.

CP/M-86 is distributed on a single-density IBM compatible 8-inch diskette using a file format which is compatible with all previous CP/M-80 operating systems. In particular, the first two tracks are reserved for operating system and bootstrap programs, while the remainder of the diskette contains directory information which leads to program and data files. CP/M-86 is distributed for operation with the Intel SBC 86/12 single-board computer connected to floppy disks through an Intel 204 Controller. The operation of CP/M-86 on this configuration serves as a model for other 8086 and 8088 environments, and is presented below.

The principal components of the distribution system are listed below:

- The 86/12 Bootstrap ROM (BOOT ROM)
- The Cold Start Loader (LOADER)
- The CP/M-86 System (CPM.SYS)

When installed in the SBC 86/12, the BOOT ROM becomes a part of the memory address space, beginning at byte location 0FF000H, and receives control when the system reset button is depressed. In a non-standard environment, the BOOT ROM is replaced by an equivalent initial loader and, therefore, the ROM itself is not included with CP/M-86. The BOOT ROM can be obtained from Digital Research or, alternatively, it can be programmed from the listing given in Appendix C or directly from the source file which is included on the distribution disk as BOOT.A86. The responsibility of the BOOT ROM is to read the LOADER from the first two system tracks into memory and pass program control to the LOADER for execution.

THE COLD START LOAD OPERATION

The LOADER program is a simple version of CP/M-86 that contains sufficient file processing capability to read CPM.SYS from the system disk to memory. When LOADER completes its operation, the CPM.SYS program receives control and proceeds to process operator input commands.

Both the LOADER and CPM.SYS programs are preceded by the standard CMD header record. The 128-byte LOADER header record contains the following single group descriptor.

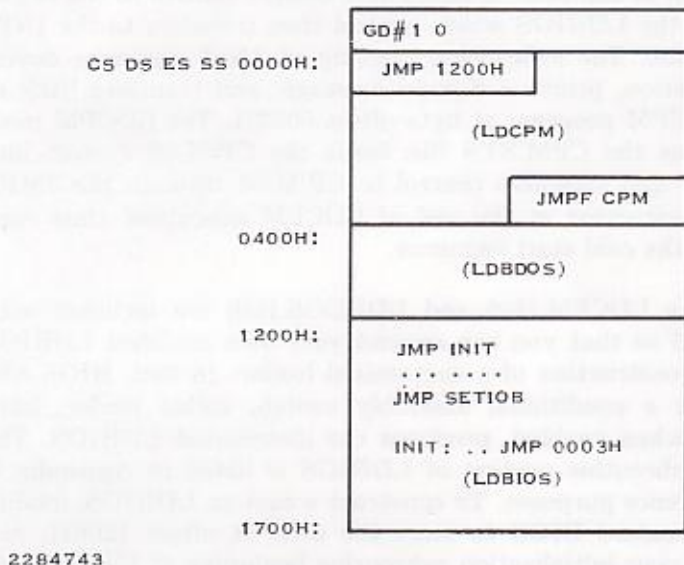
G-FORM	G-LENGTH	A-BASE	G-MIN	G-MAX
1	XXXXXXXXXX	0400	XXXXXXXX	XXXXXXXX

8B 16B 16B 16B 16B

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where G-Form = 1 denotes a code group, x fields are ignored and A-Base defines the paragraph address where the BOOT ROM begins filling memory (A-Base is the word value which is offset three bytes from the beginning of the header). Note that since only a code group is present, an 8080 memory model is assumed. Further, although the A-Base defines the base paragraph address for LOADER (byte address 04000H), the LOADER can, in fact, be loaded and executed at any paragraph boundary that does not overlap CP/M-86 or the BOOT ROM.

The LOADER itself consists of three parts: the Load CPM program (LDCPM), the Loader Basic Disk System (LDBDOS), and the Loader Basic I/O System (LDBIOS). Although the LOADER is setup to initialize CP/M-86 using the Intel 86/12 configuration, the LDBIOS can be field-altered to account for non-standard hardware using the same entry points described in a previous section for BIOS modification. The organization of LOADER is shown in the following figure:



Byte offsets from the base registers are shown at the left of the diagram. GD#1 is the Group Descriptor for the LOADER code group described above, followed immediately by a 0 group terminator. The entire LOADER program is read by the BOOT ROM, excluding the header record, starting at byte location 04000H as given by the A-Field. Upon completion of the read, the BOOT ROM passes control to location 04000H where the LOADER program commences execution. The JMP 1200H instruction at the base of LDCPM transfers control to the beginning of the LDBIOS where control then transfers to the INIT subroutine. The subroutine starting at INIT performs device initialization, prints a sign-on message, and transfers back to the LDCPM program at byte offset 0003H. The LDCPM module opens the CPM.SYS file, loads the CP/M-86 system into memory and transfers control to CP/M-86 through the JMPF CPM instruction at the end of LDCPM execution, thus completing the cold start sequence.

The files LDCPM.H86 and LDBDOS.H86 are included with CP/M-86 so that you can append your own modified LDBIOS in the construction of a customized loader. In fact, BIOS.A86 contains a conditional assembly switch, called loader_bios, which, when enabled, produces the distributed LDBIOS. The INIT subroutine portion of LDBIOS is listed in Appendix C for reference purposes. To construct a custom LDBIOS, modify your standard BIOS to start the code at offset 1200H, and change your initialization subroutine beginning at INIT to perform disk and device initialization. Include a JMP to offset 0003H at the end of your INIT subroutine. Use ASM-86 to assemble your LDBIOS.A86 program using the following command:

ASM86 LDBIOS

to produce the LDBIOS.H86 machine code file. Concatenate the three LOADER modules using PIP:

```
PIP LOADER. H86 =  
LDCPM.H86, LDBDOS. H86, LDBIOS. H86
```

to produce the machine code file for the **LOADER** program. Although the standard **LOADER** program ends at offset 1700H, your modified **LDBIOS** may differ from this last address with the restriction that the **LOADER** must fit within the first two tracks and not overlap **CP/M-86** areas. Generate the command (**CMD**) file for **LOADER** using the **GENCMD** utility:

GENCMD LOADER 8080 CODE [A400]

resulting in the file **LOADER.CMD** with a header record defining the 8080 Memory Model with an absolute paragraph address of 400H, or byte address 4000H. Use **DDT** to read **LOADER.CMD** to location 900H in your 8080 system. Then use the 8080 utility **SYSGEN** to copy the loader to the first two tracks of a disk.

A> DDT

- ILOADER.CMD

- R800

- C

A> SYSGEN

SOURCE DRIVE NAME (or return to skip) **< cr >**

DESTINATION DRIVE NAME (or return to skip) **B**

Alternatively, if you have access to an operational **CP/M-86** system, the command

LDCOPY LOADER

copies **LOADER** to the system tracks. You now have a diskette with a **LOADER** program which incorporates your custom **LDBIOS** capable of reading the **CPM.SYS** file into memory. For standardization, we assume **LOADER** executes at location 4000H. **LOADER** is statically relocatable, however, and its operating address is determined only by the value of **A-Base** in the header record.

You must, of course, perform the same function as the BOOT ROM to get LOADER into memory. The boot operation is usually accomplished in one of two ways. First, you can program your own ROM (or PROM) to perform a function similar to the BOOT ROM when your computer's reset button is pushed. As an alternative, most controllers provide a power-on boot operation that reads the first disk sector into memory. This one-sector program, in turn, reads the LOADER from the remaining sectors and transfers to LOADER upon completion, thereby performing the same actions as the BOOT ROM. Either of these alternatives is hardware-specific, so you will need to be familiar with the operating environment.

ORGANIZATION OF CPM.SYS

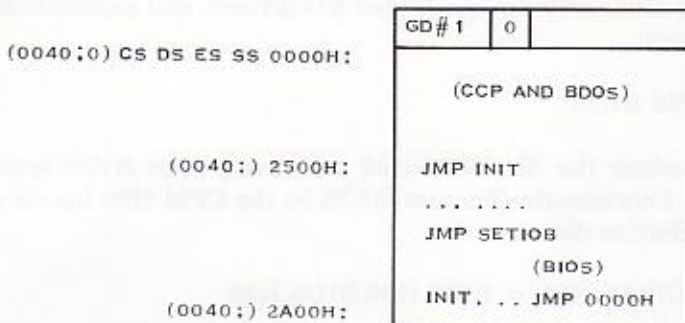
The CPM.SYS file, read by the LOADER program, consists of the CCP, BDOS, and BIOS in CMD file format, with a 128-byte header record similar to that of the LOADER program:

G-FORM	G-LENGTH	A-BAS	G-MIN	G-MAX
1	XXXXXXXXXX	040	XXXXXXXX	XXXXXXXX

8B 16B 16B 16B 16B

2284744

where, instead, the A-Base load address is paragraph 040H, or byte address 0400H, immediately following the 8086 interrupt locations. The entire CPM.SYS file appears on disk as shown in the following figure:



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where GD#1 is the Group Descriptor containing the A-Base value followed by a 0 terminator. The distributed 86/12 BIOS is listed in Appendix D, with an include statement that reads the SINGLES.LIB file containing the disk definition tables. The SINGLES.LIB file is created by GENDEF using the SINGLES.DEF statements shown below:

```
disks 2
diskdef 0,1,26,6,1024,243,64,64,2
diskdef 1,0
endef
```

The CPM.SYS file is read by the LOADER program beginning at the address given by A-Base (byte address 0400H), and control is passed to the INIT entry point at offset address 2500H. Any additional initialization not performed by LOADER takes place in the INIT subroutine and, upon completion, INIT executes a JMP 0000H to begin execution of the CCP. The actual load address of CPM.SYS is determined entirely by the address given in the A-Base field which can be changed if you wish to execute CP/M-86 in another region of memory. Note that the region occupied by the operating system must be excluded from the BIOS memory region table.

In a manner similar to the `LOADER` program, you can modify the BIOS by altering either the `BIOS.A86` or skeletal `CBIOS.A86` assembly language files which are included on your source disk. In either case, create a customized BIOS which includes your specialized I/O drivers, and assemble using `ASM-86`:

ASM86 BIOS

to produce the file `BIOS.H86` containing your BIOS machine code. Concatenate this new BIOS to the `CPM.H86` file on your distribution disk:

```
PIP CPMX.H85 = CPM.H86,BIOS.H86
```

The resulting CPMX hex file is then converted to CMD file format by executing:

```
GENCMD CPMX 8080 CODE [A40]
```

in order to produce the CMD memory image with A-Base = 40H. Finally, rename the CPMX file using the command:

```
REN CPM.SYS = CPMX.CMD
```

and place this file on your 8086 system disk. Now the tailoring process is complete: you have replaced the `BOOT ROM` by either your own customized `BOOT ROM`, or a one-sector cold start loader which brings the `LOADER` program, with your custom `LDBIOS`, into memory at byte location 04000H. The `LOADER` program, in turn, reads the `CPM.SYS` file, with your custom BIOS, into memory at byte location 0400H. Control transfers to `CP/M-86`, and you are up and operating. `CP/M-86` remains in memory until the next cold start operation takes place.

You can avoid the two-step boot operation if you construct a non-standard disk with sufficient space to hold the entire CPM.SYS file on the system tracks. In this case, the cold start brings the CP/M-86 memory image into memory at the location given by A-Base, and control transfers to the INIT entry point at offset 2500H. Thus, the intermediate LOADER program is eliminated entirely, although the initialization found in the LDBIOS must, of course, take place instead within the BIOS.

Since ASM-86, GENCMD and GENDEF are provided in both COM and CMD formats, either CP/M-80 or CP/M-86 can be used to aid the customizing process. If CP/M-80 or CP/M-86 is not available, but you have minimal editing and debugging tools, you can write specialized disk I/O routines to read and write the system tracks, as well as the CPM.SYS file.

The two system tracks are simple to access, but the CPM.SYS file is somewhat more difficult to read. CPM.SYS is the first file on the disk and thus it appears immediately following the directory on the diskette. The directory begins on the third track, and occupies the first sixteen logical sectors of the diskette, while the CPM.SYS is found starting at the seventeenth sector. Sectors are skewed by a factor of six, beginning with the directory track (the system tracks are sequential), so that you must load every sixth sector in reading the CPM.SYS file. Clearly, it is worth the time and effort to use an existing CP/M system to aid the conversion process.

Sector Blocking and Deblocking

Upon each call to the BIOS WRITE entry point, the CP/M-86 BDOS includes information that allows effective sector blocking and deblocking where the host disk subsystem has a sector size which is a multiple of the basic 128-byte unit. This appendix presents a general-purpose algorithm that can be included within your BIOS and that uses the BDOS information to perform the operations automatically.

Upon each call to WRITE, the BDOS provides the following information in register CL:

- 0 = normal sector write
- 1 = write to directory sector
- 2 = write to the first sector of a new data block

Condition 0 occurs whenever the next write operation is into a previously written area, such as a random mode record update, when the write is to other than the first sector of an unallocated block, or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128-byte sectors in sequence, and thus there is little overhead involved in either operation when blocking and deblocking records, since pre-read operations can be avoided when writing records.

This appendix lists the blocking and deblocking algorithm in skeletal form (the file is included on your CP/M-86 disk). Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer which is the size of the host disk sector. Throughout the program, values and variables which relate to the CP/M sector involved in a seek operation are prefixed by `sek`, while those related to the host disk system are prefixed by `hst`. The equate statements beginning on line 24 of Appendix F define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The `SELDSK` entry point clears the host buffer flag whenever a new disk is logged-in. Note that although the `SELDSK` entry point computes and returns the Disk Parameter Header address, it does not physically select the host disk at this point (it is selected later at `READHST` or `WRITEHST`). Further, `SETTRK`, `SETSEC`, and `SETDMA` simply store the values, but do not take any other action at this point. `SECTRAN` performs trivial function of returning the physical sector number.

The principal entry points are `READ` and `WRITE`. These subroutines take the place of your previous `READ` and `WRITE` operations.

The actual physical read or write takes place at either `WRITEHST` or `READHST`, where all values have been prepared: `hstdsk` is the host disk number, `hstrk` is the host track number, and `hstsec` is the host sector number (which may require translation to a physical sector number). You must insert code at this point which performs the full host sector read or write into, or out of, the buffer at `hstbuf` of length `hstsiz`. All other mapping functions are performed by the algorithms.

Sample Random Access Program

This appendix contains a rather extensive and complete example of random access operation. The program listed here performs the simple function of reading or writing random records upon command from the terminal. Given that the program has been created, assembled, and placed into a file labelled RANDOM.COM, the CCP level command:

```
RANDOM X.DAT
```

starts the test program. The program looks for a file by the name X.DAT (in this particular case) and, if found, proceeds to prompt the console for input. If not found, the file is created before the prompt is given. Each prompt takes the form:

```
next command?
```

and is followed by operator input, terminated by a carriage return. The input commands take the form:

```
nW nR Q
```

where n is an integer value in the range 0 to 65535, and W, R, and Q are simple command characters corresponding to random write, random read, and quit processing, respectively. If the W command is issued, the RANDOM program issues the prompt

```
type data:
```

The operator then responds by typing from 1 to 27 characters, followed by a carriage return. RANDOM then writes the character string into the X.DAT file at record n. If the R command is issued, RANDOM reads record number n and displays the string value at the console. If the Q command is issued, the X.DAT file is closed, and the program returns to the console command processor. The only error message is:

error, try again

The program begins with an initialization section where the input file is opened or created, followed by a continuous loop at the label "ready" where the individual commands are interpreted. The default file control block at offset 005CH and the default buffer at offset 0080H are used in all disk operations. The utility subroutines then follow, which contain the principal input line processor, called readc. This particular program shows the elements of random access processing, and can be used as the basis for further program development. In fact, with some work, this program could evolve into a simple data base management system.

One could, for example, assume a standard record size of 128 bytes, consisting of arbitrary fields within the record. A program, called GETKEY, could be developed which first reads a sequential file and extracts a specific field defined by the operator. For example, the command:

```
GETKEY NAMES.DAT LASTNAME 10 20
```

would cause GETKEY to read the data base file NAMES.DAT and extract the LASTNAME field from each record, starting at position 10 and ending at character 20. GETKEY builds a table in memory consisting of each particular LASTNAME field, along with its 16-bit record number location within the file. The GETKEY program then sorts this list, and writes a new file, called LASTNAME.KEY, which is an alphabetical list of LASTNAME fields with their corresponding record numbers. (This list is called an inverted index in information retrieval parlance.)

Rename the program shown above as QUERY, and enhance it a bit so that it reads a sorted key file into memory. The command line might appear as:

```
QUERY NAMES.DAT LASTNAME. KEY
```

Instead of reading a number, the QUERY program reads an alphanumeric string which is a particular key to find in the NAMES.DAT data base. Since the LASTNAME.KEY list is sorted, you can find a particular entry quite rapidly by performing a binary search, similar to looking up a name in the telephone book. That is, starting at both ends of the list, you examine the entry halfway in between and, if not matched, split either the upper half or the lower half for the next search. You'll quickly reach the item you are looking for (in $\log_2(n)$ steps) where you will find the corresponding record number. Fetch and display this record at the console, just as we have done in the program shown above.

At this point you are just getting started. With a little more work, you can allow a fixed grouping size which differs from the 128 byte record shown above. This is accomplished by keeping track of the record number as well as the byte offset within the record. Knowing the group size, you randomly access the record containing the proper group, offset to the beginning of the group within the record read sequentially, until the group size has been exhausted.

Finally, you can improve QUERY considerably by allowing Boolean expressions which compute the set of records which satisfy several relationships, such as a LASTNAME between HARDY and LAUREL, and an AGE less than 45. Display all the records which fit this description. Finally, if your lists are getting too big to fit into memory, randomly access your key files from the disk as well.

```

1: ;
2: ;*****
3: ;*
4: ;*   Sample Random Access Program for CP/M-86   *
5: ;*
6: ;*****
7: ;
8: ;   BIOS Functions
9: ;
10: coninp equ 1 ;console input function
11: conout equ 2 ;console output function
12: pstring equ 9 ;print string until '$'
13: rstring equ 10 ;read console buffer
14: version equ 12 ;return version number
15: openf equ 15 ;file open function
16: closef equ 16 ;close function
17: makef equ 22 ;make file function
18: readr equ 33 ;read random
19: writerr equ 34 ;write random
20: ;
21: ; Equates for non graphic characters
22: cr equ 0dh ;carriage return
23: lf equ 0ah ;line feed
24: ;
25: ;
26: ; load SP, ready file for random access
27: ;
28: cseq
29: pushf ;push flags in CCP stack
30: pop ax ;save flags in AX
31: cli ;disable interrupts
32: mov bx,ds ;set SS register to base
33: mov ss,bx ;set SS, SP with interrupt
34: mov sp,offset stack ; for 80888
35: push ax ;restore the flags
36: popf
37: ;
38: ; CP/M-86 initial release returns the file
39: ; system version number of 2.2; check is
40: ; shown below for illustration purposes.
41: ;
42: mov cl,version
43: call bdos
44: cmp al,20h ;version 2.0 or later?
45: jnb versok
46: ; bad version, message and go back
47: mov dx,offset badver
48: call print
49: jmp abort
50: ;
51: versok:
52: ; correct version for random access
53: mov cl,openf ;open default fct
54: mov dx,offset fcb
55: call bdos

```



```

56:      inc     al                ;err 255 becomes zero
57:      jnz     ready
58: ;
59: ;      cannot open file, so create it
60:      mov     cl,makef
61:      mov     dx,offset fcb
62:      call    bdos
63:      inc     al                ;err 255 becomes zero
64:      inz     ready
65: ;
66: ;      cannot create file, directory full
67:      mov     dx,offset nospace
68:      call    print
69:      jmp     abort            ;back to ccp
70: ;
71: ;      loop back to "ready" after each command
72: ;
73: ready:
74: ;      file is ready for processing
75: ;
76:      call    readcom          ;read next command
77:      mov     ranrec,dx        ;store input record#
78:      mov     ranovf,0h        ;clear high byte if set
79:      cmp     al,'0'          ;quit?
80:      jnz     notq
81: ;
82: ;      quit processing, close file
83:      mov     cl,closef
84:      mov     dx,offset fcb
85:      call    bdos
86:      inc     al                ;err 255 becomes 0
87:      jz      error           ;error message, retry
88:      jmps    abort           ;back to ccp
89: ;
90: ;
91: ;      end of quit command, process write
92: ;
93: ;
94: notq:
95: ;      not the quit command, random write?
96:      cmp     al,'W'
97:      inz     notw
98: ;
99: ;      this is a random write, fill buffer until cr
100:     mov     dx,offset datmsq
101:     call    print            ;data prompt
102:     mov     cx,127          ;up to 127 characters
103:     mov     bx,offset buff   ;destination
104: rloop: ;read next character to buff
105:     push    cx              ;save loop control
106:     push    bx              ;next destination
107:     call    getchr          ;character to AL
108:     pop     bx              ;restore destination
109:     pop     cx              ;restore counter
110:     cmp     al,cr          ;end of line?

```

```

111:         jz      erloop
112: ;       not end, store character
113: mov     byte ptr [bx],al
114: inc     bx           ;next to fill
115: loop   rloop        ;decrement cx ..loop if
116: erloop:
117: ;       end of read loop, store 00
118: mov     byte ptr [bx],0h
119: ;
120: ;       write the record to selected record number
121: mov     cl,writer
122: mov     dx,offset fcb
123: call   bdos
124: or     al,al         ;error code zero?
125: jz     ready        ;for another record
126: jmps   error        ;message if not
127: ;
128: ;
129: ;
130: ;       end of write command, process read
131: ;
132: ;
133: notw:
134: ;       not a write command, read record?
135: cmp     al,"R"
136: jz     ranread
137: jmps   error        ;skip if not
138: ;
139: ;       read random record
140: ranread:
141: mov     cl,readr
142: mov     dx,offset fcb
143: call   bdos
144: or     al,al         ;return code 00?
145: jz     readok
146: jmps   error
147: ;
148: ;       read was successful, write to console
149: readok:
150: call   crlf         ;new line
151: mov     cx,128      ;max 128 characters
152: mov     si,offset buff ;next to get
153: wloop:
154: lods   al           ;next character
155: and    al,07fh      ;mask parity
156: inc    wloopl
157: jmp   ready        ;for another command if
158: wloopl:
159: push  cx           ;save counter
160: push  si           ;save next to get
161: cmp   al," "       ;graphic?
162: jb   skipw        ;skip output if not graphic
163: call putchr        ;output character
164: skipw:
165: pop   si

```

```

166:      pop     cx
167:      loop   wloop           ;decrement CX and check
168:      jmp    ready
169: ;
170: ;
171: ; end of read command, all errors end-up here
172: ;
173: ;
174: error:
175:      mov     dx,offset errmsg
176:      call   @rint
177:      jmp    ready
178: ;
179: ; BIOS entry subroutine
180: bdos:
181:      int     224           ;entry to BIOS if by INT
182:      ret
183: ;
184: abort:
185:      mov     cl,0           ;return to CCP
186:      call   bdos           ;use function 0 to end e
187: ;
188: ; utility subroutines for console i/o
189: ;
190: getchr:
191:      ;read next console character to a
192:      mov     cl,conino
193:      call   bdos
194:      ret
195: ;
196: putchr:
197:      ;write character from a to console
198:      mov     cl,conout
199:      mov     dl,a          ;character to send
200:      call   bdos           ;send character
201:      ret
202: ;
203: crlf:
204:      ;send carriage return line feed
205:      mov     al,cr         ;carriage return
206:      call   putchr
207:      mov     al,lf         ;line feed
208:      call   putchr
209:      ret
210: ;
211: print:
212:      ;print the buffer addressed by dx until $
213:      push   dx
214:      call   crlf
215:      pop    dx             ;new line
216:      mov     cl,pstring
217:      call   bdos           ;print the string
218:      ret
219: ;
220: readcom:

```

```

221:          ;read the next command line to the conbuf
222:      mov     dx,offset promot
223:      call    print          ;command?
224:      mov     cl,rstring
225:      mov     dx,offset conbuf
226:      call    bdom          ;read command line
227: ;      command line is present, scan it
228:      mov     ax,0          ;start with 0000
229:      mov     bx,offset conlin
230: readc:  mov     dl,[bx]     ;next command character
231:          inc     bx        ;to next command position
232:          mov     dh,0      ;zero high byte for add
233:          or     dl,dl      ;check for end of command
234:          jnz    getnum
235:          ret
236: ;      not zero, numeric?
237: getnum:
238:      sub     dl,'0'
239:      cmp     dl,10        ;carry if numeric
240:      jnb    endrd
241:      mov     cl,10
242:      mul    cl          ;multiply accumulator by
243:          add     ax,dx      ;*diqit
244:      jmps   readc       ;for another char
245: endrd:
246: ;      end of read, restore value in a and return value
247:      mov     dx,ax        ;return value in dx
248:      mov     al,-1[bx]
249:      cmp     al,'a'      ;check for lower case
250:      jnb    transl
251:      ret
252: transl:  and     al,5fH   ;translate to upper case
253:      ret
254: ;
255: ;
256: ; Template for Page 0 of Data Group
257: ; Contains default PCB and DMA buffer
258: ;
259:      dseg
260:      org     05ch
261: fcb     rb     33        ;default file control bl
262: ranrec  rw     1        ;random record position
263: ranovf  rb     1        ;high order (overflow) b
264: buff    rb     128     ;default DMA buffer
265: ;
266: ; string data area for console messages
267: badver  db     "sorry, you need cp/m version 2$"
268: nospace db     "no directory spaces"
269: datmsg  db     "type data: $"
270: errmsg  db     "error, try again.$"
271: prompt  db     "next command? $"
272: ;
273: ;
274: ; fixed and variable data area
275: ;

```

```
276: conbuf db          conlen ;length of console buffer
277: consiz rs          1      ;resulting size after read
278: conlin rs          32     ;length 32 buffer
279: conlen equ        offset $ - offset consiz
280: ?
281:          rs          31     ;16 level stack
282: stack  rb          1
283:          db          0      ;end byte for GENCMD
284:          end
```

Listing of the Boot ROM

```

*****
*
* This is the original BOOT ROM distributed with CP/M
* for the SBC 86/12 and 204 Controller. The listing
* is truncated on the right, but can be reproduced by
* assembling ROM.A86 from the distribution disk. Note
* that the distributed source file should always be
* referenced for the latest version
*
*****
;
; ROM bootstrap for CP/M-86 on an iSBC86/12
; with the
; Intel SBC 204 Floppy Disk Controller
;
; Copyright (C) 1980,1981
; Digital Research, Inc.
; Box 579, Pacific Grove
; California, 93950
;
;*****
;* This is the BOOT ROM which is initiated
;* by a system reset. First, the ROM moves
;* a copy of its data area to RAM at loca-
;* tion 00000H, then initializes the segment
;* registers and the stack pointer. The
;* various peripheral interface chips on the
;* SBC 86/12 are initialized. The 8251
;* serial interface is configured for a 9600
;* baud asynchronous terminal, and the in-
;* terrupt controller is setup for inter-
;* rupts 10H-17H (vectors at 00040H-0005FH)
;* and edge-triggered auto-EOI (end of in-
;* terrupt) mode with all interrupt levels
;* masked-off. Next, the SBC 204 Diskette
;* controller is initialized, and track 1
;* sector 1 is read to determine the target
;* paragraph address for LOADER. Finally,
;* the LOADER on track 0 sectors 2-26 and
;* track 1 sectors 1-26 is read into the
;* target address. Control then transfers
;* to LOADER. This program resides in two
;* 2716 EPROM's (2K each) at location
;* 0FF000H on the SBC 86/12 CPU board. ROM
;* 0 contains the even memory locations, and
;* ROM 1 contains the odd addresses. BOOT
;* ROM uses RAM between 00000H and 000FFH
;* (absolute) for a scratch area, along with
;* the sector 1 buffer.
;*****

```

```

00FF      true      equ      0ffh
FF00      false     equ      not true
;
00FF      debug     equ      true
;debug = true indicates bootstrap is in same roms
;with SBC 957 "Execution Vehicle" monitor
;at FE00:0 instead of FF00:0
;
000D      cr        equ      13
000A      lf        equ      10
;
;          disk ports and commands
;
00A0      base204   equ      0a0h
00A0      fdccom    equ      base204+0
00A0      fdcstat   equ      base204+0
00A1      fdcparm   equ      base204+1
00A1      fdcrs1t   equ      base204+1
00A2      fdcrst    equ      base204+2
00A4      dmacadr   equ      base204+4
00A5      dmaccont  equ      base204+5
00A6      dmacscan  equ      base204+6
00A7      dmacsadr  equ      base204+7
00A8      dmacmode  equ      base204+8
00A8      dmacstat  equ      base204+8
00A9      fdcse1    equ      base204+9
00AA      fdcsegment equ      base204+10
00AF      reset204  equ      base204+15
;
;actual console baud rate
2580      baud_rate equ      9600
;value for 8253 baud counter
0008      baud      equ      768/(baud_rate/100)
;
000A      csts      equ      0DAh ;i8251 status port
00D8      cdata     equ      0D8h ; " data port
;
00D0      tch0      equ      0D0h ;8253 PIC channel 0
00D2      tch1      equ      tch0+2 ;ch 1 port
00D4      tch2      equ      tch0+4 ;ch 2 port
00D6      tcmd      equ      tch0+6 ;8253 command port
;
00C0      icp1      equ      0C0h ;8259a port 0
00C2      icp2      equ      0C2h ;8259a port 1
;
;
ROMSEG    IF NOT DEBUG EQU      0FF00H ;normal
ENDIF
;
ROMSEG    IF DEBUG EQU      0FE00H ;share prom with SB
ENDIF
FE00
;
;

```

```

; This long jump prom'd in by hand
; cseg 0ffffh ;reset goes to here
; JMPF BOTTOM ;boot is at bottom
; EA 00 00 00 FF ;cs = bottom of pro
; ; ip = 0
; EVEN PROM ODD PROM
; 7F8 - EA 7F8 - 00
; 7F9 - 00 7F9 - 00
; 7FA - FF ;this is not done i
;
FE00 cseg romseq
;
;First, move our data area into RAM at 0000:0200
;
0000 8CC8 mov ax,cs
0002 8ED8 mov ds,ax ;point DS to CS for source
0004 BE3F01 mov SI,drombegin ;start of data
0007 BF0002 mov DI,offset ram_start ;offset of destinat
000A B80000 mov ax,0
000D 8EC0 mov es,ax ;destination segment is 000
000F B9E600 mov CX,data_length ;how much to move i
0012 F3A4 rep movs al,al ;move out of eprom
;
0014 B80000 mov ax,0
0017 8ED8 mov ds,ax ;data segment now in RAM
0019 8ED0 mov ss,ax
001B BC2A03 mov sp,stack_offset ;initialize stack s
001E FC cld ;clear the directio
;
IF NOT DEBUG
;
;Now, initialize the console USART and baud rate
;
mov al,0Eh
out csts,al ;give 8251 dummy mode
mov al,40h
out csts,al ;reset 8251 to accept mode
mov al,4Eh
out csts,al ;normal 8 bit asynch mode,
mov al,37h
out csts,al ;enable Tx & Rx
mov al,0B6h
out tcmd,al ;8253 ch.2 square wave mode
mov ax,baud
out tch2,al ;low of the baud rate
mov al,ah
out tch2,al ;high of the baud rate
;
ENDIF
;
;Setup the 8259 Programmable Interrupt Controller
;
001F B013 mov al,l3h
0021 E6C0 out icpl,al ;8259a ICW 1 8086 mode
0023 B010 mov al,l0h

```



```

0025 F6C2          out icp2,a1      ;8259a ICW 2  vector @ 40-5
0027 B01F          mov al,1Ph
0029 E6C2          out icp2,a1      ;8259a ICW 4  auto ROI mast
002B B0FF          mov al,0FFh
002D E6C2          out icp2,a1      ;8259a OCW 1  mask all leve
;
;Reset and initialize the iSB0 204 Diskette Interfa
;
restart:          ;also come back here on fatal error
002F E6AF          out reset204,AL ;reset iSB0 204 logic and
0031 B001          mov AL,1
0033 E6A2          out fdcrst,AL   ;give 8271 FDC
0035 B000          mov al,0
0037 E6A2          out fdcrst,AL   ; a reset command
0039 BB1502        mov BX,offset specs1
003C E8E100        CALL sendcom    ;program
003F BB1B02        mov BX,offset specs2
0042 E8DB00        CALL sendcom    ; Shugart SA-800 drive
0045 BB2102        mov BX,offset specs3
0048 E8D500        call sendcom    ; characteristics
004B BB1002        homer: mov BX,offset home
004E E85800        CALL execute    ;home drive 0
;
0051 BB2A03        mov bx,sector1  ;offset for first sector DM
0054 B80000        mov ax,0
0057 8EC0          mov es,ax        ;segment " " " "
0059 E8A700        call setup_dma
;
005C B90202        mov bx,offset read0
005F E84700        call execute    ;get T0 S1
;
0062 8E062D03      mov es,ABS
0066 BB0000        mov bx,0         ;get loader load address
0069 E89700        call setup_dma  ;setup DMA to read loader
;
006C BB0602        mov bx,offset read1
006F E83700        call execute    ;read track 0
0072 BB0B02        mov bx,offset read2
0075 E83100        call execute    ;read track 1
;
0078 8C06E802      mov leap_segment,ES
;
007C C706E6020000  setup far jump vector
;
;
;
;
0082 FF2EE602      jmpf dword ptr leap_offset
;
pmsg:
0086 8A0F          mov cl,[BX]
0088 84C9          test cl,cl
008A 7476          jz return
008C E80400        call conout
008F 43           inc BX
0090 E9F3FF        jmp pmsg
;

```

```

conout:
0093 E4DA      in al,csts
0095 A801      test al,1
0097 74FA      jz conout
0099 8AC1      mov al,cl
009B E6D8      out cdata,al
009D C3        ret

;
conin:
009E E4DA      in al,csts
00A0 A802      test al,2
00A2 74FA      jz conin
00A4 E4D8      in al,cdata
00A6 247F      and al,7fh
00A8 C3        ret

;
;
execute:      ;execute command string @ [BX]
              ;<BX> points to length,
              ;followed by Command byte
              ;followed by length-1 parameter bvt

;
00A9 891E0002  mov     lastcom,BX      ;remember what it w
;retrv if not ready
00AD E87000    call    sendcom        ;execute the comman
;now, let's see wha
;of status poll was
;for that command t
;point to command s
00B0 8B1E0002  mov     BX,lastcom     ;get command op cod
00B4 8A4701    mov     AL,1[BX]      ;drop drive code bi
00B7 243F      and     AL,3fh        ;mask if it will be
00B9 B90008    mov     CX,0800h     ;see if interrupt t
00BC 3C2C      cmp     AL,2ch
00BE 720B      jb     exccpoll
00C0 B98080    mov     CX,8080h     ;else we use "not c
00C3 240F      and     AL,0fh       ;unless . . .
00C5 3C0C      cmp     AL,0ch       ;there isn't
00C7 B000      mov     AL,0
00C9 7737      ja     return        ;any result at all

;
exccpoll:    ;poll for bit in b, toggled with c
              in AL,FDCSTAT
              and AL,CH
00CB E4A0      xor     AL,CL ! JZ exccpoll

;
00D3 E4A1      in     AL,fdcrslt    ;get result registe
00D5 241E      and     AL,leh       ;look only at resul
00D7 7429      jz     return        ;zero means it was

;
00D9 3C10      cmp     al,10h
00DB 7513      jne    fatal        ;if other than "Not

;
00DD BB1302    mov     bx,offset rdstat
00E0 E83D00    call    sendcom      ;perform read statu

```

```

rd_poll:
00E3 E4A0      in al,fdc_stat
00E5 A880      test al,80h           ;wait for command n
00E7 75FA      jnz rd_poll
00E9 8B1E0002  mov bx,last_com      ;recover last attem
00ED E9BDFE      jmp retry            ;and try it over ag

;
;fatal:
;fatal error

00F0 B400      mov ah,0
00F2 8BD8      mov bx,ax             ;make 16 bits
00F4 8B9F2702  mov bx,errtbl[BX]
; print appropriate error message
00F8 E88BFF      call msg
00FB E8A0FF      call conin           ;wait for key strik
00FE 58         pop ax               ;discard unused ite
00FF E92DFE      jmp restart         ;then start all ove

;
;return:
0102 C3        RET                  ;return from EXECUT

;
;setupdma:
0103 B004      mov AL,04h
0105 E6A8      out dmacmode,AL     ;enable dmac
0107 B000      mov al,0
0109 E6A5      out dmaccont,AL     ;set first (dummy)
010B B040      mov AL,40h
010D E6A5      out dmaccont,AL     ;force read data mo
010F 8CC0      mov AX,ES
0111 E6AA      out fdcsegment,AL
0113 8AC4      mov AL,AH
0115 E6AA      out fdcsegment,AL
0117 8BC3      mov AX,BX
0119 E6A4      out dmacadr,AL
011B 8AC4      mov AL,AH
011D E6A4      out dmacadr,AL
011F C3        RET

;
;
;sendcom:      ;routine to send a command string t
0120 E4A0      in AL,fdcstat
0122 2480      and AL,80h
0124 75FA      jnz sendcom         ;insure command not busy
0126 8A0F      mov CL,[BX]         ;get count
0128 43         inc BX
0129 8A07      mov al,[BX]         ;point to and fetch command
012B E6A0      out fdccom,AL      ;send command

parmloop:
012D FEC9      dec CL
012F 74D1      jz return           ;see if any (more) paramete
0131 43         inc BX              ;point to next parameter

parmpoll:
0132 E4A0      in AL,fdcstat
0134 2420      and AL,20h
0136 75FA      jnz parmpoll        ;loop until parm not full

```

```

0138 8A07          mov AL,[BX]
013A E6A1          out fdcparm,AL ;outout next parameter
013C E9EEFF        jmp parmloop    ;go see about another
;
;
;          Image of data to be moved to RAM
013F          ; drombegin equ offset S
;
013F 0000          ; clastcom      dw      0000h ;last command
;
0141 03          ; creadstring   db      3 ;length
0142 52          ;               db      52h ;read function code
0143 00          ;               db      0 ;track #
0144 01          ;               db      1 ;sector #
;
0145 04          ; creadtrk0     db      4
0146 53          ;               db      53h ;read multiple
0147 00          ;               db      0 ;track 0
0148 02          ;               db      2 ;sectors 2
0149 19          ;               db      25 ;through 26
;
014A 04          ; creadtrk1     db      4
014B 53          ;               db      53h
014C 01          ;               db      1 ;track 1
014D 01          ;               db      1 ;sectors 1
014E 1A          ;               db      26 ;through 26
;
014F 026900        ; chome0        db      2,69h,0
0152 016C          ; crdstat0      db      1,6ch
0154 05350D        ; cspecs1       db      5,35h,0dh
0157 0808E9        ;               db      08h,08h,0e9h
015A 053510        ; cspecs2       db      5,35h,10h
015D FFFFFFFF      ;               db      255,255,255
0160 053518        ; cspecs3       db      5,35h,18h
0163 FFFFFFFF      ;               db      255,255,255
;
0166 4702          ; cerrtbl       dw      offset er0
0168 4702          ;               dw      offset er1
016A 4702          ;               dw      offset er2
016C 4702          ;               dw      offset er3
016E 5702          ;               dw      offset er4
0170 6502          ;               dw      offset er5
0172 7002          ;               dw      offset er6
0174 7F02          ;               dw      offset er7
0176 9002          ;               dw      offset er8
0178 A202          ;               dw      offset er9
017A B202          ;               dw      offset erA
017C C502          ;               dw      offset erB
017E D302          ;               dw      offset erC
0180 4702          ;               dw      offset erD
0182 4702          ;               dw      offset erE
0184 4702          ;               dw      offset erF
;
0186 0D0A4E756C6C Cer0 db cr,1f,"Null Error ??",0

```

```

204572726F72
203F3F00
0186          Cer1  equ   cer0
0186          Cer2  equ   cer0
0186          Cer3  equ   cer0
0196 0D0A436C6F63 Cer4  db    cr,lf,"Clock Error",0
      6B204572726F
      7200
01A4 0D0A4C617465 Cer5  db    cr,lf,"Late DMA",0
      20444D4100
01AF 0D0A49442043 Cer6  db    cr,lf,"ID CRC Error",0
      524320457272
      6F7200
01BE 0D0A44617461 Cer7  db    cr,lf,"Data CRC Error",0
      204352432045
      72726F7200
01CF 0D0A44726976 Cer8  db    cr,lf,"Drive Not Ready",0
      65204E6F7420
      526561647900
01E1 0D0A57726974 Cer9  db    cr,lf,"Write Protect",0
      652050726F74
      65637400
01F1 0D0A54726B20 CerA  db    cr,lf,"Trk 00 Not Found",0
      3030204E6F74
      20466F756E64
      00
0204 0D0A57726974 CerB  db    cr,lf,"Write Fault",0
      65204661756C
      7400
0212 0D0A53656374 CerC  db    cr,lf,"Sector Not Found",0
      6F72204E6F74
      20466F756E64
      00
0186          CerD  equ   cer0
0186          CerE  equ   cer0
0186          CerF  equ   cer0
;
0225          dromend equ offset $
;
00E6          data_length equ dromend-drombegin
;
;          reserve space in RAM for data area
;          (no hex records generated here)
;
0000          dseg   0
              org    0200h
;
0200          ram_start equ   $
0200          lastcom  rw    1    ;last command
0202          read0    rb    4    ;read track 0 secto
0206          read1    rb    5    ;read T0 S2-26
020B          read2    rb    5    ;read T1 S1-26
0210          home     rb    3    ;home drive 0
0213          rdstat   rb    2    ;read status
0215          specs1   rb    6

```

```

021B      specs2      rb      6
0221      specs3      rb      6
0227      errtbl     rw      16
0247      er0        rb      length cer0      ;16
0247      er1        equ     er0
0247      er2        equ     er0
0247      er3        equ     er0
0257      er4        rb      length cer4      ;14
0265      er5        rb      length cer5      ;11
0270      er6        rb      length cer6      ;15
027F      er7        rb      length cer7      ;17
0290      er8        rb      length cer8      ;18
02A2      er9        rb      length cer9      ;16
02B2      erA       rb      length cerA     ;19
02C5      erB       rb      length cerB     ;14
02D3      erC       rb      length cerC     ;19
0247      erD       equ     er0
0247      erE       equ     er0
0247      erF       equ     er0
;
02E6      leap_offset rw      1
02E8      leap_segment rw      1
;
;
02EA      ;local stack
032A      stack_offset equ     offset $;stack from here do
;
;
032A      sector1   equ     offset $
;
032A      Ty        rb      1
032B      Len       rw      1
032D      Abs       rw      1      ;ABS is all we care
032F      Min       rw      1
0331      Max       rw      1
end

```

LDBIOS Listing

```
*****
*
* This the the LOADER BIOS, derived from the BIOS
* program by enabling the "loader_bios" condi-
* tional assembly switch. The listing has been
* edited to remove portions which are duplicated
* in the BIOS listing which appears in Appendix D
* where ellipses "..." denote the deleted portions
* (the listing is truncated on the right, but can
* be reproduced by assembling the BIOS.A86 file
* provided with CP/M-86)
*
*****
```

```
*****
;
; * Basic Input/Output System (BIOS) for
; * CP/M-86 Configured for iSBC 86/12 with
; * the iSBC 204 Floppy Disk Controller
;
; * (Note: this file contains both embedded
; * tabs and blanks to minimize the list file
; * width for printing purposes. You may wish
; * to expand the blanks before performing
; * major editing.)
;*****
```

```
; Copyright (C) 1980,1981
; Digital Research, Inc.
; Box 579, Pacific Grove
; California, 93950
;
```

```
; (Permission is hereby granted to use
; or abstract the following program in
; the implementation of CP/M, MP/M or
; CP/NET for the 8086 or 8088 Micro-
; processor)
```

```
FFFF true equ -1
0000 false equ not true
```

```

*****
;*
;* Loader_bios is true if assembling the
;* LOADER BIOS, otherwise BIOS is for the
;* CPM.SYS file. Blc_list is true if we
;* have a serial printer attached to BLC8538
;* Bdos_int is interrupt used for earlier
;* versions.
;*
*****

FFFF loader_bios equ true
FFFF blc_list equ true
00E0 bdos_int equ 224 ;reserved BDOS Interrupt

        IF not loader_bios
;-----|
;|
;|
;|
;-----|
        ENDIF ;not loader_bios

        IF loader_bios
;-----|
;|
1200 bios_code equ 1200h ;start of LDBIOS
0003 ccp_offset equ 0003h ;base of CPMLOADER
0406 bdos_ofst equ 0406h ;stripped BDOS entry
;|
;-----|
        ENDIF ;loader_bios
        . . .

        cseg
org ccpoffset

ccp:
org bios_code

*****
;*
;* BIOS Jump Vector for Individual Routines
;*
*****
1200 E93C00 jmp INIT ;Enter from BOOT ROM or LOADER
1203 E96100 jmp WBOOT ;Arrive here from BDOS call 0
        . . .
1239 E96400 jmp GETIOBF ;return I/O map byte (IOBYTE)
123C E96400 jmp SETIOBF ;set I/O map byte (IOBYTE)

```



```

;*****
;*
;* INIT Entry Point, Differs for LDBIOS and *
;* BIOS, according to "Loader_Bios" value *
;*
;*****
INIT:  ;print signon message and initialize hardwa
      mov ax,cs      ;we entered with a JMPF so
      mov ss,ax     ; CS: as the initial value
      mov ds,ax     ; DS:
      mov es,ax     ; and ES:
      ;use local stack during initialization
      mov sp,offset stkbase
      cld          ;set forward direction

      IF      not loader_bios
      ;-----
      ;|
      ;|      ; This is a BIOS for the CPM.SYS file.
      ;|      . . .
      ;-----
      ENDIF    ;not loader_bios

      IF      loader_bios
      ;-----
      ;|
      ;|      ;This is a BIOS for the LOADER
      ;|      push ds      ;save data segment
      ;|      mov ax,0
      ;|      mov ds,ax     ;point to segment zero
      ;|      ;BDOS interrupt offset
      ;|      mov bdos_offset,bdos_ofst
      ;|      mov bdos_segment,CS ;bdos interrupt segment
      ;|      pop ds      ;restore data segment
      ;|
      ;-----
      ENDIF    ;loader_bios

124B 1E      mov bx,offset signon
124C B80000  call pmsg      ;print signon message
124F 8ED8    mov cl,0      ;default to dr A: on coldst
1251 C70680030604 jmp ccp      ;jump to cold start entry o
1257 8C0E8203
125B 1F

125C BB1514  mov bx,offset signon
125F E85A00  call pmsg      ;print signon message
1262 B100    mov cl,0      ;default to dr A: on coldst
1264 E99CED  jmp ccp      ;jump to cold start entry o

1267 E99FED  WBOOT: jmp ccp+6     ;direct entry to CCP at com

      IF      not loader_bios
      ;-----
      ;|
      ;|      . . .
      ;|
      ;-----
      ENDIF    ;not loader_bios

```

```

;*****
;*
;* CP/M Character I/O Interface Routines
;* Console is Usart (i8251a) on i8BC 86/12
;* at ports D8/DA
;*
;*****
CONST:          ;console status
126A E4DA      in al,csts
               . . .
const_ret:     ret          ;Receiver Data Available
1272 C3
CONIN:         ;console input
               call const
               . . .
CONOUT:        ;console output
127D E4DA      in al,csts
               . . .

LISTOUT:       ;list device output
               IF      blc_list
;-----|
1288 E80700   ;|      call LISTST
               ;|      . . .
;-----|
               ENDIF ;blc_list
1291 C3      ret

LISTST:        ;poll list status
               IF      blc_list
;-----|
1292 E441   ;|      in al,lsts
               ;|      . . .
;-----|
               ENDIF ;blc_list
129C C3      ret

PUNCH:        ;not implemented in this configuration
READER:       mov al,lah
129D B01A     ret          ;return EOF for now
129F C3

```

```

GETIOBF:
12A0 B000      mov al,0          ;TTY: for consistency
12A2 C3        ret          ;IOBYTE not implemented

SETIOBF:
12A3 C3        ret          ;iobyte not implemented

zero_ret:
12A4 2400      and al,0
12A6 C3        ret          ;return zero in AL and flag

; Routine to get and echo a console character
; and shift it to upper case

uconecho:
12A7 E8C9FF    call CONIN       ;get a console character
               .
;*****
;*
;*          Disk Input/Output Routines
;*
;*
;*****

SELOSK:
12CA BB0000    ;select disk given by register CL
               mov bx,0000h
               . . .

HOME:          ;move selected disk to home position (Track
12EB C606311500 mov trk,0       ;set disk i/o to track zero
               . . .

SETTRK:        ;set track address given by CX
1300 880E3115  mov trk,c1     ;we only use 8 bits of trac
1304 C3        ret

SETSEC:        ;set sector number given by cx
1305 880E3215  mov sect,c1    ;we only use 8 bits of sect
1309 C3        ret

SECTTRAN:     ;translate sector CX using table at [DX]
130A 8BD9      mov bx,cx
               . . .

SETDMA:        ;set DMA offset given by CX
1311 890E2A15  mov dma_adr,CX
1315 C3        ret

SETDMAB:      ;set DMA segment given by CX
1316 890E2C15  mov dma_seg,CX
131A C3        ret

;
GETSEGT:      ;return address of physical memory table
131B BB3815    mov bx,offset seq_table
131E C3        ret

```

```

;*****
;*
;* All disk I/O parameters are setup: the
;* Read and Write entry points transfer one
;* sector of 128 bytes to/from the current
;* DMA address using the current disk drive
;*
;*****
READ:
131F B012      mov al,12h      ;basic read sector command
1321 EB02      jmps r_w_common

WRITE:
1323 B00A      mov al,0ah      ;basic write sector command

r_w_common:
1325 BB2F15     mov bx,offset io_com ;point to command stri
               . . .

;*****
;*
;* Data Areas
;*
;*****
1415      data_offset equ offset S

               dseg
               org   data_offset      ;contiguous with co

               IF    loader_bios

;-----|
1415 0D0A0D0A  signon db cr,lf,cr,lf
1419 43502F4D2D38 db "CP/M-86 Version 2.2",cr,lf,0
       362056657273
       696F6E20322E
       320D0A00
;-----|
               ENDIF ;loader_bios

               IF    not loader_bios

;-----|
               . . .
;-----|
               ENDIF ;not loader_bios

142F 0D0A486F6D65 bad_hom db cr,lf,"Home Error",cr,lf,0
=
;               include singles.lib ;read in disk definitio
=               DISKS 2

```

```

= 1541          dpbase equ    $          ;Base of Disk Param
=1668 00          db      0          ;Marks End of Modul

1669          loc_stk rw 32 ;local stack for initialization
16A9          stkBase equ offset $

16A9 00          db 0          ;fill last address for GENCMD

;*****
;*
;*          Dummy Data Section
;*
;*****
0000          dseg  0          ;absolute low memory
              org    0          ;(interrupt vectors)
              . . .
              END

```

BIOS Listing

```
*****
*
* This is the CP/M-86 BIOS, derived from the BIOS *
* program by disabling the "loader_bios" condi- *
* tional assembly switch. The listing has been *
* truncated on the right, but can be reproduced *
* by assembling the BIOS.A86 file provided with *
* CP/M-86. This BIOS allows CP/M-86 operation *
* with the Intel SBC 86/12 with the SBC 204 con- *
* troller. Use this BIOS, or the skeletal CBIOS *
* listed in Appendix E, as the basis for a cus- *
* tomized implementation of CP/M-86. *
* provided with CP/M-86) *
*
*****
```

```
*****
;*****
; *
; * Basic Input/Output System (BIOS) for *
; * CP/M-86 Configured for iSBC 86/12 with *
; * the iSBC 204 Floppy Disk Controller *
; *
; * (Note: this file contains both embedded *
; * tabs and blanks to minimize the list file *
; * width for printing purposes. You may wish *
; * to expand the blanks before performing *
; * major editing.) *
;*****
;
; Copyright (C) 1980,1981
; Digital Research, Inc.
; Box 579, Pacific Grove
; California, 93950
;
; (Permission is hereby granted to use
; or abstract the following program in
; the implementation of CP/M, MP/M or
; CP/NET for the 8086 or 8088 Micro-
; processor)
```

```
FFFF true equ -1
0000 false equ not true
```

```

;*****
;*
;* Loader_bios is true if assembling the
;* LOADER BIOS, otherwise BIOS is for the
;* CPM.SYS file. Blc_list is true if we
;* have a serial printer attached to BLC8538
;* Bdos_int is interrupt used for earlier
;* versions.
;*
;*****

0000 loader_bios equ false
FFFF blc_list equ true
00E0 bdos_int equ 224 ;reserved BDOS Interrupt

IF not loader_bios
;-----
;|
2500 bios_code equ 2500h
0000 ccp_offset equ 0000h
0B06 bdos_ofst equ 0B06h ;BDOS entry point
;|
ENDIF ;not loader_bios

IF loader_bios
;-----
;|
bios_code equ 1200h ;start of LDBIOS
ccp_offset equ 0003h ;base of CPMLOADER
bdos_ofst equ 0406h ;stripped BDOS entry
;|
ENDIF ;loader_bios

00DA cstc equ 0DAh ;i8251 status port
00D8 cdata equ 0D8h ; " data port

IF blc_list
;-----
;|
0041 lstc equ 41h ;2651 No. 0 on BLC8538 stat
0040 ldata equ 40h ; " " " " data
0060 blc_reset equ 60h ;reset selected USARTS on B
;|
ENDIF ;blc_list

;*****
;*
;* Intel iSBC 204 Disk Controller Ports
;*
;*****

```

```

00A0          base204          equ 0a0h          ;SBC204 assigned ad
00A0          fdc_com          equ base204+0     ;8271 FDC out comma
00A0          fdc_stat        equ base204+0     ;8271 in status
00A1          fdc_parm        equ base204+1     ;8271 out parameter
00A1          fdc_rslt        equ base204+1     ;8271 in result
00A2          fdc_rst         equ base204+2     ;8271 out reset
00A4          dmaC_adr        equ base204+4     ;8257 DMA base addr
00A5          dmaC_cont       equ base204+5     ;8257 out control
00A6          dmaC_scan       equ base204+6     ;8257 out scan cont
00A7          dmaC_sadr       equ base204+7     ;8257 out scan addr
00A8          dmaC_mode       equ base204+8     ;8257 out mode
00A8          dmaC_stat       equ base204+8     ;8257 in status
00A9          fdc_sel         equ base204+9     ;FDC select port (n
00AA          fdc_segment     equ base204+10    ;segment address re
00AF          reset_204      equ base204+15    ;reset entire inter

000A          max_retries     equ 10          ;max retries on dis
                                           ;before perm error
000D          cr              equ 0dh          ;carriage return
000A          lf              equ 0ah          ;line feed

                cseq
                org          ccpoffset
ccp:           org          bios_code

;*****
;*
;* BIOS Jump Vector for Individual Routines *
;*
;*****

2500 E93C00    jmp INIT          ;Enter from BOOT ROM or LOADER
2503 E98400    jmp WBOOT         ;Arrive here from BIOS call 0
2506 E99000    jmp CONST        ;return console keyboard status
2509 E99600    jmp CONIN        ;return console keyboard char
250C E99D00    jmp CONOUT       ;write char to console device
250F E9A500    jmp LISTOUT      ;write character to list device
2512 E9B700    jmp PUNCH        ;write character to punch device
2515 E9B400    jmp READER       ;return char from reader device
2518 E9FF00    jmp HOME         ;move to trk 00 on cur sel drive
251B E9DB00    jmp SELDSK      ;select disk for next rd/write
251E E90E01    jmp SETTRK      ;set track for next rd/write
2521 E91001    jmp SETSEC      ;set sector for next rd/write
2524 E91901    jmp SETDMA      ;set offset for user buff (DMA)
2527 E92401    jmp READ        ;read a 128 byte sector
252A E92501    jmp WRITE       ;write a 128 byte sector
252D E99100    jmp LISTST      ;return list status
2530 E90601    jmp SECTTRAN    ;xlate logical->physical sector
2533 E90F01    jmp SETDMAB     ;set seg base for buff (DMA)
2536 E91101    jmp GETSEGT     ;return offset of Mem Desc Table
2539 E99300    jmp GETIOBF     ;return I/O map byte (IOBYTE)
253C E99300    jmp SETIOBF     ;set I/O map byte (IOBYTE)

```



```

;*****
;*
;* INIT Entry Point, Differs for LDBIOS and
;* BIOS, according to "Loader_Bios" value
;*
;*****
INIT: ;print signon message and initialize hardwa
253F 8CC8 mov ax,cs ;we entered with a JMPP so
2541 8ED0 mov ss,ax ; CS: as the initial value
2543 8ED8 mov ds,ax ; DS:
2545 8EC0 mov es,ax ; and ES:
;use local stack during initialization
2547 BCE429 mov sp,offset stkbase
254A FC cld ;set forward direction

IF not loader_bios
;-----|
; This is a BIOS for the CPM.SYS file.
; Setup all interrupt vectors in low
; memory to address trap

254B 1E push ds ;save the DS register
254C B80000 mov ax,0
254F 8ED8 mov ds,ax
2551 8EC0 mov es,ax ;set ES and DS to zero
;setup interrupt 0 to address trap routine
2553 C70600008D25 mov int0_offset,offset int_trap
2559 8C0E0200 mov int0_segment,CS
255D BF0400 mov di,4
2560 BE0000 mov si,0 ;then propagate
2563 B9FE01 mov cx,510 ;trap vector to
2566 F3A5 rep movs ax,ax ;all 256 interrupts
;BDOS offset to proper interrupt
2568 C7068003060B mov bdos_offset,bdos_ofst
256E 1F pop ds ;restore the DS register

;*****
;*
;* National "BLC 8538" Channel 0 for a serial*
;* 9600 baud printer - this board uses 8 Sig-*
;* netics 2651 Usarts which have on-chip baud*
;* rate generators.
;*
;*****
256F B0FF mov al,0FFh
2571 E660 out b/c_reset,al ;reset all usarts on 8538
2573 B04E mov al,4Eh
2575 E642 out ldata+2,al ;set usart 0 in async 8 bit
2577 B03E mov al,3Eh
2579 E642 out ldata+2,al ;set usart 0 to 9600 baud
257B B037 mov al,37h
257D E643 out ldata+3,al ;enable Tx/Rx, and set up R

```

```

;|
;-----|
                ENDIF ;not loader_bios
                IF    loader_bios
;|
                ;This is a BIOS for the LOADER
                push ds ;save data segment
                mov ax,0
                mov ds,ax ;point to segment zero
                ;BDOS interrupt offset
                mov bdos_offset,bdos_ofst
                mov bdos_segment,CS ;bdos interrupt segment
                pop ds ;restore data segment
;|
;-----|
                ENDIF ;loader_bios

257F BB4427      mov bx,offset signon
2582 E86600      call pmsg ;print signon message
2585 B100        mov cl,0 ;default to dr A: on coldst
2587 E976DA      jmp ccp ;jump to cold start entry o

258A E979DA      WBOOT: jmp ccp+6 ;direct entry to CCP at com

                IF    not loader_bios
;|
;-----|
int_trap:
                cli ;block interrupts
258D FA          mov ax,cs
258E 8CC8        mov ds,ax ;get our data segment
2590 8ED8        mov bx,offset int_trp
2592 BB7927      call pmsg
2595 E85300      hlt ;hardstop
2598 F4
;|
;-----|
                ENDIF ;not loader_bios

;*****
;*
;* CP/M Character I/O Interface Routines *
;* Console is Usart (i8251a) on iSBC 86/12 *
;* at ports D8/DA *
;* *
;*****

CONST: ;console status
2599 E4DA      in al,csts
259B 2402      and al,2
259D 7402      jz const_ret
259F 0CFF      or al,255 ;return non-zero if RDA

const_ret:
25A1 C3       ret ;Receiver Data Available

```

```

25A2 E8F4FF      CONIN:          ;console input
25A5 74FB        call const
25A7 E4D8        jz CONIN        ;wait for RDA
25A9 247F        in al,cdata     ;read data and remove parit
25AB C3          and al,7fh
ret

25AC E4DA        CONOUT:         ;console output
25AE 2401        in al,csts
25B0 74FA        and al,1        ;get console status
25B2 8AC1        jz CONOUT        ;wait for TBE
25B4 E6D8        mov al,cl
25B6 C3          out cdata,al    ;Transmitter Buffer Empty
ret              ;then return data

LISTOUT:         ;list device output

IF blc_list
;-----
;|
25B7 E80700      call LISTST
25BA 74FB        jz LISTOUT      ;wait for printer not busy
25BC 8AC1        mov al,cl
25BE E640        out ldata,al   ;send char to TI 810
;|
;-----
ENDIF ;blc_list

25C0 C3          ret

LISTST:          ;poll list status

IF blc_list
;-----
;|
25C1 E441        in al,lsts
25C3 2481        and al,81h     ;look at both TXRDY and DTR
25C5 3C81        cmp al,81h
25C7 750A        jnz zero ret   ;either false, printer is b
25C9 0CFE        or al,255     ;both true, LPT is ready
;|
;-----
ENDIF ;blc_list

25CB C3          ret

PUNCH:           ;not implemented in this configuration
READER:          ;not implemented in this configuration

25CC B01A        mov al,lah
25CE C3          ret              ;return EOF for now

GETIOBF:         ;TTY: for consistency
25CF B000        mov al,0
25D1 C3          ret              ;IOBYTE not implemented

```

```

SETIOBF:
25D2 C3          ret          ;iobyte not implemented

zero_ret:
25D3 2400       and al,0
25D5 C3          ret          ;return zero in AL and flag

; Routine to get and echo a console character
; and shift it to upper case

uconecho:
25D6 E8C9FF     call CONIN   ;get a console character
25D9 50         push ax
25DA 8AC8       mov cl,al   ;save and
25DC E8C9FF     call CONOUT
25DF 58         pop ax     ;echo to console
25E0 3C61       cmp al,'a'  ;less than 'a' is ok
25E2 7206       jb uret
25E4 3C7A       cmp al,'z'  ;greater than 'z' is ok
25E6 7702       ja uret
25E8 2C20       sub al,'a'-'A' ;else shift to caps
uret:
25FA C3          ret

; utility subroutine to print messages

pmsg:
25EB 8A07       mov al,[BX] ;get next char from message
25ED 84C0       test al,al
25EF 7428       jz return  ;if zero return
25F1 8AC8       mov CL,AL
25F3 E8B6FF     call CONOUT ;print it
25F6 43         inc BX
25F7 EBF2       jmps pmsg   ;next character and loop

;*****
;*
;* Disk Input/Output Routines
;*
;*****

SELDSK:          ;select disk given by register CL
25F9 BB0000     mov bx,0000h
25FC 80F902     cmp cl,2   ;this BIOS only supports 2
25FF 7318       jnb return ;return w/ 0000 in BX if ba
2601 B080       mov al,80h
2603 80F900     cmp cl,0
2606 7502       jne sell  ;drive 1 if not zero
2608 B040       mov al,40h ;else drive is 0
260A A26928     mov sel_mask,al ;save drive select mask
;now, we need disk paramete

260D B500       mov ch,0
260F 8BD9       mov bx,cx  ;BX = word(CL)
2611 B104       mov cl,4

```

```

2613 D3E3          shl bx,cl          ;multiply drive code * 16
                  ;create offset from Disk Parameter Base
2615 81C37C28      add bx,offset dp_base
return:
2619 C3           ret

HOME:             ;move selected disk to home position (Track
mov trk,0         ;set disk i/o to track zero
261A C6066C2800   mov bx,offset hom_com
261F BB6E28        call execute
2622 E83500        call execute
2625 74F2          jz return         ;home drive and return if 0
2627 BB6A27        mov bx,offset bad_hom ;else print
262A EB8EFF        call pmsg         ;"Home Error"
262D EBEB         jmps home         ;and retrv

SETTRK:           ;set track address given by CX
262F 880E6C28     mov trk,cl        ;we only use 8 bits of trac
2633 C3           ret

SETSEC:           ;set sector number given by cx
2634 880E6D28     mov sect,cl       ;we only use 8 bits of sect
2638 C3           ret

SECTRAN:         ;translate sector CX using table at [DX]
2639 8BD9          mov bx,cx
263B 03DA          add bx,dx
263D 8A1F          mov bl,[bx]       ;add sector to tran table a
263F C3           ret                               ;get logical sector

SETDMA:          ;set DMA offset given by CX
2640 890E6528     mov dma_adr,CX
2644 C3           ret

SETDMAB:         ;set DMA segment given by CX
2645 890E6728     mov dma_seg,CX
2649 C3           ret

;
GETSEGT:         ;return address of physical memory table
264A BB7328       mov bx,offset seq_table
264D C3           ret

;*****
;*
;* All disk I/O parameters are setup: the
;* Read and Write entry points transfer one
;* sector of 128 bytes to/from the current
;* DMA address using the current disk drive
;*
;*****

READ:            mov al,12h          ;basic read sector command
264E B012         jmps r_w_common

WRITE:

```

```

2652 B00A          mov al,0ah          ;basic write sector command

r_w_common:
2654 BB6A28      mov bx,offset io_com ;point to command string
2657 884701      mov byte ptr 1[BX],al ;put command into string
;               fall into execute and return

execute: ;execute command string.
;[BX] points to length,
;         followed by Command byte,
;         followed by length-1 parameter byte

265A 891E6328    mov last_com,BX ;save command address for r
outer_retry:
;allow some retrving
265E C60662280A  mov rtrv_cnt,max_retries

retry:
2663 8B1E6328    mov BX,last_com
2667 E88900      call send_com ;transmit command to i8271
;               check status poll

266A 8B1E6328    mov BX,last_com
266E 8A4701      mov al,1[bx] ;get command op code
2671 B90008      mov cx,0800h ;mask if it will be "int re
2674 3C2C        cmp al,2ch
2676 720B        jb exec_poll ;ok if it is an interrupt t
2678 B98080      mov cx,8080h ;else we use "not command b
267B 240F        and al,0fh
267D 3C0C        cmp al,0ch ;unless there isn't
267F B000        mov al,0
2681 7736        ja exec_exit ; any result
;               poll for bits in CH,
;               toggled with bits in CL

exec_poll:
2683 E4A0        in al,fdc_stat ;read status
2685 22C5        and al,ch
2687 32C1        xor al,cl ; isolate what we want to
2689 74F8        jz exec_poll ;and loop until it is done

;               ;Operation complete,
268B E4A1        in al,fdc_rslt ; see if result code indica
268D 241E        and al,1eh
268F 7428        jz exec_exit ;no error, then exit
;               ;some type of error occurre

2691 3C10        cmp al,10h
2693 7425        je dr_nrdy ;was it a not ready drive ?
;               ;no,

dr_rdy: ; then we just retry read or write
dec rtry_cnt
2695 FE0E6228    dec rtry_cnt
2699 75C8        jnz retry ; up to 10 times

;               ;
;               ; retries do not recover from the
;               ; hard error

269B B400        mov ah,0

```

```

269D 8BD8          mov bx,ax          ;make error code 16 bits
269F 8B9F9127     mov bx,errtbl[BX]
26A3 E845FF       call pmsg          ;print appropriate message
26A6 E4D8         in al,cdata       ;flush usart receiver buffe
26A8 E82BFF       call uconecho     ;read upper case console ch
26AB 3C43         cmp al,'C'
26AD 7425         je wboot_1        ;cancel
26AF 3C52         cmp al,'R'
26B1 74AB         je outer_retrv   ;retry 10 more times
26B3 3C49         cmp al,'I'
26B5 741A         je z_ret          ;ignore error
26B7 0CFF        or al,255         ;set code for permanent err
exec_exit:
26B9 C3          ret

dr_nrdy:          ;here to wait for drive ready
                call test_ready
                jnz retrv          ;if it's ready now we are d
26BA E81A00       call test_ready
26BD 75A4         jnz retrv          ;if not ready twice in row,
26BF E81500       call test_ready
26C2 759F         jnz retrv          ;if not ready twice in row,
26C4 BB0228       mov bx,offset nrdymsg
26C7 E821FF       call pmsg ;"Drive Not Ready"

nrdy01:          call test_ready
26CA E80A00       call test_ready
26CD 74FB         jz nrdy01         ;now loop until drive ready
26CF EB92         jmps retrv        ;then go retry without decr

zret:            and al,0
                ret              ;return with no error code

wboot_1:         ;can't make it w/ a short 1
                jmp WBOOT

26D4 E9B3FE

;*****
;*
;* The i8271 requires a read status command
;* to reset a drive-not-ready after the
;* drive becomes ready
;*
;*****

test_ready:
26D7 B640          mov dh, 40h       ;proper mask if dr 1
26D9 F066692880   test sel_mask,80h
26DE 7502         jnz nrdy2
26E0 B604          mov dh, 04h       ;mask for dr 0 status bit

nrdy2:          mov bx,offset rds_com
                call send_com

dr_poll:        in al,fdc_stat    ;get status word
                test al,80h
26E8 E4A0         jnz dr_poll      ;wait for not command busy
26EA A880         in al,fdc_rslt   ;get "special result"
26EC 75FA         test al,dh       ;look at bit for this drive
26EE E4A1
26F0 84C6

```

```

2738 E4A0          in al,fdc_stat
273A AB20          test al,20h      ;test "parameter register f
273C 75FA          inz parm_poll   ;idle until parm req not fu
273E 8A07          mov al,[BX]
2740 E6A1          out fdc_parm,al ;send next parameter
2742 EBEB          jmps parm_loop  ;go see if there are more p

```

```

;*****
;*
;*          Data Areas
;*
;*****
2744          data_offset equ offset $

          dseg
          org      data_offset      ;contiguous with co
          IF      loader_bios
;-----
;|
signon db      cr,lf,cr,lf
db      "CP/M-86 Version 2.2",cr,lf,0
;|
;-----
          ENDIF ;loader_bios

          IF      not loader_bios
;-----
;|
2744 0D0A0D0A      signon db      cr,lf,cr,lf
2748 202053797374 db      " System Generated - 11 Jan 81",c
        656D2047656E
        657261746564
        20202D203131
        204A616E2038
        310D0A00
;|
;-----
          ENDIF ;not loader_bios

276A 0D0A486F6D65 bad_hom db      cr,lf,"Home Error",cr,lf,0
        204572726F72
        0D0A00
2779 0D0A496E7465 int_trp db      cr,lf,"Interrupt Trap Halt",cr,lf,0
        727275707420
        547261702048
        616C740D0A00

2791 B127B127B127 errtbl dw er0,er1,er2,er3
        B127
2799 C127D127DE27      dw er4,er5,er6,er7
        EF27
27A1 022816282828      dw er8,er9,erA,erB
        3D28
27A9 4D28B127B127      dw erC,erD,erE,erF

```


B127

```
27B1 0D0A4E756C6C er0 db cr,lf,"Null Error ??",0
204572726F72
203F3F00
27B1 er1 equ er0
27B1 er2 equ er0
27B1 er3 equ er0
27C1 0D0A436C6F63 er4 db cr,lf,"Clock Error :",0
6B204572726F
72203A00
27D1 0D0A4C617465 er5 db cr,lf,"Late DMA :",0
20444D41203A
00
27DE 0D0A49442043 er6 db cr,lf,"ID CRC Error :",0
524320457272
6F72203A00
27EF 0D0A44617461 er7 db cr,lf,"Data CRC Error :",0
204352432045
72726F72203A
00
2802 0D0A44726976 er8 db cr,lf,"Drive Not Ready :",0
65204E6F7420
526561647920
3A00
2816 0D0A57726974 er9 db cr,lf,"Write Protect :",0
652050726F74
656374203A00
2828 0D0A54726B20 erA db cr,lf,"Trk 00 Not Found :",0
3030204E6F74
20466F756E64
203A00
283D 0D0A57726974 erB db cr,lf,"Write Fault :",0
65204661756C
74203A00
284D 0D0A53656374 erC db cr,lf,"Sector Not Found :",0
6F72204E6F74
20466F756E64
203A00
27B1 erD equ er0
27B1 erE equ er0
27B1 erF equ er0
2802 nrDymsg equ er8

2862 00 rtry_cnt db 0 ;disk error retry counter
2863 0000 last_com dw 0 ;address of last command string
2865 0000 dma_adr dw 0 ;dma offset stored here
2867 0000 dma_seg dw 0 ;dma segment stored here
2869 40 sel_mask db 40h ;select mask, 40h or 80h

; Various command strings for i8271

286A 03 io_com db 3 ;length
286B 00 rd_wr db 0 ;read/write function code
286C 00 trk db 0 ;track #
```

```

26F2 C3          ret          ;return status of ready
;*****
;*
;* Send com sends a command and parameters
;* to the i8271: BX addresses parameters.
;* The DMA controller is also initialized
;* if this is a read or write
;*
;*****

send_com:
26F3 E4A0      in al,fdc_stat
26F5 A880      test al,80h      ;insure command not busy
26F7 75FA      jnz send_com    ;loop until ready

                ;see if we have to initialize for a DMA ope

26F9 8A4701    mov al,1[bx]    ;get command byte
26FC 3C12      cmp al,12h
26FE 7504      jne write_maybe ;if not a read it could be
2700 B140      mov cl,40h
2702 EB06      jmps init_dma   ;is a read command, go set

write_maybe:
2704 3C0A      cmp al,0ah
2706 7520      jne dma_exit   ;leave DMA alone if not rea
2708 B180      mov cl,80h     ;we have write, not read

init_dma:
;we have a read or write operation, setup DMA contr
;(CL contains proper direction bit)
270A B004      mov al,04h
270C E6A8      out dmac_mode,al ;enable dmac
270E B000      mov al,00
2710 E6A5      out dmac_cont,al ;send first byte to con
2712 8AC1      mov al,c1
2714 E6A5      out dmac_cont,al ;load direction register
2716 A16528    mov ax,dma_adr
2719 E6A4      out dmac_adr,al ;send low byte of DMA
271B 8AC4      mov al,ah
271D E6A4      out dmac_adr,al ;send high byte
271F A16728    mov ax,dma_seq
2722 E6AA      out fdc_segment,al ;send low byte of segmen
2724 8AC4      mov al,ah
2726 E6AA      out fdc_segment,al ;then high segment addre

dma_exit:
2728 8A0F      mov cl,[BX]    ;get count
272A 43        inc BX
272B 8A07      mov al,[BX]    ;get command
272D 0A066928  or al,sel_mask ;merge command and drive co
2731 E6A0      out fdc_com,al ;send command byte

parm_loop:
2733 FEC9      dec cl
2735 7482      jz exec_exit   ;no (more) parameters, retu
2737 43        inc BX         ;point to (next) parameter

parm_poll:

```

```

286D 00          sect    db 0          ;sector #

286E 022900     hom_com db 2,29h,0          ;home drive command
2871 012C       rds_com db 1,2ch          ;read status command

;              System Memory Segment Table

2873 02          segtable db 2          ;2 segments
2874 DF02       dw tpa_seg          ;1st seg starts after BIOS
2876 2105       dw tpa_len          ;and extends to 08000
2878 0020       dw 2000h           ;second is 20000 -
287A 0020       dw 2000h           ;3FFFF (128k)

=
;              include singles.lib ;read in disk definitio
=
;              DISKS 2
= 287C          dpbase equ          $          ;Base of Disk Param
=287C AB280000  dpe0   dw          xlt0,0000h        ;Translate Table
=2880 00000000  dw          0000h,0000h        ;Scratch Area
=2884 C5289C28  dw          dirbuf,dpb0         ;Dir Buff, Parm Blo
=2888 64294529  dw          csv0,alv0          ;Check, Alloc Vecto
=288C AB280000  dpel   dw          xlt1,0000h        ;Translate Table
=2890 00000000  dw          0000h,0000h        ;Scratch Area
=2894 C5289C28  dw          dirbuf,dpbl         ;Dir Buff, Parm Blo
=2898 93297429  dw          csv1,alv1          ;Check, Alloc Vecto

=
;              DISKDEF 0,1,26,6,1024,243,64,64,2
= 289C          dpb0   equ          offset $          ;Disk Parameter Blo
=289C 1A00     dw          26          ;Sectors Per Track
=289E 03       db          3          ;Block Shift
=289F 07       db          7          ;Block Mask
=28A0 00       db          0          ;Extnt Mask
=28A1 F200     dw          242         ;Disk Size - 1
=28A3 3F00     dw          63         ;Directory Max
=28A5 C0       db          192        ;Alloc0
=28A6 00       db          0          ;Alloc1
=28A7 1000     dw          16         ;Check Size
=28A9 0200     dw          2          ;Offset
= 28AB          xlt0   equ          offset $          ;Translate Table
=28AB 01070D13 db          1,7,13,19
=28AF 19050B11 db          25,5,11,17
=28B3 1703090F db          23,3,9,15
=28B7 1502080E db          21,2,8,14
=28BB 141A060C db          20,26,6,12
=28BF 1218040A db          18,24,4,10
=28C3 1016     db          16,22

= 001F          als0   equ          31          ;Allocation Vector
= 0010          css0   equ          16         ;Check Vector Size
=
;              DISKDEF 1,0
= 289C          dpbl   equ          dpb0         ;Equivalent Paramet
= 001F          als1   equ          als0         ;Same Allocation Ve
= 0010          css1   equ          css0         ;Same Checksum Vect
= 28AB          xlt1   equ          xlt0         ;Same Translate Tab
=
;              ENDEF
=
;              Uninitialized Scratch Memory Follows:
= 28C5          begdat equ          offset $          ;Start of Scratch A

```

```

=28C5      dirbuf  rs      128          ;Directory Buffer
=2945      alv0   rs      als0         ;Alloc Vector
=2964      csv0   rs      css0         ;Check Vector
=2974      alv1   rs      als1         ;Alloc Vector
=2993      csv1   rs      css1         ;Check Vector
= 29A3     enddat equ      offset $    ;End of Scratch Are
= 00DE     datsiz equ      offset $-begdat ;Size of Scratch Ar
=29A3 00   db      0                ;Marks End of Modul

```

```

29A4      loc_stk rw 32 ;local stack for initialization
29E4      stkbase equ offset $

```

```

29E4      lastoff equ offset $
02DF     tpa_seg equ (lastoff+0400h+15) / 16
0521     tpa_len equ 0800h - tpa_seg
29E4 00   db 0 ;fill last address for GENCMD

```

```

;*****
;*
;*          Dummy Data Section
;*
;*****
0000      dseq    0 ;absolute low memory
          org    0 ;(interrupt vectors)
0000     int0_offset rw 1
0002     int0_segment rw 1
          ; pad to system call vector
0004      rw     2*(bdos_int-1)

0380     bdos_offset rw 1
0382     bdos_segment rw 1
          END

```

CBIOS Listing

```

*****
*
* This is the listing of the skeletal CBIOS which
* you can use as the basis for a customized BIOS
* for non-standard hardware. The essential por-
* tions of the BIOS remain, with "rs" statements
* marking the routines to be inserted.
*
*****

;*****
;*
;* This Customized BIOS adapts CP/M-86 to
;* the following hardware configuration
;* Processor:
;* Brand:
;* Controller:
;*
;*
;* Programmer:
;* Revisions :
;*
;*****

FFFF      true          equ -1
0000      false        equ not true
000D      cr           equ 0dh ;carriage return
000A      lf           equ 0ah ;line feed

;*****
;*
;* Loader_bios is true if assembling the
;* LOADER BIOS, otherwise BIOS is for the
;* CPM.SYS file.
;*
;*****

0000      loader_bios  equ false
00E0      bdos_int     equ 224 ;reserved BDOS interrupt

                IF      not loader_bios

;|
2500      bios_code    equ 2500h
0000      cpm_offset   equ 0000h
0B06      bdos_ofst    equ 0B06h ;BDOS entry point
;|
;-----

```

```

                                ENDIF ;not loader_bios

                                IF loader_bios
;-----
bios_code equ 1200h ;start of LDBIOS
ccp_offset equ 0003h ;base of CPMLoader
bdos_ofst equ 0406h ;striped BDOS entry
;-----
                                ENDIF ;loader_bios

                                cseq
                                org ccpoffset

ccp:
                                org bios_code

;*****
;*
;* BIOS Jump Vector for Individual Routines *
;*
;*****

2500 E93C00 jmp INIT ;Enter from BOOT ROM or LOADER
2503 E97900 jmp WBOOT ;Arrive here from BDOS call 0
2506 E98500 jmp CONST ;return console keyboard status
2509 E98D00 jmp CONIN ;return console keyboard char
250C E99A00 jmp CONOUT ;write char to console device
250F E9A200 jmp LISTOUT ;write character to list device
2512 E9B500 jmp PUNCH ;write character to punch device
2515 E9BD00 jmp READER ;return char from reader device
2518 E9F600 jmp HOME ;move to trk 00 on-cur sel drive
251B E9D900 jmp SELDSK ;select disk for next rd/write
251E E90101 jmp SETTRK ;set track for next rd/write
2521 E90301 jmp SETSEC ;set sector for next rd/write
2524 E90C01 jmp SETDMA ;set offset for user buff (DMA)
2527 E91701 jmp READ ;read a 128 byte sector
252A E94701 jmp WRITE ;write a 128 byte sector
252D E98F00 jmp LISTST ;return list status
2530 E9F900 jmp SECTTRAN ;xlate logical->physical sector
2533 E90201 jmp SETDMAB ;set seq base for buff (DMA)
2536 E90401 jmp GETSEGT ;return offset of Mem Desc Table
2539 E9A400 jmp GETIOBF ;return I/O map byte (IOBYTE)
253C E9A500 jmp SETIOBF ;set I/O map byte (IOBYTE)

;*****
;*
;* INIT Entry Point, Differs for LDBIOS and *
;* BIOS, according to "Loader_Bios" value *
;*
;*****

253F 8CC8 INIT: ;print signon message and initialize hardwa
                                mov ax,cs ;we entered with a JMPF so

```

```

2541 8ED0          mov ss,ax          ;CS: as the initial value o
2543 8ED8          mov ds,ax          ;DS:,
2545 8EC0          mov es,ax          ;and ES:
                    ;use local stack during initialization
2547 BC5928        mov sp,offset stkbse
254A FC           cld              ;set forward direction

                    IF      not loader_bios
;-----|
;|
;|              ; This is a BIOS for the CPM.SYS file.
;|              ; Setup all interrupt vectors in low
;|              ; memory to address trap

254B 1E           push ds           ;save the DS register
254C C606A72600   mov IOBYTE,0     ;clear IOBYTE
2551 B80000        mov ax,0
2554 8ED8          mov ds,ax
2556 8EC0          mov es,ax        ;set ES and DS to zero
                    ;setup interrupt 0 to address trap routine
2558 C70600008225 mov int0_offset,offset int_trap
255E 8C0E0200     mov int0_segment,CS
2562 BF0400        mov di,4
2565 BE0000        mov si,0         ;then propagate
2568 B9FE01        mov cx,510       ;trap vector to
256B F3A5          rep movs ax,ax   ;all 256 interrupts
                    ;BDOS offset to proper interrupt
256D C7068003060B mov bdos_offset,bdos_ofst
2573 1F           pop ds           ;restore the DS register

;|              (additional CP/M-86 initialization)
;|
;-----|
                    ENDIF ;not loader_bios

                    IF      loader_bios
;-----|
;|
;|              ;This is a BIOS for the LOADER
;|              push ds           ;save data segment
;|              mov ax,0
;|              mov ds,ax        ;point to segment zero
;|              ;BDOS interrupt offset
;|              mov bdos_offset,bdos_ofst
;|              mov bdos_segment,CS ;bdos interrupt segment
;|              (additional LOADER initialization)
;|              pop ds           ;restore data segment
;|
;-----|
                    ENDIF ;loader_bios

2574 BBB126        mov bx,offset signon
2577 E86F00        call pmsg        ;print signon message
257A B100          mov cl,0         ;default to dr A: on coldst
257C E981DA        jmp ccp          ;jump to cold start entry o

```

```

257F E984DA      WBOOT:  jmp ccp+6      ;direct entry to CCP at com
                IF      not loader_bios
                ;-----|
                ;|
                int_trap:
2582 FA          cli          ;block interrupts
2583 8CC8        mov ax,cs      ;
2585 8ED8        mov ds,ax      ;get our data segment
2587 BBD126     mov bx,offset int_trp
258A E85C00     call dmsq
258D F4          hlt          ;hardstop
                ;|
                ;-----|
                ENDIF  ;not loader_bios

                ;*****
                ;*
                ;*  CP/M Character I/O Interface Routines  *
                ;*
                ;*****

                CONST:      ;console status
258E             rs          10      ;(fill-in)
2598 C3          ret

                CONIN:      ;console input
2599 E8F2FF     call CONST
259C 74FB        jz  CONIN      ;wait for RDA
259E             rs          10      ;(fill-in)
25A8 C3          ret

                CONOUT:     ;console output
25A9             rs          10      ;(fill-in)
25B3 C3          ret          ;then return data

                LISTOUT:    ;list device output
25B4             rs          10      ;(fill-in)
25BE C3          ret

                LISTST:     ;poll list status
25BF             rs          10      ;(fill-in)
25C9 C3          ret

                PUNCH:      ;write punch device
25CA             rs          10      ;(fill-in)
25D4 C3          ret

                READER:     ;
25D5             rs          10      ;(fill-in)
25DF C3          ret

                GETIOBF:    mov al,IOBYTE
25E0 A0A726

```



```

25E3 C3                ret

SETIOBF:
25E4 880EA726         mov IOBYTE,cl        ;set iobyte
25E8 C3                ret                ;iobyte not implemented

pmsg:
25E9 8A07             mov al,[BX]          ;get next char from message
25EB 84C0             test al,al
25ED 7421             jz return           ;if zero return
25EF 8AC8             mov CL,AL
25F1 E8B5FF          call CONOUT          ;print it
25F4 43               inc BX
25F5 EBF2             jmps pmsg            ;next character and loop

;*****
;*
;*           Disk Input/Output Routines           *
;*
;*****

SELDSK:                ;select disk given by register CL
ndisks equ 2 ;number of disks (up to 16)
25F7 880EA826         mov disk,cl          ;save disk number
25FB BB0000           mov bx,0000h         ;ready for error return
25FE 80F902           cmp cl,ndisks        ;n beyond max disks?
2601 730D             jnb return           ;return if so
2603 B500             mov ch,0             ;double(n)
2605 8BD9             mov bx,cx            ;bx = n
2607 B104             mov cl,4             ;ready for *16
2609 D3E3             shl bx,cl            ;n = n * 16
260B B9F126           mov cx,offset dpbase
260E 03D9             add bx,cx            ;dpbase + n * 16
2610 C3                return: ret          ;bx = .dph

HOME:                  ;move selected disk to home position (Track
mov trk,0              ;set disk i/o to track zero
2617 rs 10              ;(fill-in)
2621 C3                ret

SETTRK:                ;set track address given by CX
mov trk,CX
2622 890EA926         mov trk,CX
2626 C3                ret

SETSEC:                ;set sector number given by cx
mov sect,CX
2627 890EAB26         mov sect,CX
262B C3                ret

SECTRAN:               ;translate sector CX using table at [DX]
mov bx,cx
262E 03DA             add bx,dx            ;add sector to tran table a
2630 8A1F             mov bl,[bx]          ;get logical sector
2632 C3                ret

SETDMA:                ;set DMA offset given by CX

```

```

2633 890EAD26      mov dma_adr,CX
2637 C3           ret

SETDMAB: ;set DMA segment given by CX
2638 890EAF26      mov dma_seg,CX
263C C3           ret

;
GETSEGT: ;return address of physical memory table
263D BBE826        mov bx,offset seg_table
2640 C3           ret

;*****
;*
;* All disk I/O parameters are setup:
;* DISK is disk number (SELDISK)
;* TRK is track number (SETTRK)
;* SECT is sector number (SETSEC)
;* DMA_ADR is the DMA offset (SETDMA)
;* DMA_SEG is the DMA segment (SETDMAB)
;* READ reads the selected sector to the DMA
;* address, and WRITE writes the data from
;* the DMA address to the selected sector
;* (return 00 if successful, 01 if perm err)
;*
;*****

2641 READ:        rs      50      ;fill-in
2673 C3           ret

2674 WRITE:       rs      50      ;(fill-in)
26A6 C3           ret

;*****
;*
;* Data Areas
;*
;*****
26A7 data_offset equ offset $

      dseg
      org      data_offset ;contiguous with co
26A7 00 IOBYTE db      0
26A8 00 disk db      0 ;disk number
26A9 0000 trk dw      0 ;track number
26AB 0000 sect dw      0 ;sector number
26AD 0000 dma_adr dw     0 ;DMA offset from DS
26AF 0000 dma_seg dw     0 ;DMA Base Segment

      IF loader_bios
;-----
;|
signon db cr,lf,cr,lf

```

```

                                db      'CP/M-86 Version 1.0',cr,lf,0
;|
;-----
                                ENDIF ;loader_bios
                                IF     not loader_bios
;-----
;|
26B1 0D0A0D0A                signon db      cr,lf,cr,lf
26B5 53797374656D            db      'System Generated 00/00/00'
                                2047656E6572
                                617465642030
                                302F30302F30
                                30
26CE 0D0A00                  db      cr,lf,0
;|
;-----
                                ENDIF ;not loader_bios

26D1 0D0A                    int_trp db      cr,lf
26D3 496E74657272            db      'Interrupt Trap Halt'
                                757074205472
                                61702048616C
                                74
26E6 0D0A                    db      cr,lf
;
                                System Memory Segment Table

26E8 02                      seqtable db 2 ;2 segments
26E9 C602                    dw tpa_seg ;1st seg starts after BIOS
26EB 3A05                    dw tpa_len ;and extends to 08000
26ED 0020                    dw 2000h ;second is 20000 -
26EF 0020                    dw 2000h ;3FFFF (128k)

=
                                include singles.lib ;read in disk definitio
=
                                ; DISKS 2
= 26F1                        dpbase equ      S ;Base of Disk Param
=26F1 20270000                dpe0 dw      x1t0,0000h ;Translate Table
=26F5 00000000                dw      0000h,0000h ;Scratch Area
=26F9 3A271127                dw      dirbuf,dpb0 ;Dir Buff, Parm Blo
=26FD D927BA27                dw      csv0,alv0 ;Check, Alloc Vecto
=2701 20270000                dpel dw      x1t1,0000h ;Translate Table
=2705 00000000                dw      0000h,0000h ;Scratch Area
=2709 3A271127                dw      dirbuf,dpb1 ;Dir Buff, Parm Blo
=270D 0828E927                dw      csv1,alv1 ;Check, Alloc Vecto
=
                                ; DISKDEF 0,1,26,6,1024,243,64,64,2
= 2711                        dpb0 equ      offset S ;Disk Parameter Blo
=2711 1A00                    dw      26 ;Sectors Per Track
=2713 03                      db      3 ;Block Shift
=2714 07                      db      7 ;Block Mask
=2715 00                      db      0 ;Extnt Mask
=2716 F200                    dw      242 ;Disk Size - 1
=2718 3F00                    dw      63 ;Directory Max
=271A C0                      db      192 ;Alloc0
=271B 00                      db      0 ;Alloc

```


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